# Mutation Operators for Ada 

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#### Abstract

Mutation analysis is a method for testing software. It provides a method for assessing the adequacy of test data. This report describes the mutation operators defined for the Ada programming language. The mutation operators are categorized using syntactic criteria, in a form suitable for an implementor of a mutation-based system, or a tester wishing to understand how mutation analysis can be used to test Ada programs.

Each mutation operator is carefully defined, and when appropriate, implementation notes and suggestions are provided. We include operators for all syntactic elements of Ada, including exception handling, generics, and tasking. A summary table listing all operators for Ada, and compared with C and Fortran operators is also provided. The design described here is the result of deliberations among the authors in which all aspects of the Ada language and software development in Ada were considered. These operators can also be viewed as the culmination of previous mutation operator definitions for other languages. This report is intended to serve as a manual for the Ada mutation operators.


## 1 MUTATION APPLIED TO ADA

Although mutation has been applied to several languages, the Ada programming language presents several new features, both syntactic and semantic. Most of the mutation operators for other languages are still meaningful for Ada, with appropriate changes in the names. In this section, we discuss these new features and make recommendations for handling them.

### 1.1 Strong Typing in Ada

Ada's strong typing requires that mutation operators be applied more strictly than in previous languages. Ada's static semantic restrictions, which enhance program reliability and readability, make brute-force ap-
plications of mutation operators likely to generate many invalid programs. For example, determining if the replacement of one arithmetic operator by another will yield a valid Ada expression is complicated by Ada's overloading rules.

In our mutation operator definitions, we point out situations when certain operators should not be applied because of the strong typing rules. We also divide the operators up at a fine-grain level to separate types. Although applying these restrictions require more care when building a mutation system, each program will have fewer mutants.

### 1.2 Syntactically Incorrect Mutants

A stillborn mutant is a mutant that is syntactically illegal. A trivial mutant is a mutant that is killed by almost any test case. An Ada mutation system will be expected to avoid generating stillborn mutants as often as is possible, and a well-designed set of operators will avoid most trivial mutants. While avoiding stillborn mutants will require more care on the part of the developer, and more analysis on the part of the mutation system, it will result in fewer Ada mutants, which will reduce the cost of testing.

In our definitions of mutation operators, we make several decisions to avoid trivial mutants. These decisions are mentioned explicitly.

### 1.3 Overloading in Ada

Ada allows operations to be "overloaded", that is, the same name can be used to represent more than one operation. This is primarily used when the same semantic operation is used for different types. For example, a print operation might be written for numeric types and character types, and the same name can be used to implement the operation for both types.

The only time that this feature impacts our operators is when one variable is being replaced by another. In this case, if the types are different, the replacement is normally not done. However, if the replacement results in an operation is valid because of overloading, then the replacement is done.

### 1.4 Pointers

In Fortran mutation systems such as Mothra [?], there were no pointers. In C, pointers are pervasive and untyped, and a pointer can be derived for any variable by "dereferencing". Additionally, C allows arithmetic operations to be applied to pointers, and even allows pointers to be mixed in expressions with integers and other types. Thus, C mutation systems such as PiSCES [?] must deal with many complexities because of the language.

In Ada, pointers are typed, and a limited number of operators can be applied to pointers. Thus, we are able to, in most cases, treat pointer types as just another kind of operand. A pointer reference is treated just as a variable reference, and we define Operand Replacement Operators to make appropriate substitutions.

### 1.5 Exception Handling in Ada

Ada includes a novel exception handling mechanism. Each program unit can have an associated exception handler, which defines statements to be executed when a particular exception is raised. When an exception is raised, if there is a handler for that particular exception, the associated statements are executed, otherwise the exception is propagated to the calling program unit. Some exceptions are built in to the language, and others can be defined by the programmer. The most common mistake is to use the wrong handler, or have the incorrect exception being raised. This situation is explicitly handled by a mutation operator (the SER operator).

### 1.6 Generic Packages in Ada

Ada allows packages (and procedures) to be generic, so that the package can be instantiated with parameters such as types etc. Our opinion is that mutations to generic package headers or instantiations would result in trivial mutants, thus we do explicitly mutate for generics.

It is possible that the testing process should be modified to reflect generics, so that a generic package is tested with various instantiations, but this is beyond the scope of mutation testing.

### 1.7 Tasking in Ada

Tasking is a major feature of the Ada language that is new to mutation systems. As a result, there is little experience with handling such constructs. We present a discussion of the issues involved with handling tasking, and include some initial recommendations for mutation operators in the mutation operator section.

The most obvious and relevant aspect of programs that use tasking is that the execution can be nondeterministic. Thus, executing a non-deterministic program on a test case will generate one among a potentially large set of correct outputs. We call this set of correct outputs of a test case on a program the feasible output set, and each correct execution will produce one element of this set. A mutant is equivalent if it generates the same feasible output set as the original program for all inputs. Unfortunately, this means that equivalent non-deterministic programs can produce different output. A mutant is weakly equivalent if its feasible output set is a subset of the feasible output set of the original program. That is, if a mutant always produces some correct output, but is incapable of producing all correct outputs, it is considered to be weakly equivalent.

The non-deterministic nature of tasking programs requires a modified definition of correctness. The output of a mutant program on a test case is correct if the output is in the feasible output set of the program. With deterministic programs, mutation systems can easily determine the correctness of a given output; they merely compare the output with the output of the original program on the same test case. With non-deterministic programs, however, that is not possible. We suggest the following approximation.

Run the original program $N$ times on the same test case $t$ to create $N$ outputs $O_{i}, 1 \leq i \leq N$. The set $\Omega=\left\{O_{1}, O_{2}, \ldots, O_{N}\right\}$ is an approximation of the feasible output set. Run the mutant program $m$ on the test case to create $O(m, t)$. If $O(m, t) \in \Omega$, then the mutant is left alive, else it is killed. Note that if $P$ is a deterministic program, mutant output checking is a special case of the non-deterministic case, where $\Omega$ has a single element.

Of course, how well $\Omega$ approximates the true feasible output set of the program depends largely on the value of $N$. To get the true feasible output set, we may need to run the program an infinite number of times. A value for $N$ can be estimated at several points during testing:

1. Determined by the mutation system as a constant for all program,
2. Set by the tester for each application,
3. Estimated by the system for each test case. Run the original repeatedly until a small number of executions are made without creating a new output.

The third option is more precise, and should result in a more accurate estimation of the feasible output set, but will be more expensive than the other two options. Moreover, the third option will not always terminate if each execution produces a different output. For example, if the feasible output set is infinite and each element is equally likely to be generated, we can expect each execution to produce a new output. In this case, the approximation approach will not work anyway, because each mutant can be expected to generate a unique output.

The options above are all automatable, and based on enumerating all or part of the feasible output set. Another method of determining whether a mutant's output is correct is based on a semi-automated method. If the tester can describe the feasible output set in some way, then the output can be checked to see if it matches the description. Unfortunately, this description depends on the application program. It is possible that a general-purpose language could be devised that would allow a tester to enter a description of the feasible output set.

## 2 ADA MUTATION OPERATORS

For a program $P$, mutation testing produces a set of alternate programs. Each alternate program, $P_{i}$, known as a mutant of $P$, is formed by modifying a single statement of $P$ according to some predefined modification rule. These modification rules are called mutation operators. The syntactic change itself is called the mutation, and the resulting program is the mutant program, or simply mutant. The original program plus the mutant programs are collectively known as the program neighborhood, $N$, of $P$ [?].

