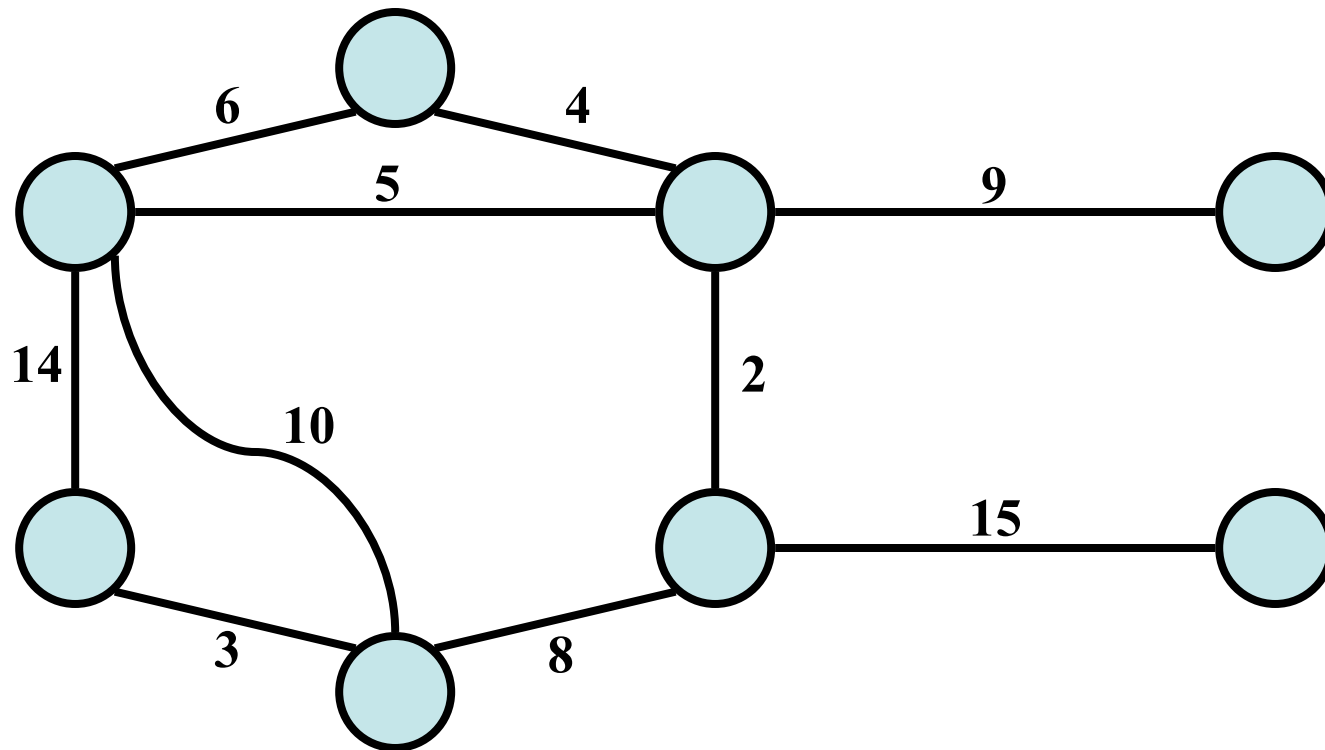


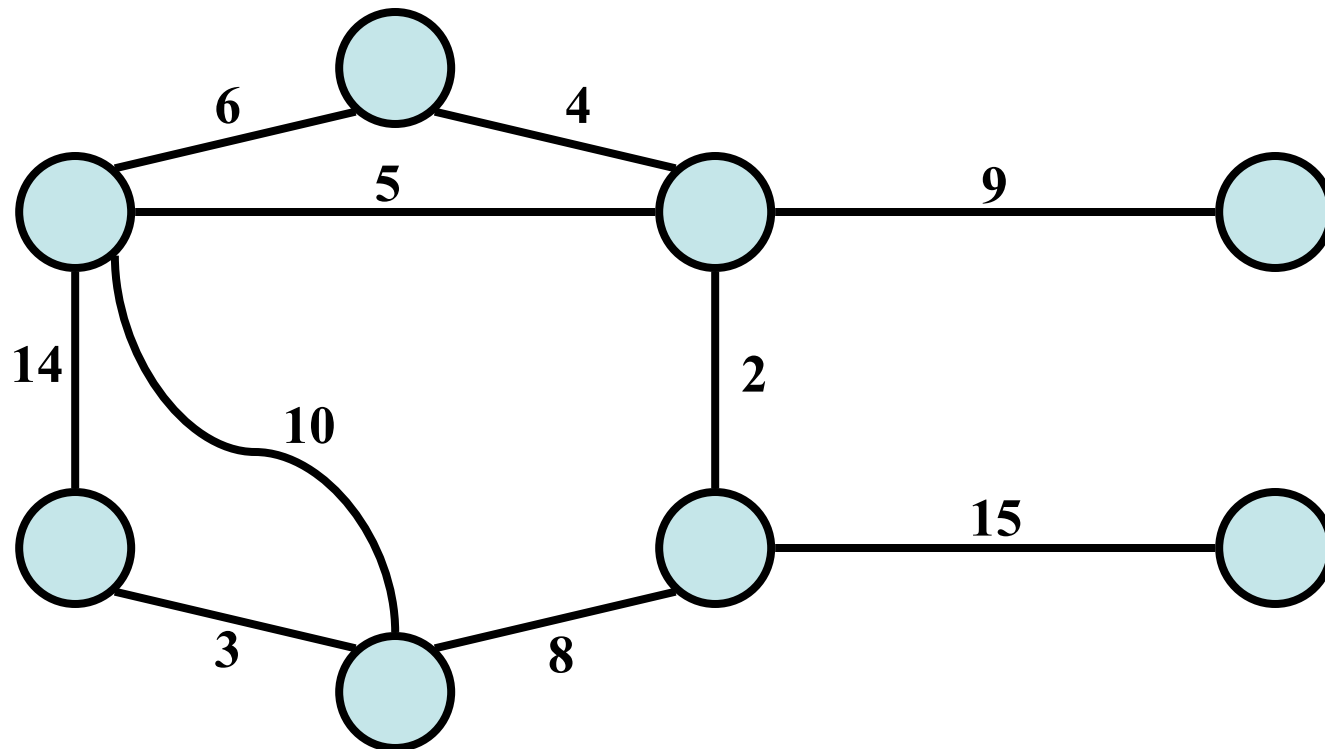
# Minimum Spanning Tree

- Problem: given a connected, undirected, weighted graph:



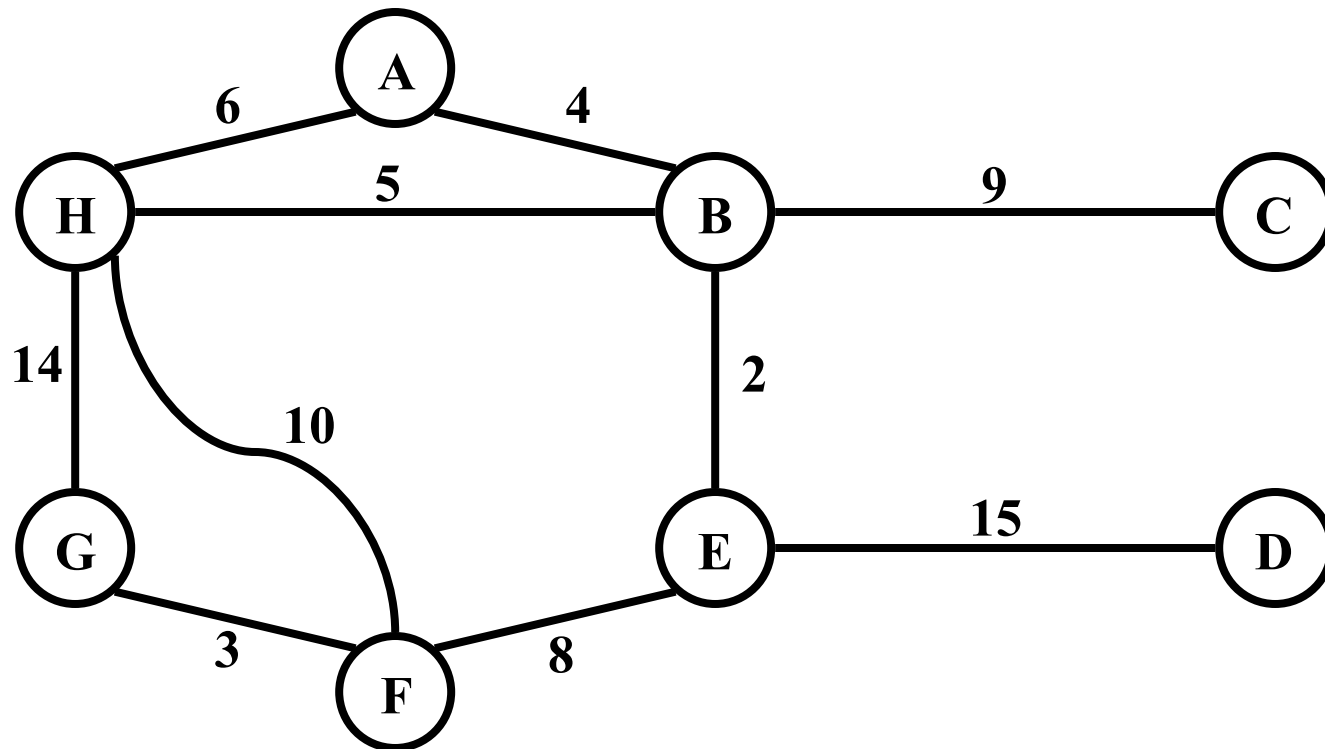
# Minimum Spanning Tree

- Problem: given a connected, undirected, weighted graph, find a *spanning tree* using edges that minimize the total weight



