Solutions to A Previous Final Examination

December 18, 2006

Print Your Name:

Read the following now.

• Write your name on all pages.
• You have 120 minutes to earn up to 180 points.
• For problems involving calculation, show intermediate steps to ensure partial credit (that is, in case your answers are incorrect).
• Brief and concise answers will be favored in grading.
• Write down your answers clearly. I reserve the right to take off points due to poor writing or English structures.
• One blank page is provided at the end for your convenience.

STOP! Do not turn to the next page until instructed to do so.
1. (35pt in total) Physical and Data Link Layers

(a) (7pt) Give the delay modulation encoding of the bit stream 0, 1, 0, 1, 0, 0, 0, 1. The signal starts with a low value.

```
  1  0  1  0  1  0  0  0  1
```

(b) (8pt) Describe the purpose of medium access control in two sentences.

To resolve the issue which station has the right to access a shared broadcast medium when more than one station need to transmit at the same time.
(c) Consider a communication link with bandwidth $H=5,000$ Hz and $S/N=30$ dB.

i. (10pt) Calculate its maximum data rate according Shannon’s theorem.

Straight $S/N = 1000$.
Max data rate $= 5000 \times \log_2(1 + 1000) \approx 49836$ bps

ii. (10pt) Calculate its maximum data rate according to Nyquest’s theorem with a QAM encoding that uses 4 phase shifts and 4 amplitudes.

QAM is not covered in this semester and thus will not be in your final. The point here is $V = 16$.
Max data rate $= 2 \times 5000 \times \log_2 16 = 40,000$ bps
2. (40pt total) Routing

(a) (20pt) Perform Dijkstra’s algorithm on the network below, using node 0 as the starting point. Give the contents of the `dist` and `nhop` arrays after the completion of 4 iterations.

```
<table>
<thead>
<tr>
<th>u</th>
<th>0</th>
<th>1</th>
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<td>dist</td>
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<td>nhop</td>
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<td>nhop</td>
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<td>nhop</td>
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<td>3</td>
</tr>
</tbody>
</table>
```
(b) (16pt) The distance vector of a router $G$ is given below:

$$(A, 1, 1), (B, 0, 1), (C, 0, 4), (E, 2, 1), (F, 0, 6), (G, -1, 0), (H, 0, 7), (I, 1, 3)$$

where an entry $(Y, p, c)$ indicates $X$ reaches $Y$ through port $p$ in $c$ hops. $X$ receives via port 1 a vector

$$(((A, -1, 0), (B, 0, 2), (C, 0, 1), (D, 3, 1), (E, 2, 1), (F, 3, 3), (G, 1, 1), (H, 4, 1), (I, 4, 6)$$

Answer the following questions.

i. Give the vector of $X$ after processing the incoming vector. Please show entries in the alphabetic order of routers.

First, there is a typo in the question. The router is named $G$, not $X$.

**Answer:**

$(A, 1, 1)$  
$(B, 0, 1)$  
$(C, 1, 2)$  
$(D, 1, 2)$  
$(E, 2, 1)$  
$(F, 1, 4)$  
$(G, -1, 0)$  
$(H, 1, 2)$  
$(I, 1, 7)$

ii. (4pt) Circle the neighboring routers of $G$.

$$A \quad B \quad C \quad D \quad E \quad F \quad G \quad H \quad I$$

**Neighbors are A, B, E**
(c) (40pt in total) Internet Architecture

(a) (8pt) Describe the purpose of the ARP protocol and how it works.

The ARP (Address Resolution Protocol) translates network layer addresses, used by network applications, to hardware addresses required for transmission over LANs. A requesting host broadcasts an ARP request message that includes its own IP and hardware addresses and the IP address in question. The owner of the inquired address responds with its hardware address. The reply is to the requesting host only.

(b) (8pt) Give 4 items in the bootstrap configurations of an Internet host.

IP address, DNS address, gateway router address, and network mask.
(c) (8pt) Explain the purpose of the DHCP protocol and how it works.

The DHCP (Dynamic Host Configuration Protocol) provides hosts with the information they need to operate on the Internet. A booting host broadcasts a DHCP request on its LAN and a DHCP server attached to that LAN responds with its bootstrap configurations as given in the above question.

(d) (6pt) A TCP segment is entering a private network that uses a gateway PC to perform IP masquerading. Circle the field(s) in TCP/IP headers of the packet that has/have to be updated by the gateway PC.

- source IP address
- source port number
- destination IP address — Y
- destination port number — Y
- window size advertisement
- IP header checksum — Y
(e) (5pt) Circle correct statement(s) about Internet data delivery.

• Datagrams may be lost in transit. — Y
• Datagrams may be fragmented in transit. — Y
• Datagrams suffering transmission errors must be discarded by routers.
• Segments of a TCP connection always arrive at the destination in the right order.
• Segments of a TCP connection always use the same path to reach destination.

(f) (5pt) Circle the functions that are performed by the network layer in the Internet model.