This report defines a set of mutation operators for the Ada programming language. Our operators are partially based on the previous operators defined for Ada [?], the C operators [?], and the Fortran-77 [?] operators used by Mothra [?]. These operators are complete for the Ada language as defined in the Ada Reference Manual [?]. This is as distinct from the previous Ada operators, which did not cover the entire language. Additionally, the previous operators were designed with very little experience with the Ada language, or writing Ada programs. As a result, our operators are significantly more extensive than the earlier set. Because of the extensive experience we have had with the Fortran mutation operators as implemented in the Mothra testing system, they have influenced our Ada operators most heavily.

We organize our operators differently than authors of previous sets of operators. In particular, we separate our operators primarily on the basis of what type of lexical elements are modified; this gives us four types of operators. Mutation operators within these groups have reasonably uniform semantics and rules for applications. Also, the number of mutants produced are on the same order of magnitude for all operators within our types. We also include one type of operators (Coverage), specifically to include branch coverage testing strategies.

The five types of mutation operators for Ada are:

- Operand Replacement Operators (30 operators)
- Statement Operators (14 operators)
- Expression Operators (14 operators)
- Coverage Operators (4 operators)
- Tasking Operators (3 operators)

We have a total of 65 operators. In the following section, we define each operator in turn. Following that, we present all of our Ada operators in one comprehensive table, shown with the correlating Fortran and C operators, if any. Lastly, we show an Ada package with example mutants.

## 3 ADA MUTATION OPERATOR DEFINITIONS

This section comprises the major part of this report, both in technical terms, and in bulk. Each type of mutation operator is discussed in a separate subsection, and each individual operator is defined. Each subsection starts with a general discussion about the operator type, then a table is given listing all the operators of that type. Next, each operator is defined in turn.

### 3.1 Operand Replacement Operators

Each operand replacement operator starts with the letter 0 . There are 29 operand replacement operators. These operators cause each operand to be replaced by each other syntactically legal operand. There are five kinds of operands in Ada:

1. Variables
2. Constants
3. Array References
4. Record References
5. Pointer References

Although in Ada there is no real difference between record and pointer references, we define separate operators so as to have uniform definitions. Replacing these five kinds of operands result in 25 operators; there are four additional operators for three structured types, and one additional operator for variable initialization.

| Operand Replacement Operators |  |
| :--- | :--- |
| OVV | Variable replaced by a variable. |
| OVC | Variable replaced by a constant. |
| OVA | Variable replaced by an array reference. |
| OVR | Variable replaced by a record reference. |
| OVP | Variable replaced by a pointer reference. |
| OVI | Variable initialization elimination. |
| OCV | Constant replaced by a variable. |
| OCC | Constant replaced by a constant. |
| OCA | Constant replaced by an array reference. |
| OCR | Constant replaced by a record reference. |
| OCP | Constant replaced by a pointer reference. |
| OAV | Array reference replaced by a variable. |
| OAC | Array reference replaced by a constant. |
| OAA | Array reference replaced by an array reference. |
| OAR | Array reference replaced by a record reference. |
| OAP | Array reference replaced by a pointer reference. |
| OAN | Array name replaced by an array name. |
| ORV | Record reference replaced by a variable. |
| ORC | Record reference replaced by a constant. |
| ORA | Record reference replaced by an array reference. |
| ORR | Record reference replaced by a record reference. |
| ORP | Record reference replaced by a pointer reference. |
| ORF | Record field replaced by a record field. |
| ORN | Record name replaced by a record name. |
| OPV | Pointer reference replaced by a variable. |
| OPC | Pointer reference replaced by a constant. |
| OPA | Pointer reference replaced by an array reference. |
| OPR | Pointer reference replaced by a record reference. |
| OPP | Pointer reference replaced by a pointer reference. |
| OPN | Pointer name replaced by a pointer name. |

## Notes:

The strong typing rules of Ada will drastically reduce the number of mutants of this type that are generated.

Do mutate initializations (only OCC).

Do mutate references of enumerated types.

Do not mutate types.
Do not mutate declarations.

Do not mutate CASE constants.

Do not mutate loop parameters on FOR statements (it is a declaration).
Variables that are of a type that is declared externally and private are considered to be scalar.

The following named objects are considered as CONSTANT and are mutated using OVC, OC?, ORC, OAC, and OPC:

- Objects declared with the keyword CONSTANT
- Loop parameters
- Parameters of class IN

1. The 25 simple replacement operators are all uniform and merely replace one type with another.
2. OVI: Variable initialization elimination.

Eliminate the initialization part of each variable initialization.
3. OAN: Array name replaced by an array name.

Replace just the array name in an array reference by other array names when the base types are the same, and the index types are the same.
4. ORF: Record field replaced by a record field.

Replace a record field reference by another field name of the same record when the second field is of the same type.
5. ORN: Record name replaced by a record name.

Replace just the record name in a record reference by other record names when the field names and types are the same.
6. OPN: Pointer name replaced by a pointer name.

Replace just the pointer name in a pointer reference by other pointer names when the field names and types are the same.

### 3.2 Statement Modification Operators

Each statement modification operator starts with the letter S. There are 13 statement modification operators. These operators modify entire statements and modify the control structures of Ada. The relevant control structures are:

## 1. BLOCK

2. CASE
3. EXIT
4. FOR
5. GOTO
6. IF
7. LOOP
8. RAISE
9. RETURN
10. WHILE

We summarize the operators in a table, then discuss each operator in detail.

| Statement Modification Operators |  |
| :--- | :--- |
| SEE | Exception on execution. |
| SRN | Replace with NULL. |
| SRR | Return statement replacement. |
| SGL | GOTO label replacement. |
| SRE | Replace with EXIT. |
| SWR | Replace WHILE with repeat-until. |
| SRW | Replace repeat-until with WHILE. |
| SZI | Zero iteration loop. |
| SOI | One iteration loop. |
| SNI | N iteration loop. |
| SRI | Reverse iteration loop. |
| SES | END shift. |
| SCA | CASE alternative replacement. |
| SER | RAISE exception handler replacement. |

1. SEE: Exception on execution.

Replace the first statement in each basic block with:
RAISE mut_trap;
mut_trap is a mutation-defined exception. We need to get the mutant number to the handler, which can be in the local procedure or the main program. We recommend having the mutant call a subroutine:

Except_on_Exec (42);
which then raises the exception.

## Notes:

Do not replace if elimination of the statement would result in a compile-time error, for example, if the statement is the only RETURN in a function.

SEE should be applied to statements and block statements. For example, the following structure will result in four mutants:

```
WHILE (e1) LOOP
    IF (e2) THEN
        s1;
    ELSE
        s2;
    END IF;
END LOOP;
Mutant 1: Except_on_Exec (n);
Mutant 2: WHILE (e1) LOOP
    Except_on_Exec (n);
        END LOOP;
```

```
Mutant 3: WHILE (e1) LOOP
    IF (e2) THEN
                        Except_on_Exec (n);
        ELSE
            s2;
        END IF;
END LOOP;
Mutant 4: WHILE (e1) LOOP
        IF (e2) THEN
            s1;
        ELSE
            Except_on_Exec (n);
        END IF;
END LOOP;
```

2. SRN: Replace with NULL.

Replace each statement with NULL.
The replacement should be done according to the rules of the SEE mutation operator, except the replacement should be done on each statement, not each basic block.

Do not replace if elimination of the statement would result in a compile-time error, for example, if the statement is the only RETURN in a function.
3. SRR: Return statement replacement.

Replace each statement in a FUNCTION or PROCEDURE with RETURN.
For parameterized RETURN statements (in functions), replace each statement with every RETURN that appears in the function. Do not replace RETURN statements.

Do not replace if elimination of the statement would result in a compile-time error.
4. SGL: GOTO label replacement.