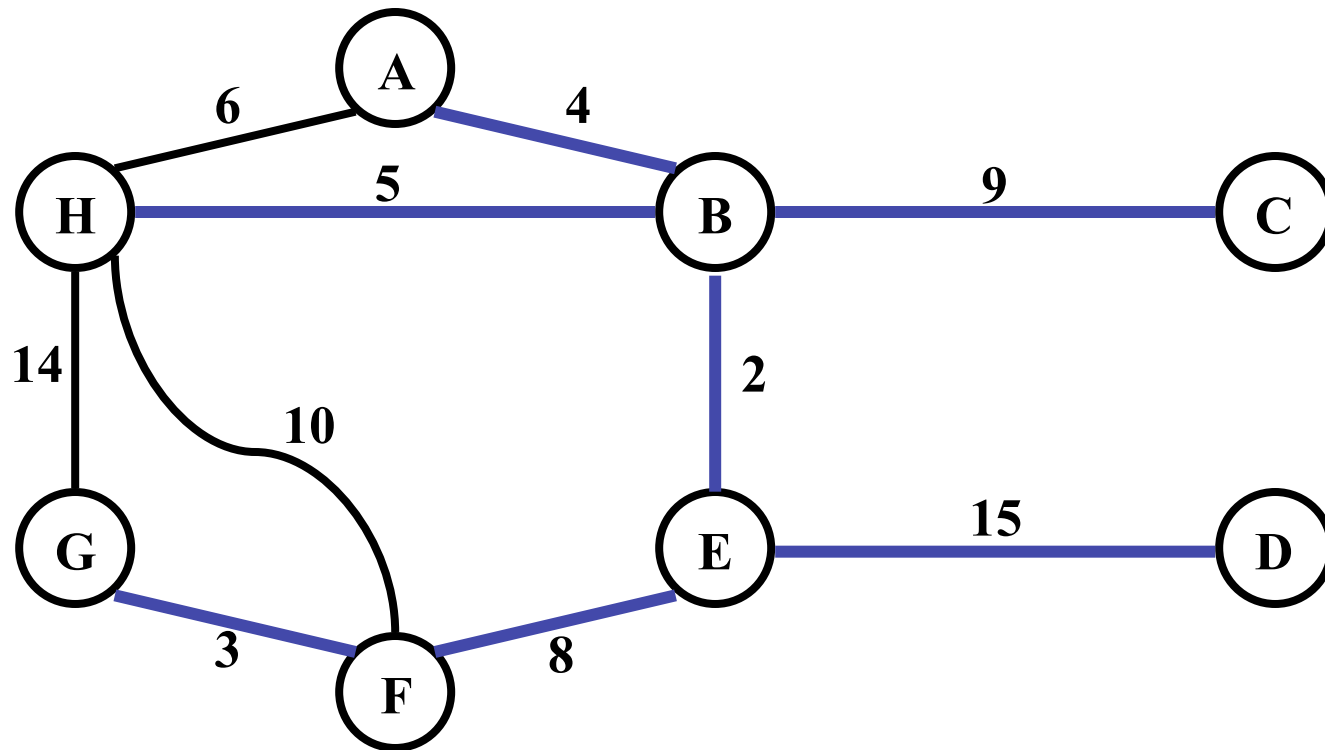
# Minimum Spanning Tree

- Which edges form the minimum spanning tree (MST) of the below graph?



# Minimum Spanning Tree

- Answer:



# Minimum Spanning Tree

- MSTs satisfy the *optimal substructure* property: an optimal tree is composed of optimal subtrees
- Let  $T$  be an MST of  $G$  with an edge  $(u,v)$  in the middle  
Removing  $(u,v)$  partitions  $T$  into two trees  $T_1$  and  $T_2$ 
  - Claim:  $T_1$  is an MST of  $G_1 = (V_1, E_1)$ , and  $T_2$  is an MST of  $G_2 = (V_2, E_2)$  (*Do  $V_1$  and  $V_2$  share vertices? Why?*)
  - Proof:  $w(T) = w(u,v) + w(T_1) + w(T_2)$   
(There can't be a better tree than  $T_1$  or  $T_2$ , or  $T$  would be suboptimal)

# Minimum Spanning Tree

- Thm:

Let  $T$  be MST of  $G$ , and let  $A \subseteq T$  be subtree of  $T$

Let  $(u,v)$  be min-weight edge connecting  $A$  to  $V-A$

Then  $(u,v) \in T$

# Minimum Spanning Tree

- Thm:  
Let  $T$  be MST of  $G$ , and let  $A \subseteq T$  be subtree of  $T$   
Let  $(u,v)$  be min-weight edge connecting  $A$  to  $V-A$   
Then  $(u,v) \in T$
- Proof: in book (see Thm 23.1)

# Prim's Algorithm

```
MST-Prim( $G, w, r$ )
   $Q = V[G];$ 
  for each  $u \in Q$ 
     $key[u] = \infty;$ 
   $key[r] = 0;$ 
   $p[r] = \text{NULL};$ 
  while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q);$ 
    for each  $v \in \text{Adj}[u]$ 
      if ( $v \in Q$  and  $w(u, v) < key[v]$ )
         $p[v] = u;$ 
         $key[v] = w(u, v);$ 
```



# Prim's Algorithm

```
MST-Prim( $G, w, r$ )
```

```
   $Q = V[G];$ 
```

```
  for each  $u \in Q$ 
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```
     $key[u] = \infty;$ 
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   $key[r] = 0;$ 
```

```
   $p[r] = \text{NULL};$ 
```

```
  while ( $Q$  not empty)
```

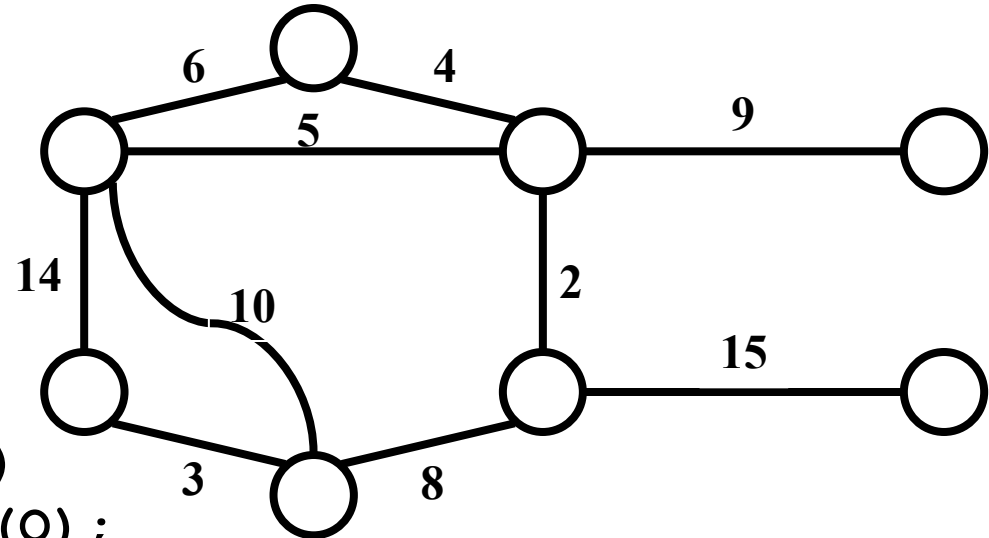
```
     $u = \text{ExtractMin}(Q);$ 
```

```
    for each  $v \in \text{Adj}[u]$ 
```

```
      if ( $v \in Q$  and  $w(u, v) < key[v]$ )
```

```
         $p[v] = u;$ 
```

```
         $key[v] = w(u, v);$ 
```



# Prim's Algorithm

```
MST-Prim( $G, w, r$ )
```

```
   $Q = V[G];$ 
```

```
  for each  $u \in Q$ 
```

```
     $key[u] = \infty;$ 
```

```
   $key[r] = 0;$ 
```

```
   $p[r] = \text{NULL};$ 
```

```
  while ( $Q$  not empty)
```

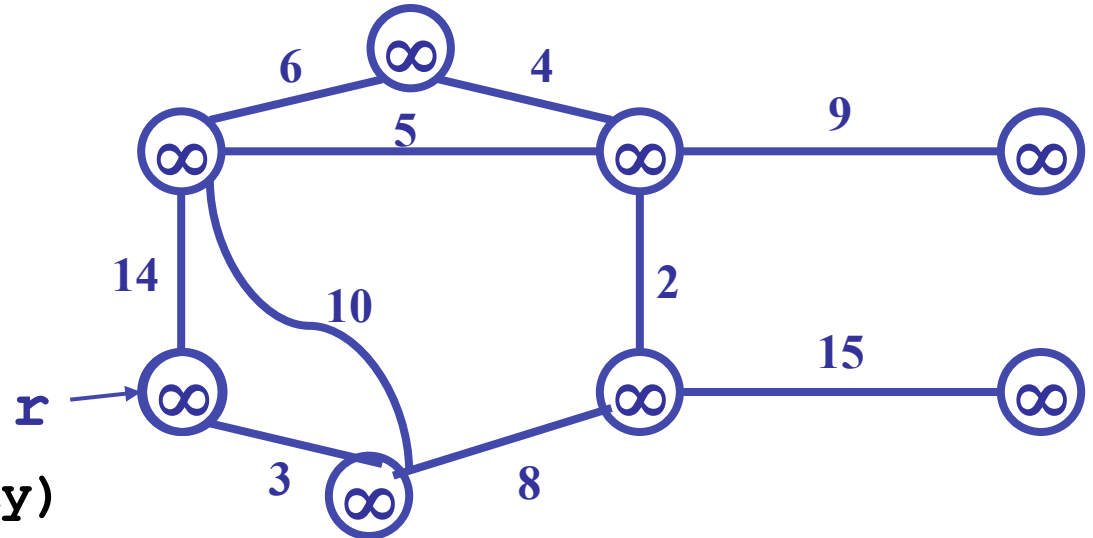
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     $u = \text{ExtractMin}(Q);$ 
```

```
    for each  $v \in \text{Adj}[u]$ 
```

```
      if ( $v \in Q$  and  $w(u, v) < key[v]$ )
```

```
         $p[v] = u;$ 
```

```
         $key[v] = w(u, v);$ 
```



Pick a start vertex  $r$

# Prim's Algorithm

```
MST-Prim( $G, w, r$ )
```

```
   $Q = V[G];$ 
```

```
  for each  $u \in Q$ 
```

```
     $key[u] = \infty;$ 
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   $key[r] = 0;$ 
```