• DNS
• fragmentation — Y
• flow control
• congestion control
• routing — Y
• data encryption
(d) (35pt total) Transmission Control Protocol

(a) In this question, you use the RTT estimation that accounts for sample variations, to process three RTT samples: 150, 200 and 210 in that order. Initial RTT=200 and initial DEV=50. Answer the following questions.

i. (3pt) Show the initial Timeout (that is, the value before processing the three samples.)

\[
\text{Timeout} = \text{RTT} + 4 \times \text{DEV} = 200 + 4 \times 50 = 400 \text{ ms.}
\]

ii. (12pt) Give the RTT and Timeout after processing the three samples.

**Processing Sample 1: 150**
\[
\begin{align*}
\text{DIFF} &= 150 - 200 = -50 \\
\text{RTT} &= 200 + \frac{1}{8} \times (-50) = 193.75 \\
\text{DEV} &= 50 + \frac{1}{4} \times (| -50 | - 50) = 50 \\
\text{Timeout} &= 193.75 \times 4 \times 50 = 393.75
\end{align*}
\]

**Processing Sample 2: 200**
\[
\begin{align*}
\text{DIFF} &= 200 - 193.75 = 6.25 \\
\text{RTT} &= 193.75 + \frac{1}{8} \times 6.25 = 194.53125 \\
\text{DEV} &= 50 + \frac{1}{4} \times (| 6.25 | - 50) = 39.0625 \\
\text{Timeout} &= 194.53125 + 4 \times 39.0625 = 350.78125
\end{align*}
\]

**Processing Same 3: 210**
\[
\begin{align*}
\text{DIFF} &= 210 - 194.53125 = 15.46875 \\
\text{RTT} &= 194.53125 + \frac{1}{8} \times 15.46875 = 196.4648375 \\
\text{DEV} &= 39.0625 + \frac{1}{4} \times (|15.46875 | - 39.0625) = 33.1640625 \\
\text{Timeout} &= 196.4648375 + 4 \times 33.1640625 = 329.1210879
\end{align*}
\]

**Comments:** Notice that the last two samples 200 and 210 are close to the RTT and hence reduce the average difference DEV, which in turn results in better confidence of the RTT value and tighter (smaller) timeout intervals.
(b) (15pt) Show the acknowledgments in response to the following segments, arriving in that order:

- Seg(2500,500, with SYN=1) — Ack (3001)
- Seg(4000,200) — Ack (3001)
- Seg(4200,300) — Ack (3001)
- Seg(3000,1000) — Ack (4500)
- Seg(4500,100, with FIN=1) — Ack(4601)

where Seg(X,Y) denotes a segment with sequence number X and Y bytes of data.

The answers above may seem peculiar. That is because byte 3000 was transmitted twice. It was the last byte in the first segment and also the first byte in the 4th segment. Such overlapping does not affect acknowledgments.

(c) (5pt) Circle correct statement(s) of TCP connections.

- initial sequence number is 0 — Y
- initial window size is 1 (segment) — Y
- segments of a connection follows the same path
- segments of a connection arrive at the destination “machine” in the right order
- segments of a connection arrive at the destination “application” in the right order — Y
(e) (30pt) Security
Consider the following remote login procedure:

Step 1. the client asks the user to enter username and password.
Step 2. the client asks the server to provide an IDEA key.
Step 3. the server chooses a random IDEA key $k$ and computes
$k' = \text{RSA}_{\text{encrypt}}(k, \text{private_key(server)})$.
Step 4. the server sends $k'$ to the client.
Step 5. the client computes the digest of the password: $D = \text{MD5(password)}$.
Step 6. the client computes $M' = \text{IDEA}_{\text{encrypt}}(M, k)$, where $M = \text{username}|D$.
Step 7. the client computes $k'' = \text{RSA}_{\text{encrypt}}(k, \text{private_key(client)})$.
Step 8. the client sends $(M', k'')$ to the server.

Answer the following questions.

(a) (8pt) Give the computation/equation the server uses to recover the plaintext request $M$.

$$M = \text{IDEA}_{\text{decrypt}}(M', k)$$

Note: the server needs not to recover $k$. It chose $k$ in the first place.

(b) (8pt) Give the computation/equation the server uses to authenticate the message, that is, to ensure it comes from the client machine.

$$k' = \text{RSA}_{\text{decrypt}}(k'', \text{public_key(client)})$$
(c) (8pt) Give the computation/equation the server uses to check the correctness of the password. We assume that the server maintains a file of usernames and passwords. Use the notation Passwd($x$) to refer to the password of user $x$ stored in that file.

$$\text{MD5(Passwd(username))} == D$$

where previously recovered $M = \text{username}||D$

(d) (6pt) Suppose a third part intercepts the message ($M', k''$). Argue that a replay of the message will not allow the third party to login successfully.

For the second login, the server would use a different $k$ value. The server cannot recover $M$, causing login failures.
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