Replace each GOTO label with all other visible, legal labels.
The innermost sequence of statements that encloses the target statement must also enclose the GOTO statement (note that the GOTO statement can be a statement of an inner sequence). Furthermore, if a GOTO statement is enclosed by an ACCEPT statement or the body of a program unit, then the target statement must not be outside this enclosing construct; conversely, it follows form the previous rule that if the target statement is enclosed by such a construct, then the GOTO statement cannot be outside.
5. SRE: Replace with EXIT.

This operator replaces statements within loops with EXIT statements. There are three variations.
(a) Replace each statement in a loop with EXIT;
(b) Replace each statement in a loop with an EXIT name; for each named enclosing loop.
(c) Replace each statement in a loop with each EXIT WHEN ...; that appears in the loop.

## Notes:

If there is only one statement in the loop, this change would be equivalent to SRN, so do not generate. The C operator SBR only did the second of the three variations.
6. SWR: Replace WHILE with repeat-until.

Although there is no explicit repeat-until statement in Ada, the construct is commonly built using a LOOP and an EXIT. Using the incorrect kind of loop is a common programming mistake. The format of the change is:

7. SRW: Replace repeat-until with WHILE.

This is the opposite of SWR. The format of the change will be:

```
    ORIGINAL MUTANT
LOOP WHILE (NOT e) LOOP
    . .
EXIT WHEN e;
END LOOP;
END LOOP;
```

Rather than only applying this operation to loops where the EXIT WHEN statement is the last statement in the loop body, it is applied to all EXITs in the loop.
8. Definite loop mutations.

We have four goals for mutating definite loops (FOR).
(a) Bypass the loop entirely (zero iterations).
(b) Cause the loop to iterate once (one iteration).
(c) Cause the loop to be iterated more than once ( N iterations).
(d) Cause the loop to be executed in reverse (reverse iteration).

The first three goals are satisfied by introducing a new loop counter for each loop. For a loop $i$, associate the counter loop_i_count. loop_i_count is initialized to zero before the loop begins. It is incremented by one each iteration through the loop.
(a) SZI: Zero iterations

After the loop, if loop_i_count $=0$, then RAISE Mut_Trap;
(b) SOI: One iteration

After the loop, if loop_i_count $=1$, then RAISE Mut_Trap;
(c) SNI: N iterations

After the loop, if loop_i_count $>1$, then RAISE Mut_Trap;
(d) SRI: reverse iteration

Add the keyword REVERSE to the loop if it is not there, remove it if it is there.
9. SES: END shift.

Move each END statement up and down one statement. This applies to END statements occurring in BLOCK and LOOP statements, but not CASE statements and subprograms.
10. SCA: CASE alternative replacement.

First, each case statement alternative with multiple choices is separated into alternatives where each alternative contains only one choice. A range (e.g., 5..20) is considered to be only one choice. Next, substitute each statement sequence with each other sequence in the CASE statement.

Do not mutate choices.
Example:

```
CASE Var1 is
    WHEN A | B => statements_1;
    WHEN C => statements_2;
    WHEN OTHERS => statements_3;
END CASE;
```

This case statement creates 8 mutants:

```
CASE Var1 is
    WHEN A => statements_2;
    WHEN B => statements_1;
    WHEN C => statements_2;
    WHEN OTHERS => statements_3;
END CASE;
CASE Var1 is
    WHEN A => statements_3;
    WHEN B => statements_1;
    WHEN C => statements_2;
    WHEN OTHERS => statements_3;
END CASE;
CASE Var1 is
    WHEN A => statements_1;
    WHEN B => statements_2;
    WHEN C => statements_2;
    WHEN OTHERS => statements_3;
END CASE;
CASE Var1 is
    WHEN A => statements_1;
    WHEN B => statements_3;
    WHEN C => statements_2;
    WHEN OTHERS => statements_3;
END CASE;
CASE Var1 is
    WHEN A => statements_1;
    WHEN B => statements_1;
    WHEN C => statements_1;
    WHEN OTHERS => statements_3;
END CASE;
CASE Var1 is
    WHEN A => statements_1;
    WHEN B => statements_1;
    WHEN C => statements_3;
    WHEN OTHERS => statements_3;
END CASE;
CASE Var1 is
    WHEN A => statements_1;
    WHEN B => statements_1;
    WHEN C => statements_2;
    WHEN OTHERS => statements_1;
END CASE;
CASE Var1 is
    WHEN A => statements_1;
```

```
    WHEN B => statements_1;
    WHEN C => statements_2;
    WHEN OTHERS => statements_2;
END CASE;
```

11. SER: RAISE exception handler replacement.

For each explicit RAISE statement, replace the name of the exception by other exceptions. Replace programmer-defined exceptions only by other programmer-defined exceptions, and built-in exceptions by other built-in exceptions.

### 3.3 Expression Modification Operators

Each expression modification operator starts with the letter E. There are 14 expression modification operators. These operators modify expression operators and entire expressions. We summarize the operators in a table, then discuss each operator in detail.

| Expression Modification Operators |  |
| :--- | :--- |
| EAI | Absolute value insertion. |
| ENI | Neg-Absolute value insertion. |
| EEZ | Exception on zero. |
| EOR | Arithmetic operator replacement. |
| ERR | Relational operator replacement. |
| EMR | Membership test replacement. |
| ELR | Logical operator replacement. |
| EUI | Unary operator insertion. |
| EUR | Unary operator replacement. |
| ESR | Subprogram operator replacement. |
| EDT | Domain twiddle. |
| EAR | Attribute replacement. |
| EEO | Exception on overflow. |
| EEU | Exception on underflow. |

1. EAI: Absolute value insertion.

Insert the unary operator ABS in front of every arithmetic expression and subexpression. Do not mutate if the expression can be statically determined to be greater than or equal to zero. For example, we can determine this for the following cases:

- Constants
- The type is a nonnegative subtype of Integer (for example, Natural or Positive)
- Loop variable where the lower bound is greater than or equal to zero.

Do not mutate if the change will make a discrete range in a FOR statement have a NULL range (this would be equivalent to an SZI mutant).
Do not mutate CONSTANTS (this would be equivalent to a EUI mutant).
2. ENI: Neg-absolute value insertion.

Insert - ABS in front of every arithmetic expression and subexpression. Do not mutate if the expression can be statically determined to be less than or equal to zero. For example, we can determine this for the following cases:

- Constants
- The type is a negative subtype of Integer.

Do not mutate if the change will make a discrete range in a FOR statement have a NULL range (this would be equivalent to an SZI mutant).
Do not mutate CONSTANTS (this would be equivalent to a EUI mutant).
3. EEZ: Exception on zero.

Insert the subprogram Except_on_Zero in front of every arithmetic expression and subexpression. Except_on_Zero(E); raises EEZ_Exception if E is 0, else it returns E. Do not mutate if the expression can be statically determined to be not equal to 0 . For example, we can determine this for the following cases:

- Constants.
- The type is a subtype of Integer that does not include 0 (for example, Positive).
- Loop variable where the range does not include 0 .

4. EOR: Arithmetic operator replacement.

Replace each binary arithmetic operator $(+,-, *, /, \mathrm{MOD}, \mathrm{REM}, * *)$ with each other binary arithmetic operator that is syntactically legal.
Strong typing notes:

- MOD and REM are only defined for Integer types.
- ** requires the right operand to be Integer.
-     * allows Fixed Point and Integer to be mixed.
- / allows Fixed Point on the left and Integer on the right.

5. ERR: Relational operator replacement.

Replace each relational operator with each other relational operator that is syntactically legal.
$<,>,>=,<=$ are only defined for scalar and discrete array types.
6. EMR: Membership test replacement.

Replace each IN with NOT IN and each NOT IN with IN.
Note: This operator is subsumed by the CDE operator and should not be used if CDE is.
7. ELR: Logical operator replacement.