```
   $p[r] = \text{NULL};$ 
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```
  while ( $Q$  not empty)
```

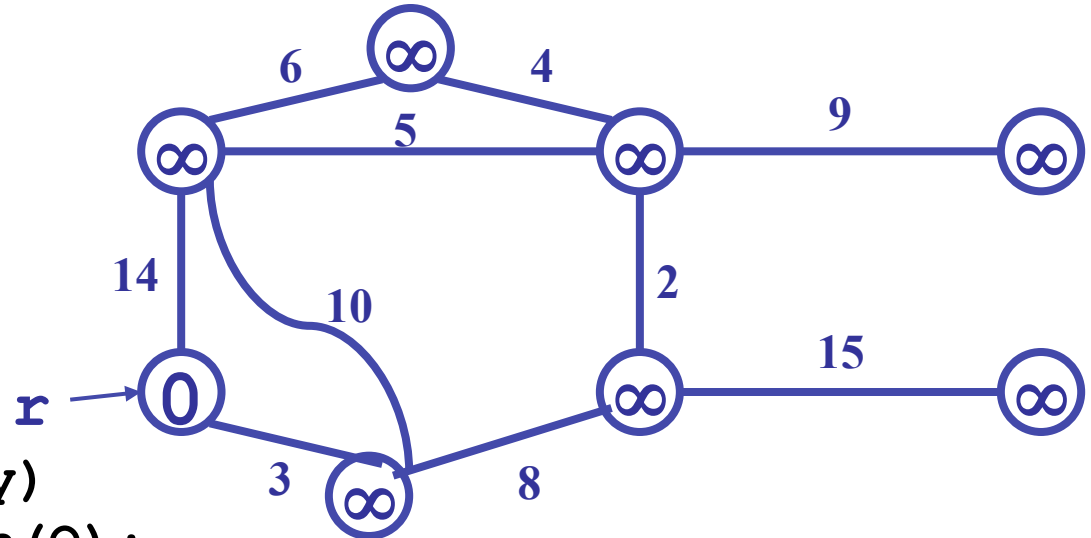
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     $u = \text{ExtractMin}(Q);$ 
```

```
    for each  $v \in \text{Adj}[u]$ 
```

```
      if ( $v \in Q$  and  $w(u, v) < key[v]$ )
```

```
         $p[v] = u;$ 
```

```
         $key[v] = w(u, v);$ 
```



Pick a start vertex  $r$

# Prim's Algorithm

```
MST-Prim( $G, w, r$ )
```

```
   $Q = V[G];$ 
```

```
  for each  $u \in Q$ 
```

```
     $key[u] = \infty;$ 
```

```
   $key[r] = 0;$ 
```

```
   $p[r] = \text{NULL};$ 
```

```
  while ( $Q$  not empty)
```

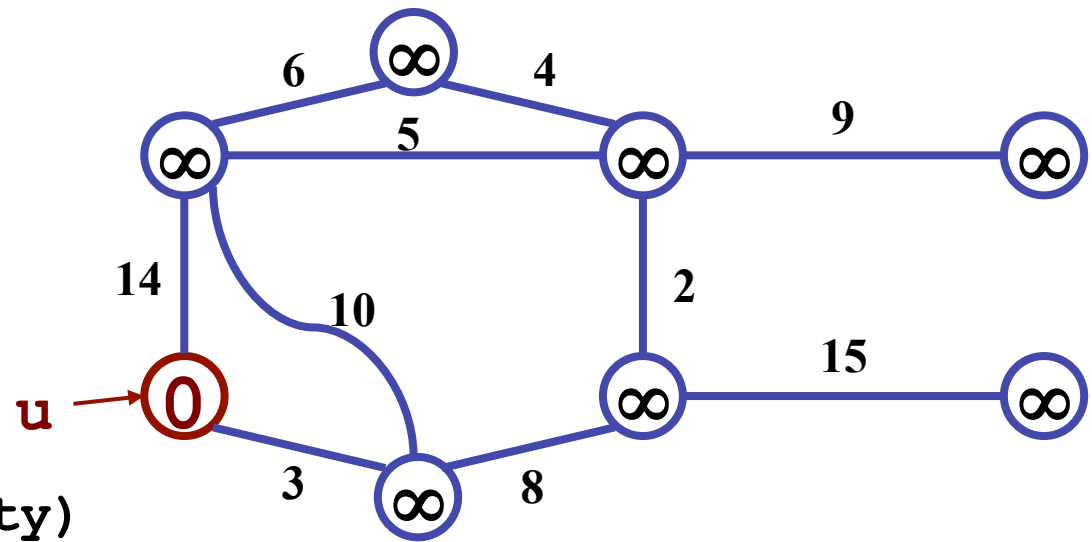
```
     $u = \text{ExtractMin}(Q);$ 
```

```
    for each  $v \in \text{Adj}[u]$ 
```

```
      if ( $v \in Q$  and  $w(u, v) < key[v]$ )
```

```
         $p[v] = u;$ 
```

```
         $key[v] = w(u, v);$ 
```



**Red vertices have been removed from Q**

# Prim's Algorithm

```
MST-Prim( $G, w, r$ )
```

```
   $Q = V[G];$ 
```

```
  for each  $u \in Q$ 
```

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     $key[u] = \infty;$ 
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   $key[r] = 0;$ 
```

```
   $p[r] = \text{NULL};$ 
```

```
  while ( $Q$  not empty)
```

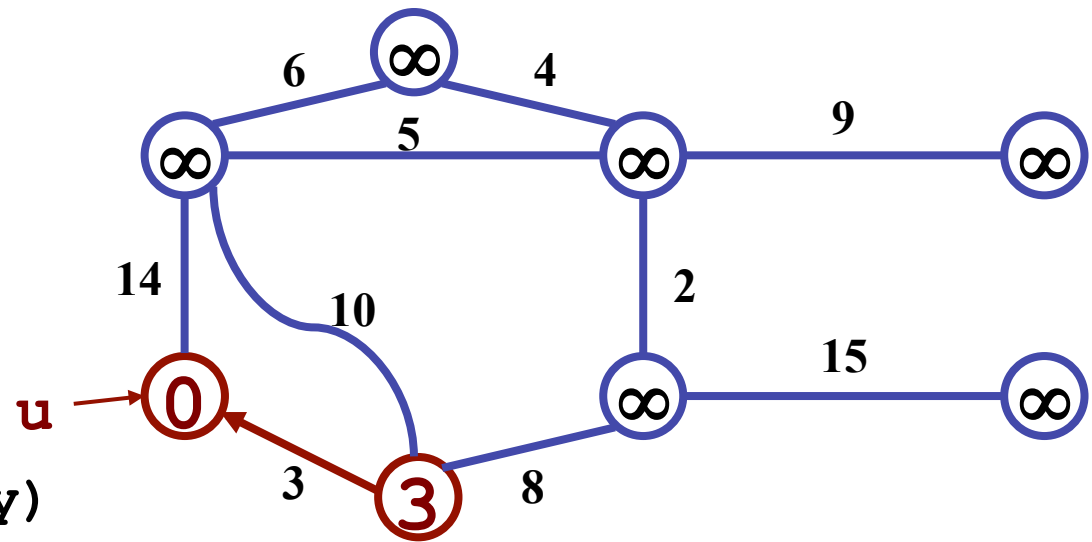
```
     $u = \text{ExtractMin}(Q);$ 
```

```
    for each  $v \in \text{Adj}[u]$  Red arrows indicate parent pointers
```

```
      if ( $v \in Q$  and  $w(u, v) < key[v]$ )
```

```
         $p[v] = u;$ 
```

```
         $key[v] = w(u, v);$ 
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   $p[r] = \text{NULL};$ 
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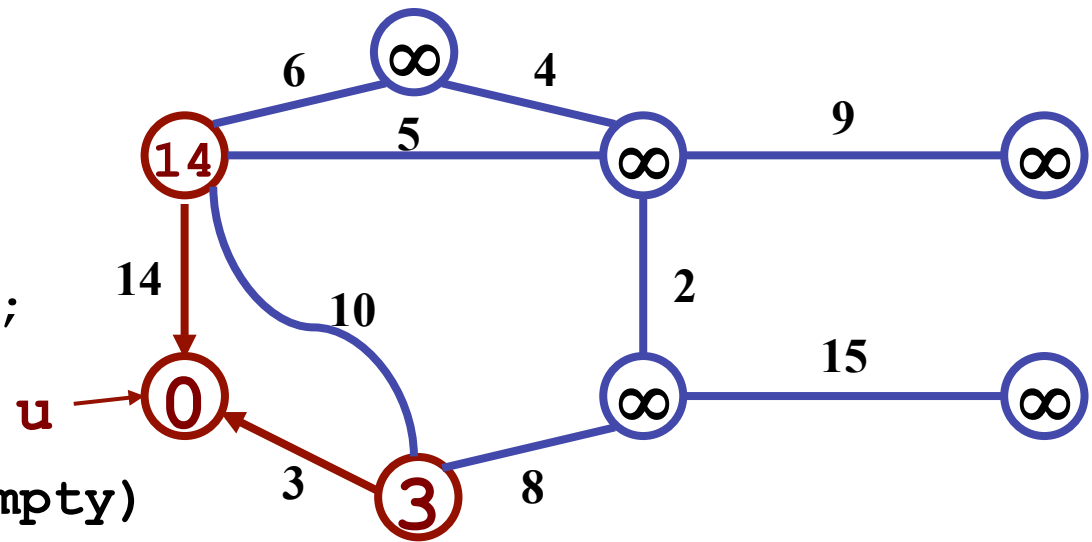
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```
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      if ( $v \in Q$  and  $w(u, v) < key[v]$ )
```

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         $p[v] = u;$ 
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```
         $key[v] = w(u, v);$ 
```