Replace each logical operator (AND, OR, XOR, AND THEN, OR ELSE) with each other logical operator.
AND, OR, and XOR are defined for Boolean expressions and one-dimensional arrays of type Boolean.
8. EUI: Unary operator insertion.

Insert the unary operator - in front of each arithmetic expression and subexpression.
Note: the unary operator + is the identity operation.
9. EUR: Unary operator replacement.

Replace each unary operator $(+,-, A B S)$ with each other unary operator.
Expressions should be fully parenthesized, since ABS has higher precedence than + and - .
10. ESR: Subprogram operator replacement.

Replace each function and subroutine name with each other function or subroutine name that has the same syntactic signature and comes from the same package.
Also replace with $=$ and $/=$ if the signature is appropriate $(=$ and $/=$ are implicitly defined for all types).
Do not consider the parameter class in the signature comparison.

## Example:

$\underline{\text { Package Matrix Specification: }}$

```
Matrix_Type ...
"+" (M1, M2: Matrix_Type) RETURN Matrix_Type;
"*" (M1, M2: Matrix_Type) RETURN Matrix_Type;
"<" (M1, M2: Matrix_Type) RETURN Boolean;
```

Matrix Use:

```
A, B, C : Matrix_Type;
    :
C := A + B; ==> mutation ==> C := A * B;
IF (A < B) ... ==> mutation ==> IF (A = B) ...
    =>> mutation ==> IF (A /= B) ...
```

11. EDT: Domain twiddle.

Each innermost expression (operand: constant, variable, array reference, record reference, pointer reference) is twiddled, that is, modified by a small amount. For each operand, the modification produces two mutants, one where the modification is in a positive direction, the other in a negative direction. This amount depends on the type:

| TYPE |  | MODIFICATION |
| :--- | :--- | :--- |
| Integer | $-\quad+1$ and -1 |  |
| Float | $-\quad{ }^{*} 1.05$ and ${ }^{*} .95$ |  |
| Fixed Point | $-\quad+$ Delta and -Delta |  |
| Character Types | - T'SUCC and T'PRED |  |
| Enumerated Types | $-\quad$ T'SUCC and T'PRED |  |

The twiddle must not create a value out of the range for that type. For example, if X is of type Natural and has the value 0 , the mutant -1 is not generated.

## Notes:

Do not twiddle array subscripts - most changes would cause an out-of-bounds failure.
Do not mutate if the change would result in mutant that is equivalent to another twiddle on the same expression (for example, $(\mathrm{X}+\mathrm{Y})==>((\mathrm{X}-1)+\mathrm{Y})$ and $((\mathrm{X}+1)+\mathrm{Y})$, but not $(\mathrm{X}+(\mathrm{Y}-1))$ and $(\mathrm{X}+(\mathrm{Y}+1))$ ).

Do not mutate loop parameters if the change would result in a NULL range (this would be equivalent to an OCC mutant).
12. EAR: Attribute replacement.

Each attribute is replaced by each other syntactically legal attribute. Attributes are defined in Appendix A of the Ada Reference Manual.
13. EEO: Exception on overflow.

Insert the subprogram Except_on_OverFlow in front of every arithmetic expression. Except_on_OverFlow(E) raises EEO_Exception if the expression results in an overflow, else it returns the value of the expression. Do not mutate if the expression can be statically determined to not overflow. For example, we can determine this for the following cases:

- Constants.
- Loop variable.


## Notes:

It might be possible to allow the Ada runtime system to detect overflow problems, and define a handler for the overflow. The ADA reference manual, section 4.5.7, paragraph 7, says:
"If the result overflows, NUMERIC_ERROR should be raised, but will not necessarily be raised. That is, it is not strictly required."

The ADA reference manual, section 13.7.3, says:
"If an overflow occurs, and there is no NUMERIC_ERROR, T'MACHINE-OVERFLOWS is FALSE, else TRUE."

This does not make sense to me, for two reasons:
(a) The value for T'MACHINE_OVERFLOWS seems to be backwards.
(b) NUMERIC_ERROR is not required because detecting overflow is hard in some situations. But setting this attribute requires overflow to be detected.
14. EEU: Exception on underflow.

Insert the subprogram Except_on_UnderFlow in front of every arithmetic expression. Except_on_UnderFlow(E) raises EEU_Exception if the expression results in an underflow, else it returns the value of the expression. Do not mutate if the expression can be statically determined to not underflow. For example, we can determine this for the following cases:

- Constants.
- Loop variable.


### 3.4 Coverage Operators

The previous operators do not cover the branch coverage criteria [?] as do the Mothra operators [?]. For the Ada operators, we have chosen to define separate operators expressly for this purpose. This is so that the tester can explicitly choose to cover one or more of the branch coverage criteria, without having to use other operators.

The coverage criteria we consider are based on the following definitions:

Definition: A Condition in a program is a pair of algebraic expressions related by one of the relational operators $\{>,<,=, \geq, \leq, \neq\}$.

Conditions evaluate to one of the binary values TRUE or FALSE and can be modified by the negation operator NOT.

Definition: A Decision is a list of one or more conditions connected by the two logical operators AND and OR and used in a statement that affects the flow of control of the program. Decisions represent branches in the control flow of the program.

Statement Coverage (SC) requires that every statement in the program be executed at least once. Decision Coverage ( $D C$ ) requires that every decision evaluate to both TRUE and FALSE at least once. DC is also known as branch testing and all-edges [?]. Condition Coverage ( $C C$ ) requires that each condition in each decision evaluate to both TRUE and FALSE at least once. Decision / Condition Coverage ( $D C C$ ) requires that each condition in each decision evaluate to both TRUE and FALSE at least once, and that every decision evaluate to both TRUE and FALSE at least once. DCC combines DC and CC. Modified Condition / Decision Coverage ( $M C / D C$ ) requires that every decision and every condition within the decision has taken every outcome at least once, and every condition has been shown to independently affect its decision. Multiple-Condition Coverage (MCC) requires that all possible combinations of condition outcomes in each decision be covered, that is, the entire truth table for the decision has been satisfied. MCC is also known as extended branch coverage [?].

We have designed four Ada operators specifically to cover these coverage criteria. Each coverage operator starts with the letter C. We summarize the four coverage operators in a table, then discuss each operator in detail. The SEE operator satisfies statement coverage, so it is not included in the coverage operators set.

| Coverage Operators |  |
| :--- | :--- |
| CDE | Decision coverage. |
| CCO | Condition coverage. |
| CDC | Decision/condition coverage. |
| CMC | Multiple condition coverage. |

1. CDE: Decision coverage

Each decision must evaluate to both TRUE and FALSE. Replace each decision by TRUE and FALSE.
2. CCO: Condition coverage

Each condition must evaluate to both TRUE and FALSE. Replace each condition by TRUE and FALSE.

Note: there will be some redundancy. This could be reduced by having the mutation system suppress some mutants.
3. CDC: Decision/condition coverage

Decision/condition coverage combines decision and condition. The CDC operator simply turns on CCO and CDE. We define it separately as a convenience.
4. CMC: Multiple condition coverage

All combinations of conditions must be exercised separately, which yields, for a decision with $n$ conditions, $2^{n}$ combinations. Another way of stating the MCC requirement is that the entire truth table for the decision must be covered.
We define the CMD operator in two separate ways. The second way is more efficient.

## 1. Mothra-like definition

(a) Replace each condition with TRUE and FALSE (CCO), and
(b) Replace each subset of conditions with TRUE and FALSE, and
(c) Replace each subset of logical connectors with $\neq(\mathrm{XOR})$ and $=($ NOT XOR $)$.

This will satisfy MCC, but will result in a fair amount of redundancy. There will be far more than $2^{n}$ mutants.
2. Pisces-like implementation At each decision, create $2^{N}$ mutants. Assume the decision is $D$, with $N$ conditions $C_{i}, 1 \leq i \leq N$.

Create the entire truth table for $D, T T$, where $T T_{j}$ is the truth assignment needed for mutant $j$, $1 \leq j \leq 2^{N} . T A$ is the current truth assignment for $D$.
The implementation of the mutant is as follows:

```
IF (TA=TTi}) then
    IF applying weak mutation
        Kill mutant j
    ELSE
            RETURN NOT (TA)
        END IF
END IF
```

This could be done in an evaluative way, as Pisces does, or using Schema.
Note: If CMC is used, CDE, CCO, and CDC are redundant and should not be used.

### 3.5 Tasking Operators

We have designed three Ada operators specifically to cover tasking. Each coverage operator starts with the letter T. We summarize the tasking operators in a table, then discuss each operator in detail.

| Tasking Operators |  |
| :--- | :--- |
| TEM | ENTRY statement modification. |
| TAR | ACCEPT statement replacement. |
| TSA | SELECT alternative replacement. |

1. TEM: ENTRY statement modification

Each ENTRY call is modified just as procedure calls are modified by the ESR operator. Replace each ENTRY call name with each other ENTRY name that has the same syntactic signature and comes from the same task.
Also replace conditional and timed entry calls by simple entries.
2. TAR: ACCEPT statement replacement

Replace entry names by other visible entries of the same time.
3. TSA: SELECT alternative replacement

Each SELECT alternative is modified just as the CASE statement is modified by the SCA operator. First, each SELECT statement alternative with multiple choices is separated into alternatives where each alternative contains only one choice. Next, substitute each statement sequence with each other sequence in the SELECT statement.

## 4 COMPARISON OF ADA, C AND FORTRAN-77 OPERATORS

This section contains a table that attempts to relate our Ada mutation operators with the previous operators for Ada [?], and the C [?] and Fortran-77 [?] operators. The character $\sim$ means that there is no corresponding operator.

| Operand Replacement Operators |  |  |  |
| :--- | :--- | :--- | :--- |
| Ada | Description | C | Fortran-77 |
| OVV | Variable replaced by a variable. | Vsrr | Svr |
| OVC | Variable replaced by a constant. | Vsrr | csr |
| OVA | Variable replaced by an array reference. | Vsrr | asr |
| OVR | Variable replaced by a record reference. | Vsrr | $\sim$ |
| OVP | Variable replaced by a pointer reference. | Vsrr | $\sim$ |
| OVI | Variable initialization elimination. | $\sim$ | $\sim$ |
| OCV | Constant replaced by a variable. | Ccsr | scr |
| OCC | Constant replaced by a constant. | Cccr | src |
| OCA | Constant replaced by an array reference. | $\sim$ | acr |
| OCR | Constant replaced by a record reference. | $\sim$ | $\sim$ |
| OCP | Constant replaced by a pointer reference. | $\sim$ | $\sim$ |
| OAV | Array reference replaced by a variable. | Varr | sar |
| OAC | Array reference replaced by a constant. | Varr | car |
| OAA | Array reference replaced by an array reference. | Varr | aar |
| OAR | Array reference replaced by a record reference. | Varr | $\sim$ |
| OAP | Array reference replaced by a pointer reference. | Varr | $\sim$ |
| OAN | Array name replaced by an array name. | Varr | cnr |
| ORV | Record reference replaced by a variable. | Vtrr | $\sim$ |
| ORC | Record reference replaced by a constant. | Vtrr | $\sim$ |
| ORA | Record reference replaced by an array reference. | Vtrr | $\sim$ |
| ORR | Record reference replaced by a record reference. | Vtrr | $\sim$ |
| ORP | Record reference replaced by a pointer reference. | Vtrr | $\sim$ |
| ORF | Record field replaced by a record field. | VSCR | $\sim$ |
| ORN | Record name replaced by a record name. | $\sim$ | $\sim$ |
| OPV | Pointer reference replaced by a variable. | Vprr | $\sim$ |
| OPC | Pointer reference replaced by a constant. | Vprr | $\sim$ |
| OPA | Pointer reference replaced by an array reference. | Vprr | $\sim$ |
| OPR | Pointer reference replaced by a record reference. | Vprr | $\sim$ |
| OPP | Pointer reference replaced by a pointer reference. | Vprr | $\sim$ |
| OPN | Pointer name replaced by a pointer name. | $\sim$ | $\sim$ |
|  |  |  |  |


| Statement Modification Operators |  |  |  |
| :---: | :---: | :---: | :---: |
| Ada | Description | C | Fortran-77 |
| SEE | Exception on execution. | STRP | SAN |
| SRN | Replace with NULL. | SSDL | SDL |
| SRR | Return statement replacement. | SRSR | RSR |
| SGL | GOTO label replacement. | SGLR | GLR |
| SRE | Replace with EXIT. | SBR | $\sim$ |
| SWR | Replace WHILE with repeat-until. | SWDD | $\sim$ |
| SRW | Replace repeat-until with WHILE. | SDWD | $\sim$ |
| SZI | Zero iteration loop. | $\sim$ | der |
| SOI | One iteration loop. | $\sim$ | $\sim$ |
| SNI | N iteration loop. | SMTT | $\sim$ |
| SRI | Reverse iteration loop. | $\sim$ | $\sim$ |
| SES | END shift. | SMVB | $\sim$ |
| SCA | CASE alternative replacement. | SSWM | $\sim$ |
| SER | RAISE exception handler replacement. | $\sim$ | $\sim$ |
| Expression Modification Operators |  |  |  |
| Ada | Description | C | Fortran-77 |
| EAI | Absolute value insertion. | VDTR | ABS |
| ENI | Neg-absolute value insertion. | VDTR | ABS |
| EEZ | Exception on zero. | VDTR | ABS |
| EOR | Arithmetic operator replacement. | ORAN | AOR |
| ERR | Relational operator replacement. | ORRN | ROR |
| EMR | Membership test replacement. | $\sim$ | ~ |
| ELR | Logical operator replacement. | OBBN | LCR |
| EUI | Unary operator insertion. | Uuor | UOI |
| EUR | Unary operator replacement. | Uuor | $\sim$ |
| ESR | Subprogram operator replacement. | $\sim$ | $\sim$ |
| EDT | Domain twiddle. | VTWD | crp |
| EAR | Attribute replacement. | $\sim$ | $\sim$ |
| EEO | Exception on overflow. | $\sim$ | $\sim$ |
| EEU | Exception on underflow. | $\sim$ | $\sim$ |
| Coverage Operators |  |  |  |
| Ada | Description | C | Fortran-77 |
| CDE | Decision coverage. | Oior | lcr |
| CCO | Condition coverage. | $\sim$ | ror |
| CDC | Decision/condition coverage. | $\sim$ | lcr, ror |
| CMC | Multiple condition coverage. | $\sim$ | ror, lcr |
| Tasking Operators |  |  |  |
| Ada | Description | C | Fortran-77 |
| TEM | ENTRY statement modification. | $\sim$ | $\sim$ |
| TAR | ACCEPT statement replacement. | $\sim$ | $\sim$ |
| TSA | SELECT alternative replacement. | $\sim$ | $\sim$ |

## 5 ADA SELECTIVE MUTATION

The Fortran- 77 mutation system, Mothra [?], uses 22 mutation operators, of which the 6 most populous account for 40 to $60 \%$ of all mutants. Recent experimental research [?, ?] has indicated that of the 22
mutation operators used by Mothra, 17 of them (including the 6 most populous) seem to be in some sense redundant; that is, test sets that are generated to kill only mutants generated from the other 5 mutant operators are very effective in killing mutants generated from the 17. Selective mutation is an approximation technique that selects only mutants that are truly distinct from other mutants [?]. In experimental trials, selective mutation provides almost the same coverage as non-selective mutation, with significant reductions in cost.