# Prim's Algorithm

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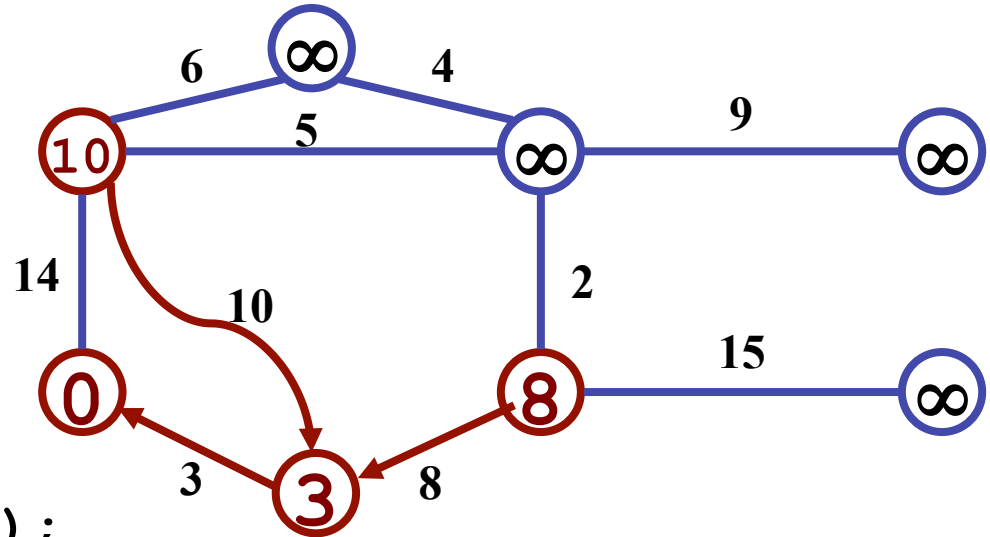
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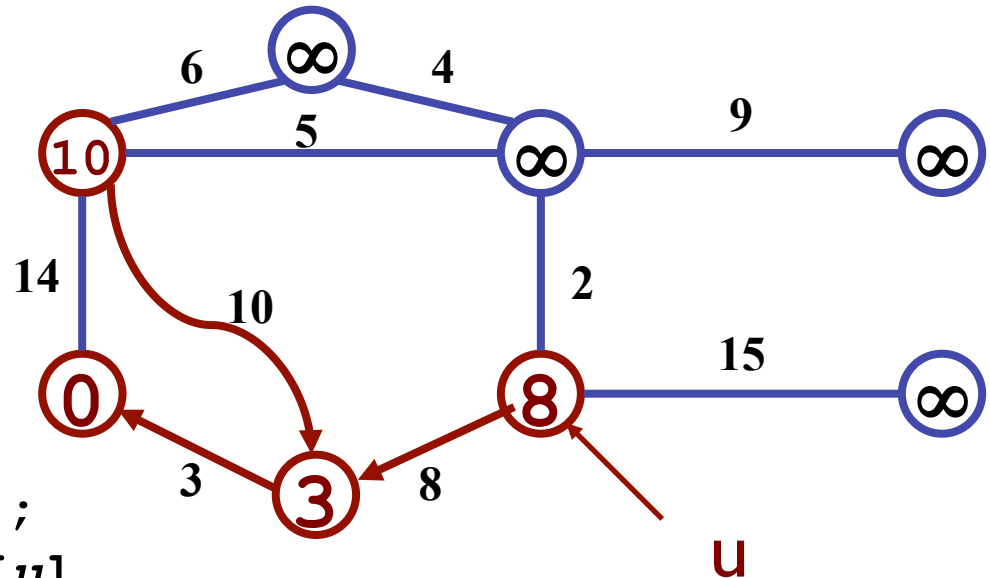
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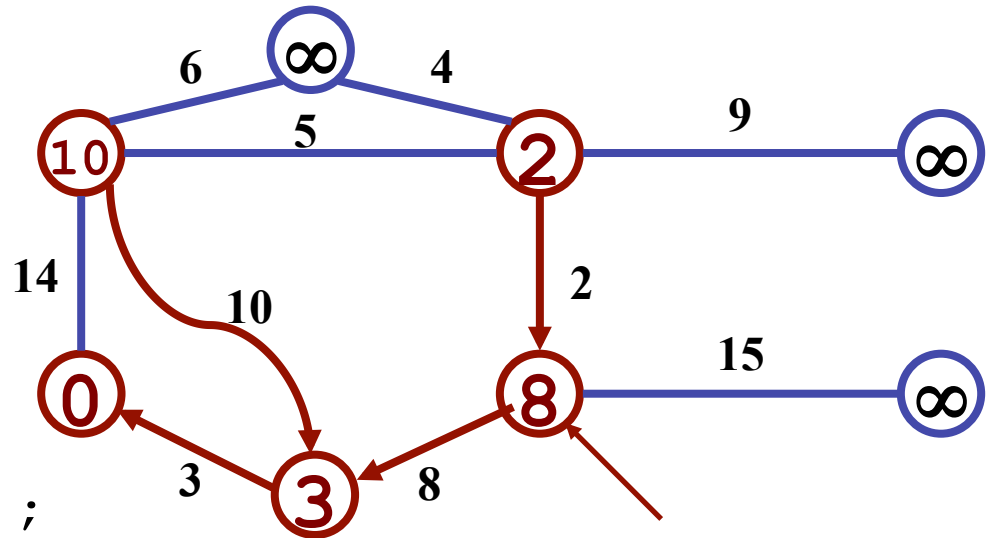
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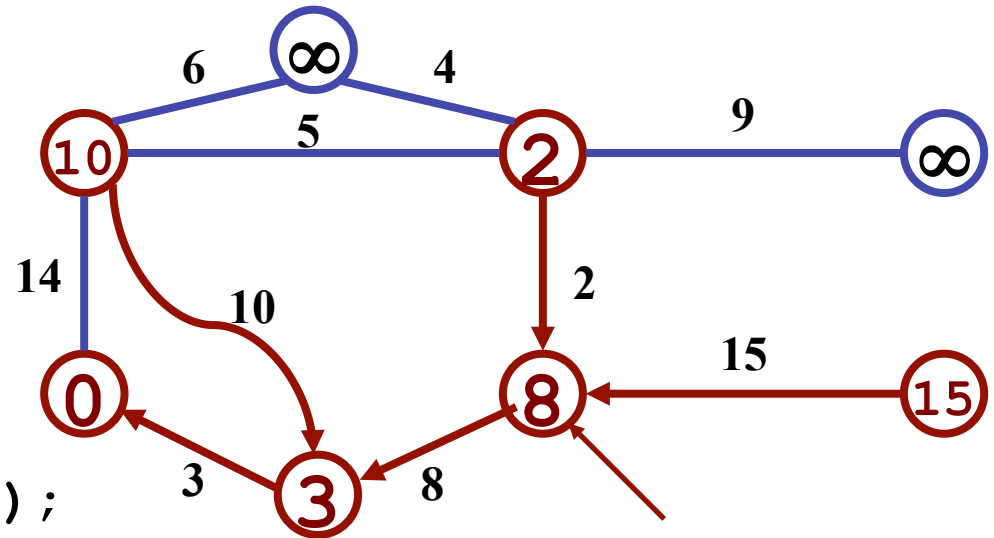
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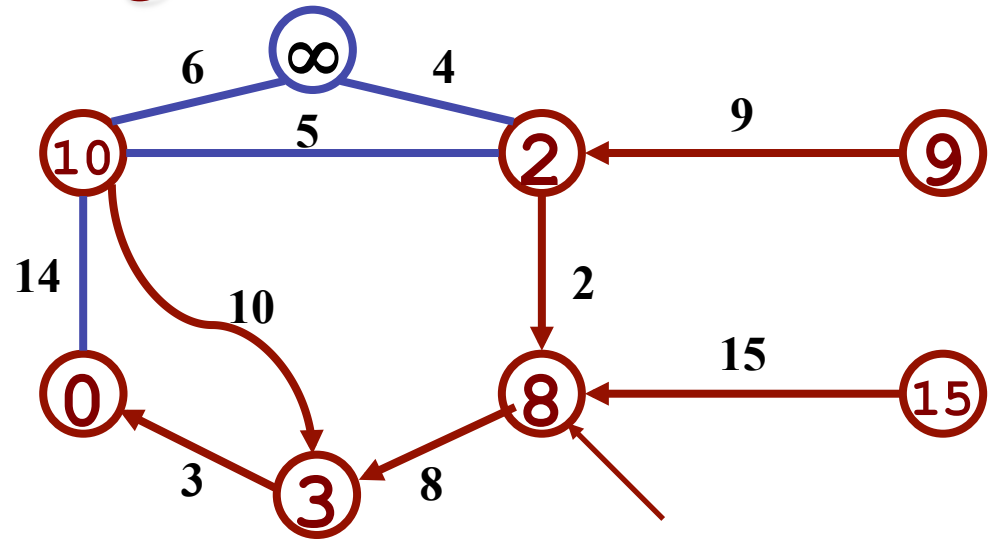
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# Prim's Algorithm

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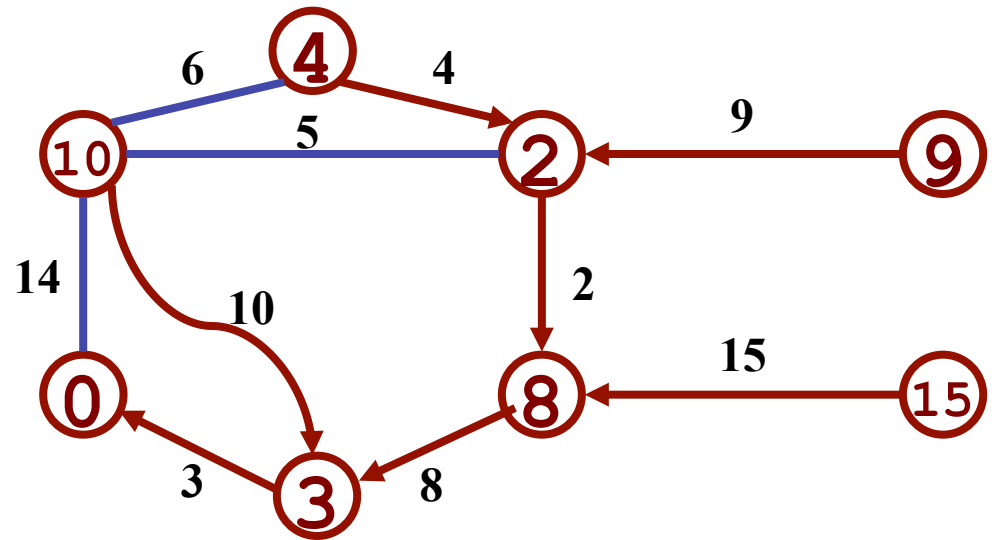
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# Prim's Algorithm

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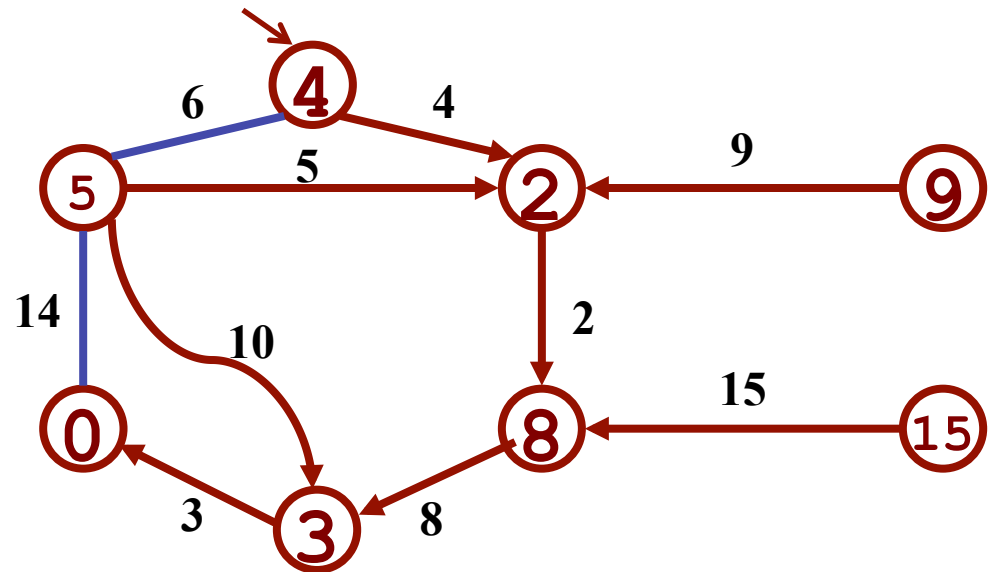
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    for each  $v \in \text{Adj}[u]$ 
```