Specifically, the results indicate that the mutation operators that replace all operands with all syntactically legal operands add very little to the effectiveness of mutation testing. Additionally, the mutation operators that modify entire statements add very little. The 5 selective operators for Fortran- 77 are ABS, which forces each arithmetic expression to take on the value 0 , a positive value, and a negative value, AOR, which replaces each arithmetic operator with every syntactically legal operator, LCR, which replaces each logical connector (AND and OR) with several kinds of logical connectors, ROR, which replaces relational operators with other relational operators, and UOI, which inserts unary operators in front of expressions. This report lists the mutation operators for Ada that should be included in the selective set.

### 5.1 List of Selective Operators

We leave out all operand replacement operators, and most of the statement operators. Most of the expression operators are included in the selective set. Because there has been no experience with tasking mutation operators, we leave them in. Further experimentation is needed to verify whether these are necessary.

| Expression Modification Operators |  |  |
| :--- | :--- | :---: |
| EAI | Absolute value insertion. |  |
| ENI | Neg-absolute value insertion. |  |
| EEZ | Exception on zero. |  |
| EOR | Arithmetic operator replacement. |  |
| ERR | Relational operator replacement. |  |
| EMR | Membership test replacement. |  |
| ELR | Logical operator replacement. |  |
| EUI | Unary operator insertion. |  |
| EUR | Unary operator replacement. |  |
| ESR | Subprogram operator replacement. |  |
| EEO | Exception on oVerflow. |  |
| EEU | Exception on underflow. |  |
| Coverage Operators |  |  |
| CMC | Multiple Condition coverage. |  |
| Tasking Operators |  |  |
| TEM | ENTRY statement modification. |  |
| TAR | ACCEPT statement replacement. |  |
| TSA | SELECT alternative replacement. |  |

## 6 COMPREHENSIVE ADA MUTATION EXAMPLE

In this section, we present an example of a mutated Ada program. We show a small Ada program with all mutants displayed "in-line", that is, with the changes shown in the text of the program. The program
reads two matrices, adds and multiplies them together, and prints the results. The lines in the program are numbered (for simplicity, all text are numbered, including comments and blanks), and mutants are shown in the program.

Each mutant is represented by the modified statement in the program. The mutated statement is shown just below the original statement. Each mutated statement includes the mutation operator name (e.g., OVV, OVC), a unique integer that identifies the mutant, and the string -->. For example, line 43 contains the header of a FOR loop. The first mutant of that line is shown immediately below it. It is an OCC (OperandConstant replaced by a Constant) mutant, number 42, and the constant 1 is replaced by the named constant MATSIZE. The program has 71 executable statements and 559 mutants.

```
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4 1
```

-- Programmer : Jeff Offutt

```
-- Programmer : Jeff Offutt
Mrogram : Jeff Offutt - 
Mrogram : Jeff Offutt - 
Program : Matrix --
Program : Matrix --
-- Purpose : Add and multiply two matrices. --
-- Purpose : Add and multiply two matrices. --
-- Date : 2/21/94 --
-- Date : 2/21/94 --
-----------------------------------------------------------------
-----------------------------------------------------------------
WITH Text_IO; USE Text_IO;
PROCEDURE Matrix IS
    -- Package instantiations.
    PACKAGE Int_IO IS NEW Integer_IO (Integer);
    USE Int_IO;
    -- Constant declarations.
    INFILE : CONSTANT String := "p1.in";
    OUTFILE : CONSTANT String := "p1.out";
    MATSIZE : CONSTANT Integer := 3;
    -- Type declarations.
    TYPE Matrix_Type IS ARRAY (1..MATSIZE,1..MATSIZE) OF Integer;
    -- Variable declarations.
    Input_File, Output_File : File_Type;
    Matrix_In1, Matrix_In2 : Matrix_Type;
    Mat_Sum, Mat_Prod : Matrix_Type;
    ---------------------------------------------------------------
    -- Mrocedure : Read_Mat (Input_File : IN File_Type; 隹 M Matrix_Type) --
    -- Purpose : Read a matrix from a file. --
    -- Params : Input_File : The file to read from. --
    -- : Mat : The matrix to read into. --
    -- Pre : Input_File must be open --
    -- : The matrix is MATSIZE lines --
    -- : MATSIZE numbers per row --
    ---------------------------------------------------------------
    PROCEDURE Read_Mat (Input_File : IN File_Type; Mat : OUT Matrix_Type) IS
        ==>szi_lci1 : Natural;
        ==>szi_lci2 : Natural;
```

```
```

    ==>szi_lci1 := 0; -- Initialize loop counter.
    ```
```

    ==>szi_lci1 := 0; -- Initialize loop counter.
    FOR i IN 1..MATSIZE LOOP
    FOR i IN 1..MATSIZE LOOP
    -->FOR i IN MATSIZE..MATSIZE LOOP
    -->FOR i IN MATSIZE..MATSIZE LOOP
    -->FOR i IN 1..1 LOOP
    -->FOR i IN 1..1 LOOP
    -->Except_On_Exec (SEE+0); -- First SEE mutant=320 + offset.
    -->Except_On_Exec (SEE+0); -- First SEE mutant=320 + offset.
    -->NULL; -- Replaces statements 43,44,45,46,47
    -->NULL; -- Replaces statements 43,44,45,46,47
    -->RETURN; -- Replaces statements 43,44,45,46,47
    -->RETURN; -- Replaces statements 43,44,45,46,47
    -->EXIT; -- Replaces statements 43,44,45,46,47
    -->EXIT; -- Replaces statements 43,44,45,46,47
    -->FOR i IN REVERSE 1..MATSIZE LOOP
    -->FOR i IN REVERSE 1..MATSIZE LOOP
    4 4
4 4
OCC 44
OCC 44
OCC 45
OCC 45
OCC 46
OCC 46
OCC 47
OCC 47
SEE }31
SEE }31
SRN 328
SRN 328
SRR 359
SRR 359
SRE }38
SRE }38
SRI }44
SRI }44
4 5
4 5
OCV 26
OCV 26
OCC 48
OCC 48
OCC 49
OCC 49
OCC 50
OCC 50
OCC 51
OCC 51
OCC 52
OCC 52
OCC 53
OCC 53
OCA 158
OCA 158
OCA 159
OCA 159
OAN 230
OAN 230
OAN 231
OAN 231
OAN 232
OAN 232
OAN 233
OAN 233
SEE 314
SEE 314
SRN 329
SRN 329
SRR 360
SRR 360
SRE 385
SRE 385
4 6
4 6
SZI }41
SZI }41
SOI }42
SOI }42
SNI }43
SNI }43
4 7
4 7
SZI 415
SZI 415
SOI }42
SOI }42
SNI }43
SNI }43
4 8
4 8
4 9
4 9
50
50
51
51
52
52
53
53
4 3
4 3
OCC 42
OCC 42
OCC 43
OCC 43
SEE }31
SEE }31
SRN 327
SRN 327
SRR 358
SRR 358
SRE }38
SRE }38
SRI 441
SRI 441
SZI

```
```

    SZI
    ```
```



```
    SZI
```

```
    SZI
```

```
    SZI
```

```
    ==>szi_lci2 := 0; -- Initialize loop counter.
```

    ==>szi_lci2 := 0; -- Initialize loop counter.
    FOR j IN 1..MATSIZE LOOP
    FOR j IN 1..MATSIZE LOOP
    -->FOR j IN MATSIZE..MATSIZE LOOP
    -->FOR j IN MATSIZE..MATSIZE LOOP
    -->FOR j IN i..MATSIZE LOOP
    -->FOR j IN i..MATSIZE LOOP
    -->FOR j IN 1..1 LOOP
    -->FOR j IN 1..1 LOOP
    -->FOR j IN 1..i LOOP
    -->FOR j IN 1..i LOOP
    -->Except_On_Exec (SEE+1);
    -->Except_On_Exec (SEE+1);
    -->NULL; -- Replaces statements 44,45,46
    -->NULL; -- Replaces statements 44,45,46
    -->RETURN; -- Replaces statements 44,45,46
    -->RETURN; -- Replaces statements 44,45,46
    -->EXIT; -- Replaces statements 44,45,46
    -->EXIT; -- Replaces statements 44,45,46
    -->FOR j IN REVERSE 1..MATSIZE LOOP
    -->FOR j IN REVERSE 1..MATSIZE LOOP
            Get (Input_File, Mat (i,j));
            Get (Input_File, Mat (i,j));
            -->Get (Output_File, Mat (i,j));
            -->Get (Output_File, Mat (i,j));
            -->Get (Input_File, Mat (1,j));
            -->Get (Input_File, Mat (1,j));
            -->Get (Input_File, Mat (j,j));
            -->Get (Input_File, Mat (j,j));
            -->Get (Input_File, Mat (MATSIZE,j));
            -->Get (Input_File, Mat (MATSIZE,j));
            -->Get (Input_File, Mat (i,1));
            -->Get (Input_File, Mat (i,1));
            -->Get (Input_File, Mat (i,i));
            -->Get (Input_File, Mat (i,i));
            -->Get (Input_File, Mat (i,MATSIZE));
            -->Get (Input_File, Mat (i,MATSIZE));
            -->Get (Input_File, Mat (Mat(i,j),j));
            -->Get (Input_File, Mat (Mat(i,j),j));
            -->Get (Input_File, Mat (i,Mat(i,j)));
            -->Get (Input_File, Mat (i,Mat(i,j)));
            -->Get (Input_File, Matrix_In1 (i,j));
            -->Get (Input_File, Matrix_In1 (i,j));
            -->Get (Input_File, Matrix_In2 (i,j));
            -->Get (Input_File, Matrix_In2 (i,j));
            -->Get (Input_File, Mat_Sum (i,j));
            -->Get (Input_File, Mat_Sum (i,j));
            -->Get (Input_File, Mat_Prod (i,j));
            -->Get (Input_File, Mat_Prod (i,j));
            -->Except_On_Exec (SEE+2);
            -->Except_On_Exec (SEE+2);
            -->NULL;
            -->NULL;
            -->RETURN;
            -->RETURN;
            -->EXIT;
            -->EXIT;
        END LOOP;
        END LOOP;
    -->IF (szi_lci1 = 0) RAISE Mut_Trap;
    -->IF (szi_lci1 = 0) RAISE Mut_Trap;
    -->IF (szi_lci1 = 1) RAISE Mut_Trap;
    -->IF (szi_lci1 = 1) RAISE Mut_Trap;
    -->IF (szi_lci1 > 1) RAISE Mut_Trap;
    -->IF (szi_lci1 > 1) RAISE Mut_Trap;
    END LOOP;
    END LOOP;
    -->IF (szi_lci2 = 0) RAISE Mut_Trap;
    -->IF (szi_lci2 = 0) RAISE Mut_Trap;
    -->IF (szi_lci2 = 1) RAISE Mut_Trap;
    -->IF (szi_lci2 = 1) RAISE Mut_Trap;
    -->IF (szi_lci2 > 1) RAISE Mut_Trap;
    -->IF (szi_lci2 > 1) RAISE Mut_Trap;
    END Read_Mat;
END Read_Mat;
-----------------------------------------------------------
-----------------------------------------------------------
-- Procedure : Write_Mat (Output_File : IN File_Type; --
-- Procedure : Write_Mat (Output_File : IN File_Type; --
-- Mat : IN Matrix_Type) --
-- Mat : IN Matrix_Type) --
-- Purpose : Write a matrix to a file.

```
-- Purpose : Write a matrix to a file.
```

| 54 | -- Params : Output_File : The file to write to. | -- |
| :---: | :---: | :---: |
| 55 | : Mat : The matrix to write from. |  |
| 56 | -- Pre : Output_File must be open | -- |
| 57 |  |  |
| 58 | PROCEDURE Write_Mat (Output_File : IN File_Type; Mat | IN Matrix_Type) IS |
| SZI | ==>szi_lci3 : Natural; |  |
| SZI | ==>szi_lci4 : Natural; |  |
| 59 | BEGIN |  |
| SZI | ==>szi_lci3 := 0; -- Initialize loop counter. |  |
| 60 | FOR i IN 1..MATSIZE LOOP |  |
| OCC 54 | -->FOR i IN MATSIZE..MATSIZE LOOP |  |
| OCC 55 | -->FOR i IN 1..1 LOOP |  |
| SEE 315 | -->Except_On_Exec (SEE+3); |  |
| SRN 330 | -->NULL; -- Replaces statements 60-65 |  |
| SRR 361 | -->RETURN; -- Replaces statements 60-65 |  |
| SRE 386 | -->EXIT; -- Replaces statements 60-65 |  |
| SRI 443 | -->FOR i IN REVERSE 1..MATSIZE LOOP |  |
| EDT 522 | -->FOR i IN O..MATSIZE LOOP |  |
| EDT 523 | -->FOR i IN 2..MATSIZE LOOP |  |
| EDT 524 | -->FOR i IN 0..mATSIZE-1 LOOP |  |
| EDT 525 | -->FOR i IN O..MATSIZE+1 LOOP |  |
| SZI | ==>szi_lci4 := 0; -- Initialize loop counter. |  |
| 61 | FOR j IN 1..MATSIZE LOOP |  |
| OCC 56 | -->FOR j IN MATSIZE..MATSIZE LOOP |  |
| OCC 57 | -->FOR j IN i..MATSIZE LOOP |  |
| OCC 58 | -->FOR j IN 1..1 LOOP |  |
| OCC 59 | -->FOR j IN 1..i LOOP |  |
| SEE 316 | -->Except_On_Exec (SEE+4); |  |
| SRN 331 | -->NULL; -- Replaces statements 61-63 |  |
| SRR 362 | -->RETURN; -- Replaces statements 61-63 |  |
| SRE 387 | -->EXIT; -- Replaces statements 61-63 |  |
| SRI 444 | -->FOR j IN REVERSE 1..MATSIZE LOOP |  |
| SES 450 | -->END LOOP; -- Line 63 moved below 61. |  |
| EDT 526 | -->FOR j IN O..MATSIZE LOOP |  |
| EDT 527 | -->FOR j IN 2..MATSIZE LOOP |  |
| EDT 528 | -->FOR j IN 1..MATSIZE-1 LOOP |  |
| EDT 529 | -->FOR j IN 1..MATSIZE+1 LOOP |  |
| 62 | Put (Output_File, Mat (i,j)) ; |  |
| OCV 27 | -->Put (Input_File, Mat (i,j)) ; |  |
| OCC 60 | -->Put (Output_File, Mat (1,j)) ; |  |
| OCC 61 | -->Put (Output_File, Mat (j,j)) ; |  |
| OCC 62 | -->Put (Output_File, Mat (MATSIZE,j)); |  |
| OCC 63 | -->Put (Output_File, Mat (i,1)) |  |
| OCC 64 | -->Put (Output_File, Mat (i,i)); |  |
| OCC 65 | -->Put (Output_File, Mat (i,MATSIZE)); |  |
| OCA 160 | -->Put (Output_File, Mat (Mat (i,j),j)) ; |  |
| OCA 161 | -->Put (Output_File, Mat (i,Mat (i,j))) ; |  |
| OAC 201 | -->Put (Output_File, i) ; |  |
| OAC 202 | -->Put (Output_File, j) ; |  |
| OAC 203 | -->Put (Output_File, 1) ; |  |
| OAN 234 | -->Put (Output_File, Matrix_In1 (i,j)) ; |  |
| OAN 235 | -->Put (Output_File, Matrix_In2 (i,j)) ; |  |
| OAN 236 | -->Put (Output_File, Mat_Sum (i,j)) ; |  |