```
      if ( $v \in Q$  and  $w(u, v) < key[v]$ )
```

```
         $p[v] = u;$ 
```

```
         $key[v] = w(u, v);$ 
```



# Prim's Algorithm

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   $Q = V[G];$ 
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```
  for each  $u \in Q$ 
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     $key[u] = \infty;$ 
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   $key[r] = 0;$ 
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```
   $p[r] = \text{NULL};$ 
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  while ( $Q$  not empty)
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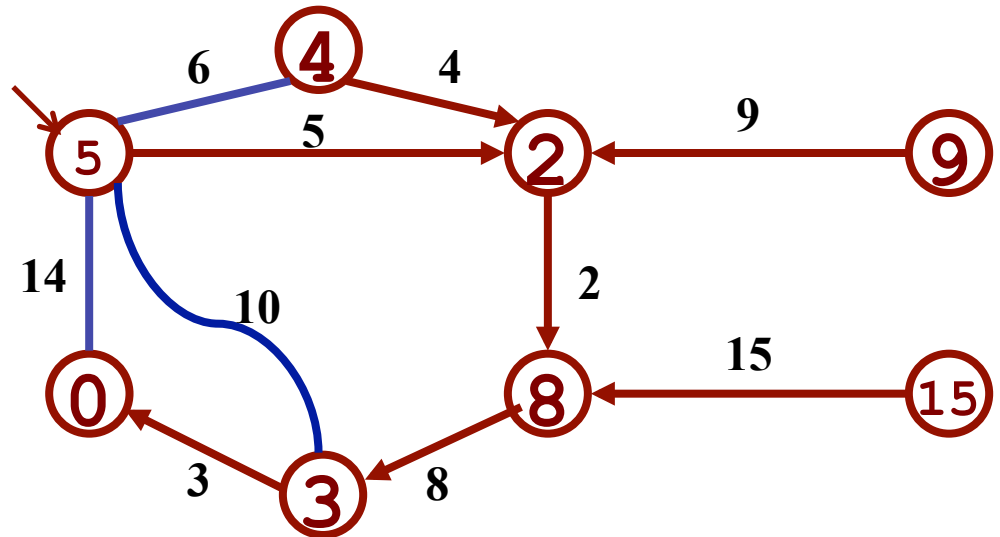
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    for each  $v \in \text{Adj}[u]$ 
```

```
      if ( $v \in Q$  and  $w(u, v) < key[v]$ )
```

```
         $p[v] = u;$ 
```

```
         $key[v] = w(u, v);$ 
```





# Prim's Algorithm

```
MST-Prim( $G, w, r$ )
```

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   $Q = V[G];$ 
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  for each  $u \in Q$ 
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```

```
   $p[r] = \text{NULL};$ 
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```
  while ( $Q$  not empty)
```

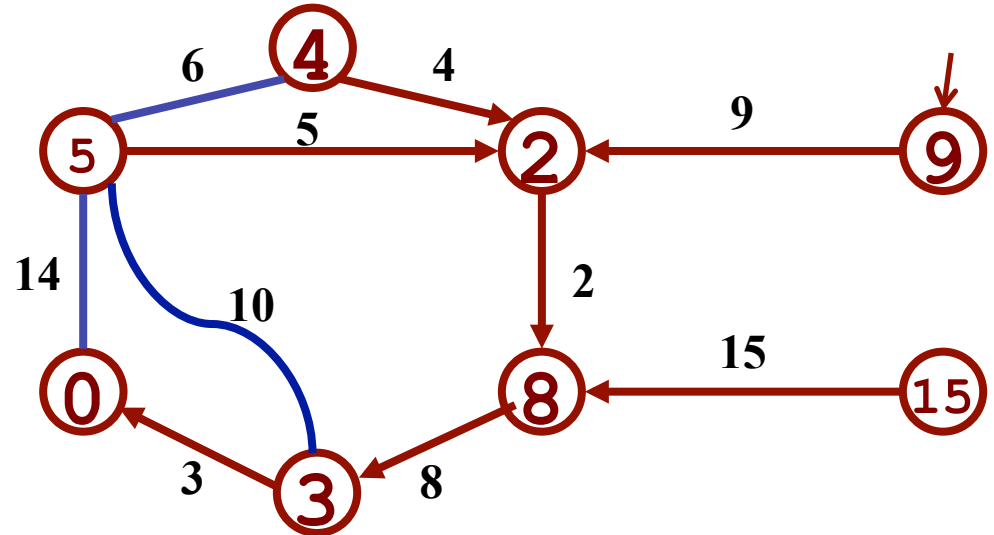
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     $u = \text{ExtractMin}(Q);$ 
```

```
    for each  $v \in \text{Adj}[u]$ 
```

```
      if ( $v \in Q$  and  $w(u, v) < key[v]$ )
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```
         $p[v] = u;$ 
```

```
         $key[v] = w(u, v);$ 
```



# Prim's Algorithm

```
MST-Prim(G, w, r)
```

```
  Q = V[G];
```

```
  for each  $u \in Q$ 
```

```
    key[u] =  $\infty$ ;
```

```
  key[r] = 0;
```

```
  p[r] = NULL;
```

```
  while (Q not empty)
```

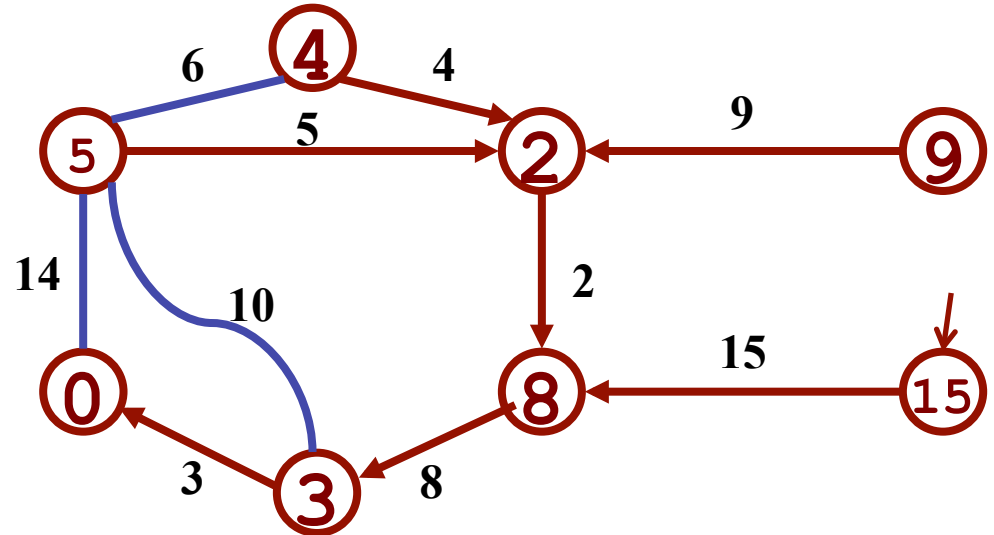
```
    u = ExtractMin(Q);
```

```
    for each  $v \in \text{Adj}[u]$ 
```

```
      if ( $v \in Q$  and  $w(u, v) < \text{key}[v]$ )
```

```
        p[v] = u;
```

```
        key[v] =  $w(u, v)$ ;
```



# Review: Prim's Algorithm

```
MST-Prim( $G, w, r$ )
   $Q = V[G];$ 
  for each  $u \in Q$ 
     $key[u] = \infty;$ 
   $key[r] = 0;$ 
   $p[r] = \text{NULL};$ 
  while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q);$ 
    for each  $v \in \text{Adj}[u]$ 
      if ( $v \in Q$  and  $w(u, v) < key[v]$ )
         $p[v] = u;$ 
         $\text{DecreaseKey}(v, w(u, v));$ 
```

# Review: Prim's Algorithm

```
MST-Prim(G, w, r)
```

```
  Q = V[G];
```

```
  for each  $u \in Q$ 
```

```
    key[u] =  $\infty$ ;
```

**How often is ExtractMin() called?**

```
  key[r] = 0;
```

**How often is DecreaseKey() called?**

```
  p[r] = NULL;
```

```
  while (Q not empty)
```

```
    u = ExtractMin(Q);
```

```
    for each  $v \in \text{Adj}[u]$ 
```

```
      if ( $v \in Q$  and  $w(u, v) < \text{key}[v]$ )
```

```
        p[v] = u;
```

```
        DecreaseKey(v, w(u, v));
```

ExtractMin total number of calls  $O(V \log V)$

DecreaseKey total number of calls  $O(E \log V)$

# Review: Prim's Algorithm

```
MST-Prim(G, w, r)
  Q = V[G];
  for each u ∈ Q
    key[u] = ∞;
  key[r] = 0;
  p[r] = NULL;
  while (Q not empty)
    u = ExtractMin(Q);
    for each v ∈ Adj[u]
      if (v ∈ Q and w(u, v) < key[v])
        p[v] = u;
        key[v] = w(u, v);
```

**What will be the running time?**

**A: Depends on queue**

**binary heap:  $O(E \lg V)$**

**Fibonacci heap:  $O(V \lg V + E)$**

ExtractMin total number of calls  $O(V \log V)$

DecreaseKey total number of calls  $O(E \log V)$

Total number of calls  $O(V \log V + E \log V) = O(E \log V)$

Think why we can combine things in the expression above

# Kruskal's Algorithm

Run the algorithm:

```
Kruskal ()
```

```
{
```

```
  T =  $\emptyset$ ;
```

```
  for each  $v \in V$ 
```

```
    MakeSet(v);
```

```
  sort E by increasing edge weight w
```

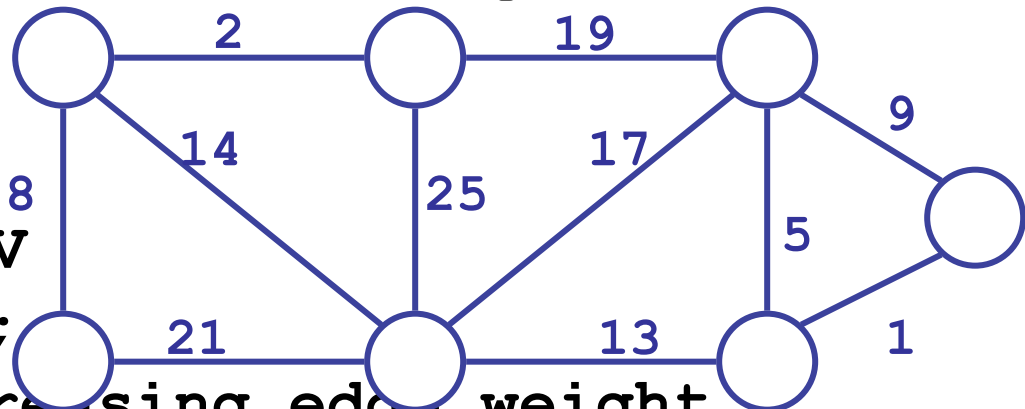
```
  for each  $(u,v) \in E$  (in sorted order)
```

```
    if FindSet(u)  $\neq$  FindSet(v)
```

```
      T = T  $\cup$  {{u,v}};
```

```
      Union(FindSet(u), FindSet(v));
```

```
}
```



# Kruskal's Algorithm

Run the algorithm:

```
Kruskal ()
```

```
{
```

```
  T =  $\emptyset$ ;
```

```
  for each v  $\in$  V
```

```
    MakeSet(v);
```

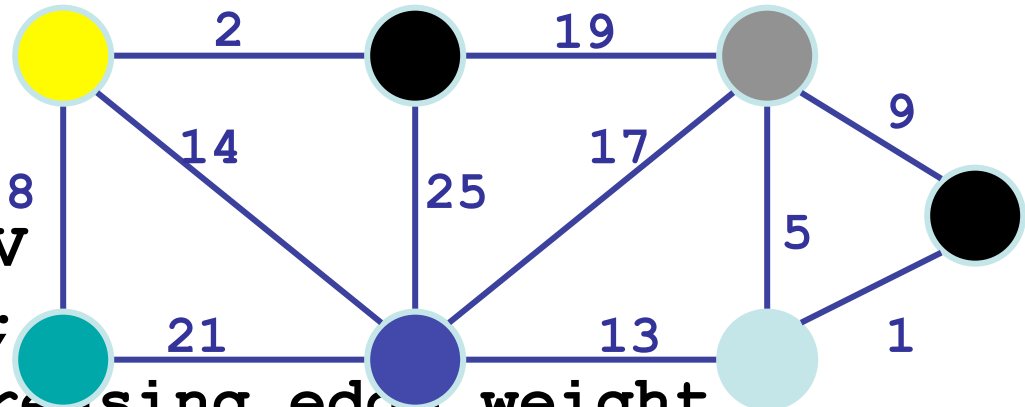
```
  sort E by increasing edge weight w  
  for each (u,v)  $\in$  E (in sorted order)
```

```
    if FindSet(u)  $\neq$  FindSet(v)
```

```
      T = T  $\cup$  {(u,v)};
```

```
      Union(FindSet(u), FindSet(v));
```

```
}
```



# Kruskal's Algorithm

Run the algorithm:

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Kruskal ()
```

```
{
```

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  T =  $\emptyset$ ;
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  for each  $v \in V$ 
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    MakeSet(v);
```

```
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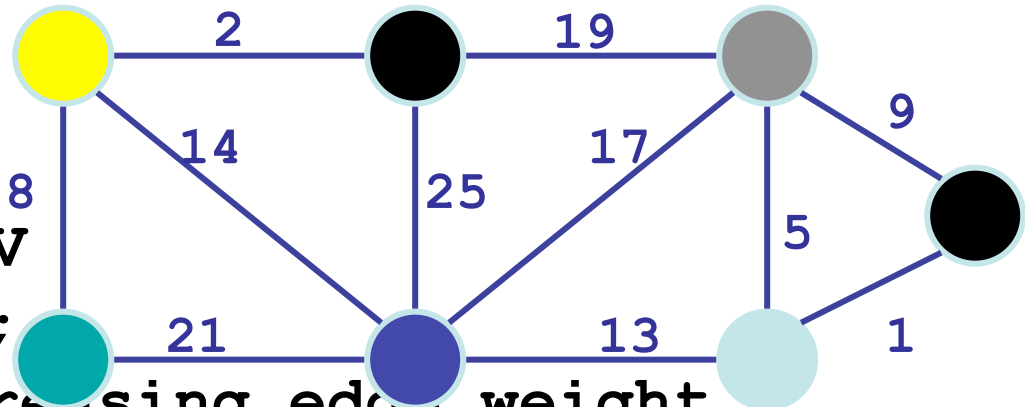
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      T = T  $\cup$  {{u,v}};
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      Union(FindSet(u), FindSet(v));
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# Kruskal's Algorithm

Run the algorithm:

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Kruskal()
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  for each  $v \in V$ 
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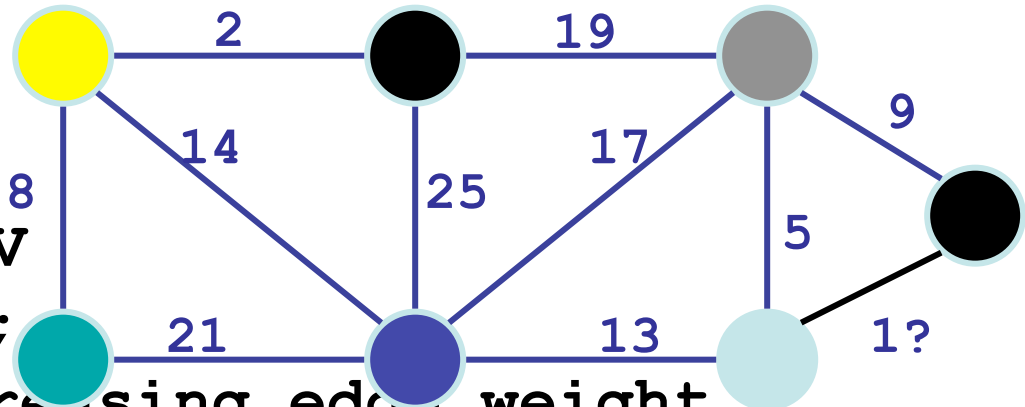
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  sort E by increasing edge weight w  
  for each  $(u,v) \in E$  (in sorted order)
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    if FindSet(u)  $\neq$  FindSet(v)
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      T = T U {{u,v}};
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      Union(FindSet(u), FindSet(v));
```

```
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```



# Kruskal's Algorithm

Run the algorithm:

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Kruskal()
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  for each  $v \in V$ 
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    MakeSet(v);
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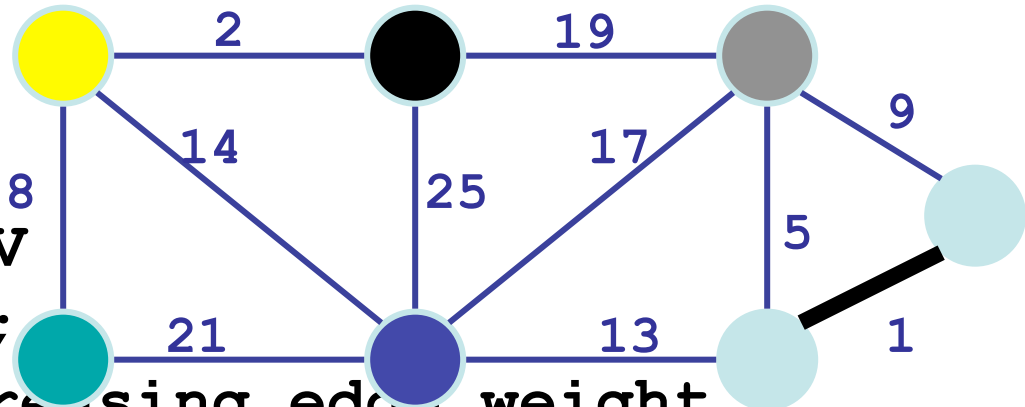
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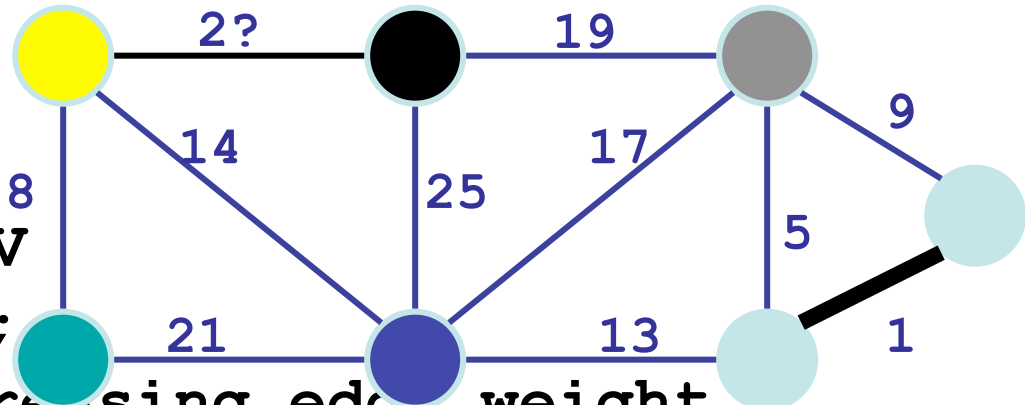
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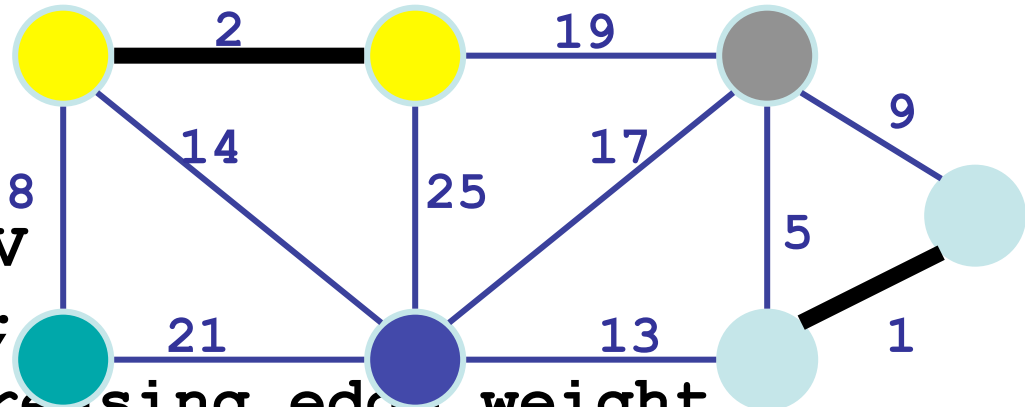
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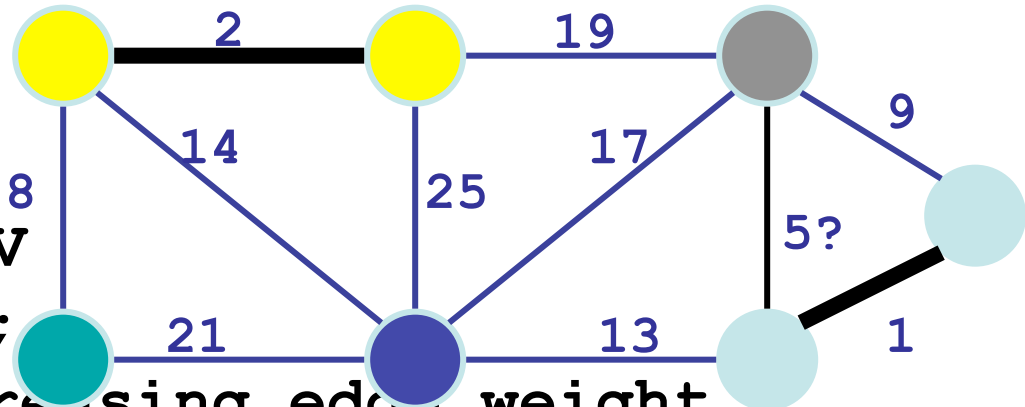
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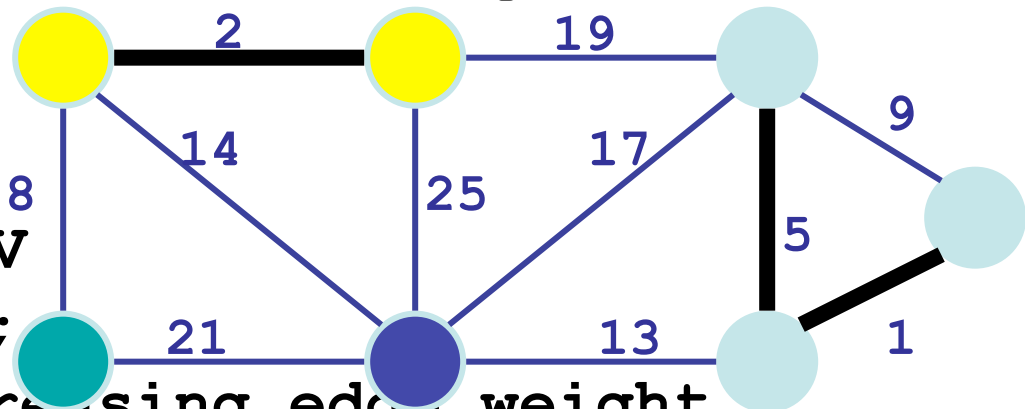
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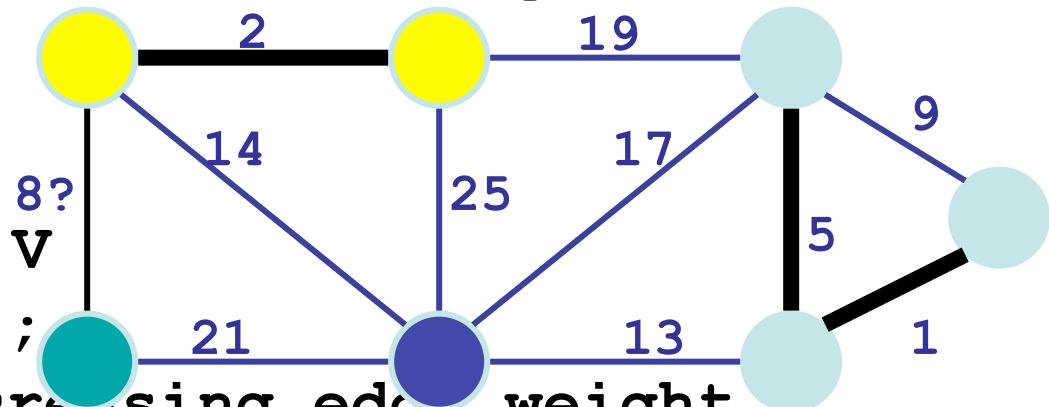
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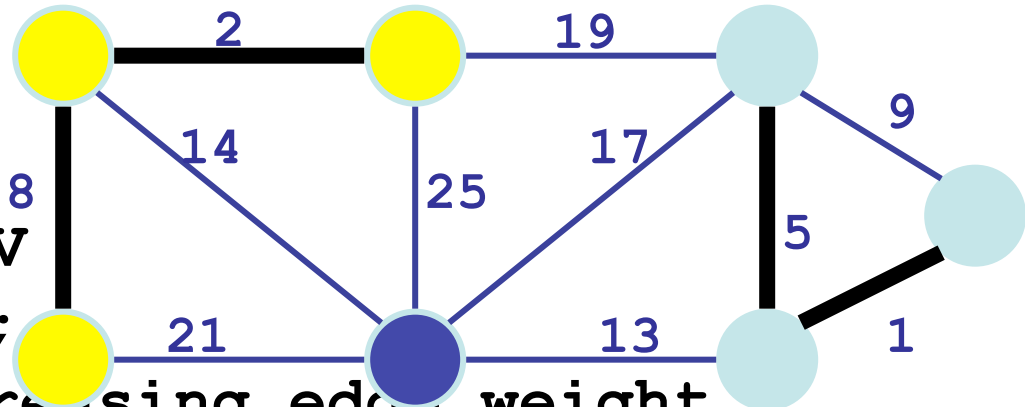
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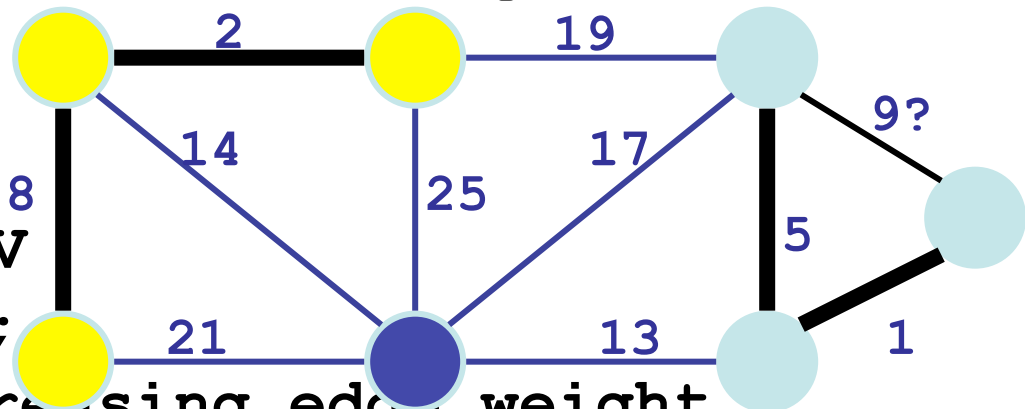
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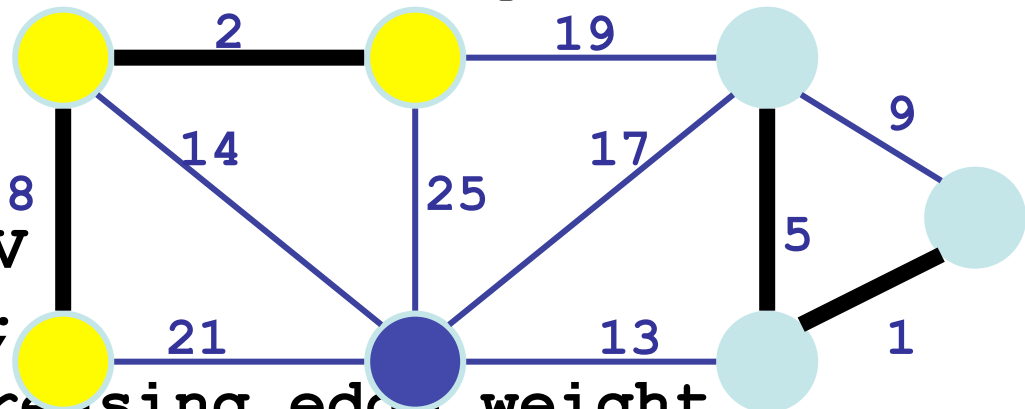
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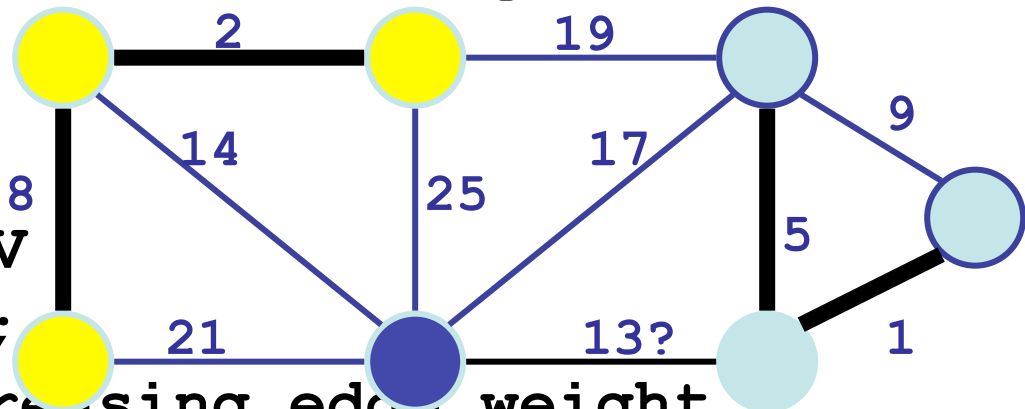
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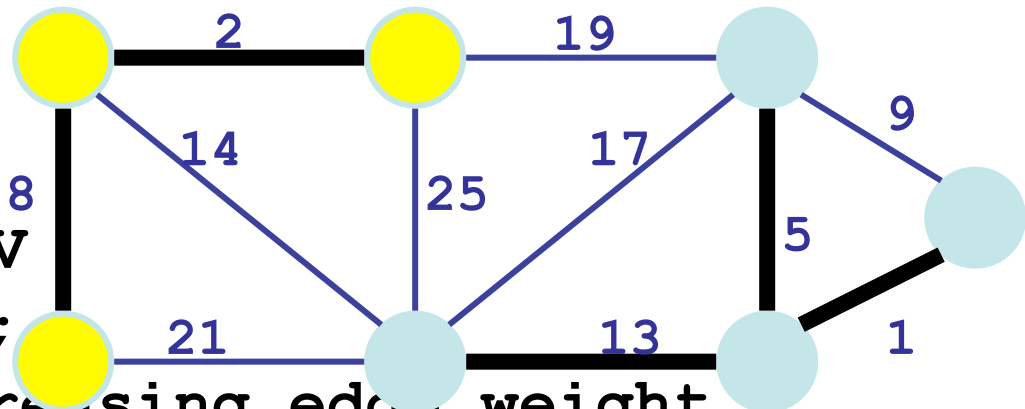
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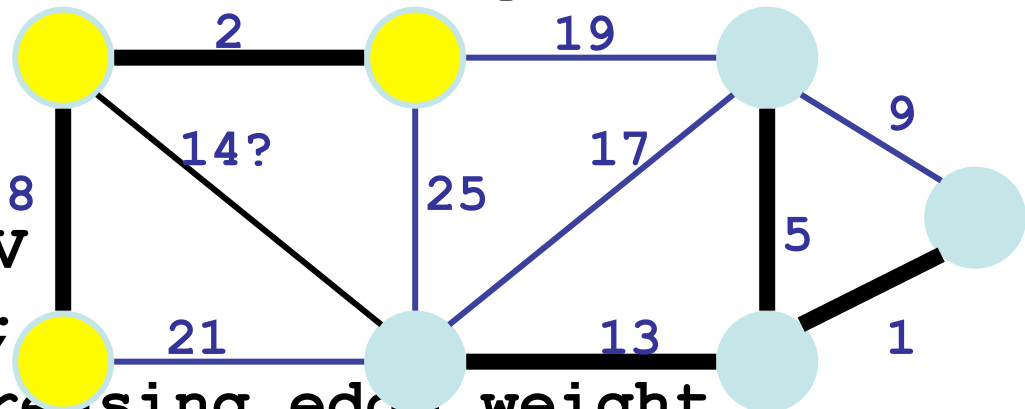
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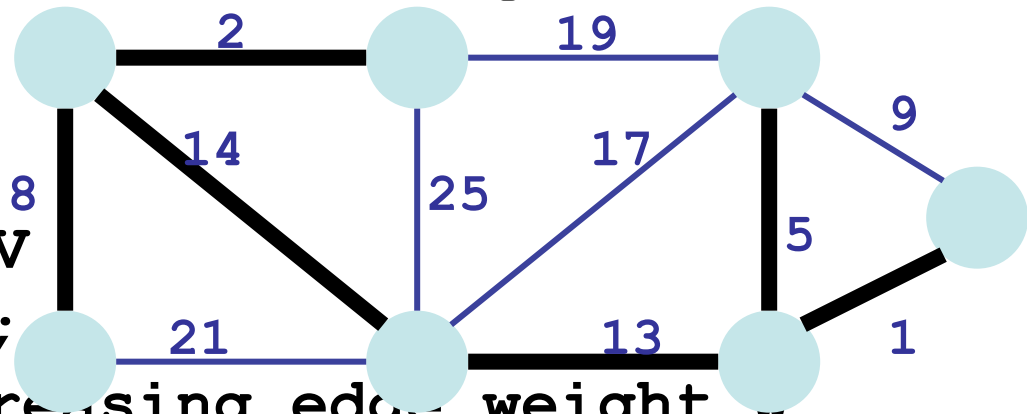
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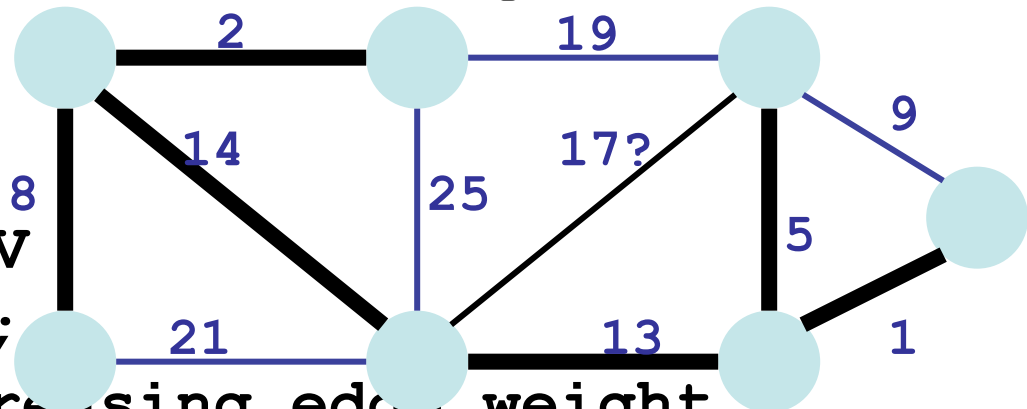
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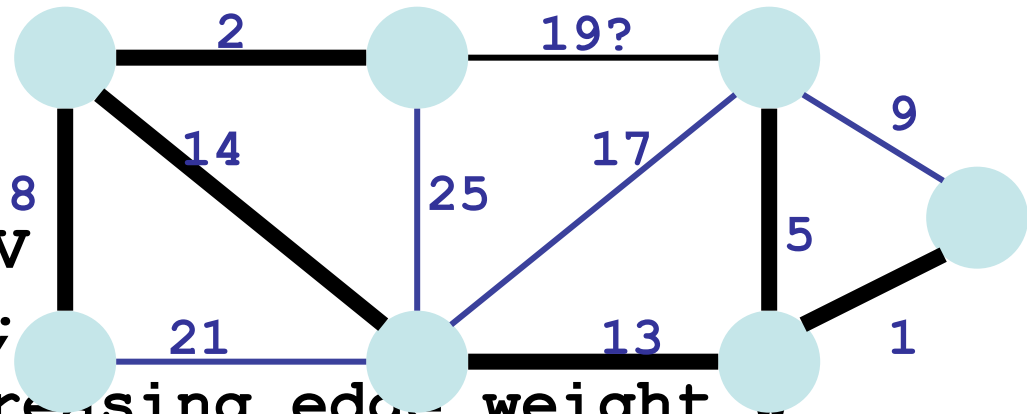
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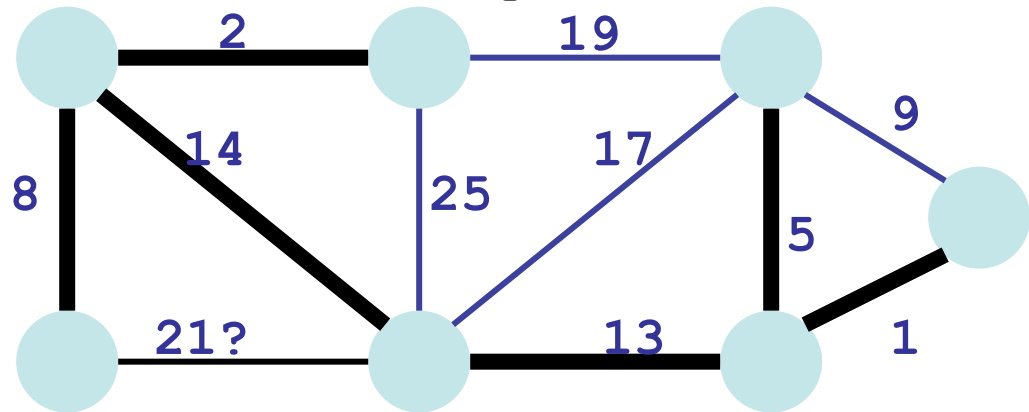
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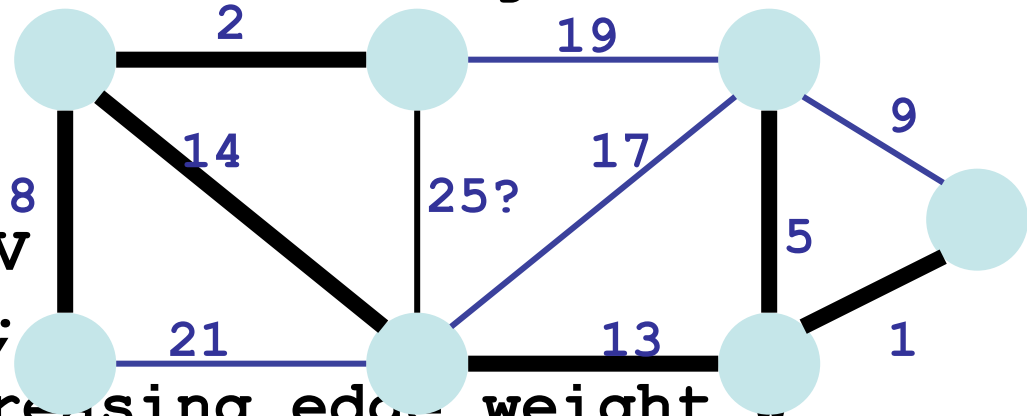
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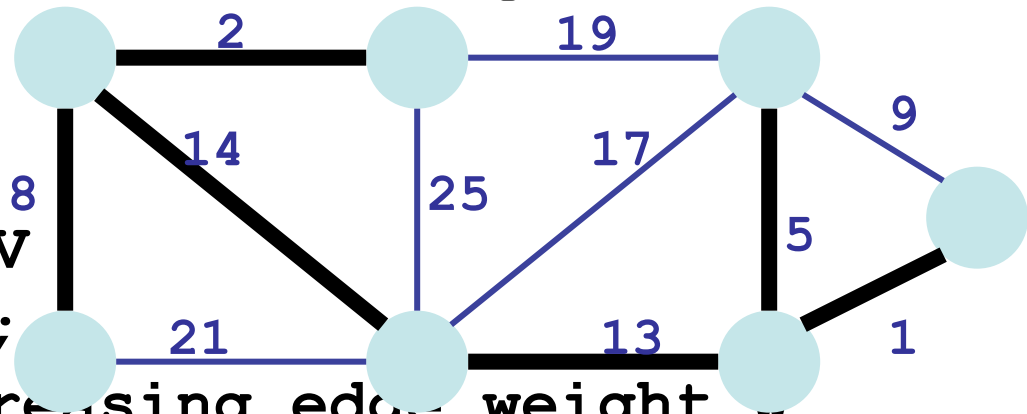
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