| OAN 237 | -->Put (Output_File, Mat_Prod (i,j)) ; |  |
| :---: | :---: | :---: |
| SEE 317 | -->Except_On_Exec (SEE+5); |  |
| SRN 332 | -->NULL; |  |
| SRR 363 | -->RETURN; |  |
| SRE 388 | -->EXIT; |  |
| EAI 458 | -->Put (Output_File, ABS (Mat (i,j))) ; |  |
| ENI 468 | -->Put (Output_File, - ABS (Mat (i,j))) ; |  |
| EEZ 478 | -->Put (Output_File, EEZ(Mat (i,j))) ; |  |
| EUI 506 | -->Put (Output_File, -Mat (i,j)); |  |
| EDT 530 | -->Put (Output_File, Mat (i,j)+1) ; |  |
| EDT 531 | -->Put (Output_File, Mat (i,j)-1) ; |  |
| 63 | END LOOP; |  |
| SZI 416 | -->IF (szi_lci3 = 0) RAISE Mut_Trap; |  |
| SOI 425 | -->IF (szi_lci3 = 1) RAISE Mut_Trap; |  |
| SNI 434 | -->IF (szi_lci3 > 1) RAISE Mut_Trap; |  |
| 64 | New_Line (Output_File); |  |
| OCV 28 | -->New_Line (Input_File); |  |
| SEE 318 | -->Except_On_Exec (SEE+6); |  |
| SRN 333 | -->NULL; |  |
| SRR 364 | -->RETURN; |  |
| SRE 389 | -->EXIT; |  |
| SES 451 | -->END LOOP; -- Line 63 moved below 64. |  |
| 65 | END LOOP; |  |
| SZI 417 | -->IF (szi_lci4 = 0) RAISE Mut_Trap; |  |
| SOI 426 | -->IF (szi_lci4 = 1) RAISE Mut_Trap; |  |
| SNI 435 | -->IF (szi_lci4 > 1) RAISE Mut_Trap; |  |
| 66 | END Write_Mat; |  |
| 67 |  |  |
| 68 |  |  |
| 69 | -- Function : Add_Mat (Mat1, Mat2 : IN Matrix_Type) | -- |
| 70 | -- Purpose : Add two matrices |  |
| 71 | -- Params : Mat1, Mat2 : The matrices to add. | -- |
| 72 | -- Return : The matrix sum. | -- |
| 73 | -- Pre : Mat1 and Mat2 are initialized. | -- |
| 74 |  |  |
| 75 | FUNCTION Add_Mat (Mat1, Mat2 : IN Matrix_Type) RETURN Matrix_Type IS |  |
| 76 | rslt_mat : Matrix_Type; |  |
| SZI | ==>szi_lci5 : Natural; |  |
| SZI | ==>szi_lci6 : Natural; |  |
| 77 | BEGIN |  |
| SZI | ==>szi_lci5 := 0; -- Initialize loop counter. |  |
| 78 | FOR i IN 1..MATSIZE LOOP |  |
| OCC 66 | -->FOR i IN MATSIZE..MATSIZE LOOP |  |
| OCC 67 | -->FOR i IN 1..1 LOOP |  |
| SEE 319 | -->Except_On_Exec (SEE+7); |  |
| SRN 334 | -->NULL; -- Replaces statements 78-82 |  |
| SRR 365 | -->RETURN; -- Replaces statements 78-82 |  |
| SRR 366 | -->RETURN (rslt_mat); -- Replaces statements 78-82 |  |
| SRE 390 | -->EXIT; -- Replaces statements 78-82 |  |
| SRI 445 | -->FOR i IN REVERSE 1..MATSIZE LOOP |  |
| EDT 532 | -->FOR i IN O..MATSIZE LOOP |  |
| EDT 533 | -->FOR i IN 2..MATSIZE LOOP |  |
| EDT 534 | -->FOR i IN 1..MATSIZE-1 LOOP |  |


| SZI | ==>szi_lci6 := 0; -- Initialize loop counter. |
| :---: | :---: |
| 79 | FOR j IN 1..MATSIZE LOOP |
| OCC 68 | -->FOR j IN MATSIZE..MATSIZE LOOP |
| OCC 69 | -->FOR j IN i..MATSIZE LOOP |
| OCC 70 | -->FOR j IN 1..1 LOOP |
| OCC 71 | -->FOR j IN 1..i LOOP |
| SEE 320 | -->Except_On_Exec (SEE+8); |
| SRN 335 | -->NULL; -- Replaces statements 79-81 |
| SRR 367 | -->RETURN; -- Replaces statements 79-81 |
| SRR 368 | -->RETURN (rslt_mat); -- Replaces statements 79-81 |
| SRE 391 | -->EXIT; -- Replaces statements 79-81 |
| SRI 446 | -->FOR j IN REVERSE 1..MATSIZE LOOP |
| SES 452 | -->END LOOP; -- Line 81 moved below 79. |
| EDT 536 | -->FOR j IN O..MATSIZE LOOP |
| EDT 537 | -->FOR j IN 2..MATSIZE LOOP |
| EDT 538 | -->FOR j IN 1..MATSIZE-1 LOOP |
| EDT 539 | -->FOR j IN 1...MATSIZE+1 LOOP |


| lt_mat (i,j) := Mat1 (i,j) + Mat2 (i,j) ; |  |
| :---: | :---: |
| >rslt_ma | $(1, j) ~:=~ M a t 1 ~(i, j) ~+~ M a t 2 ~(i, j) ~ ; ~$ |
| >rslt_m | (j,j) := Mat1 (i,j) + Mat2 (i,j) ; |
| >rslt_m | (i,MATSIZE) := Mat1 (i, j) + Mat2 (i,j) ; |
| >rslt_m | $(i, 1) ~:=~ M a t 1 ~(i, j) ~+~ M a t 2 ~(i, j) ~ ; ~$ |
| -->rslt_m | (i,i) := Mat1 (i,j) + Mat2 (i,j) ; |
| >rslt_m | $(i, j) ~:=~ M a t 1 ~(M A T S I Z E, ~ j) ~+~ M a t 2 ~(i, j) ~ ; ~$ |
| >rslt_m | $(\mathrm{i}, \mathrm{j}):=\mathrm{Mat1}(1, \mathrm{j})+\mathrm{Mat2}(\mathrm{i}, \mathrm{j})$; |
| -->rslt_ma | $(\mathrm{i}, \mathrm{j}):=\mathrm{Mat1}(\mathrm{j}, \mathrm{j})+\mathrm{Mat2}(\mathrm{i}, \mathrm{j})$; |
| >rslt_m | $(i, j) ~:=~ M a t 1 ~(i, M A T S I Z E) ~+~ M a t 2 ~(i, j) ~ ; ~$ |
| >rslt_m | $(i, j) ~:=~ M a t 1 ~(i, 1) ~+~ M a t 2 ~(i, j) ~ ; ~$ |
| >rslt_ma | $(i, j) ~:=~ M a t 1 ~(i, i) ~+~ M a t 2 ~(i, j) ~ ; ~$ |
| >rslt_m | $(\mathrm{i}, \mathrm{j}) ~:=$ Mat1 ( $\mathrm{i}, \mathrm{j}$ ) + Mat2 (MATSIZE, j$)$; |
| >rslt_ma | $(i, j) ~:=~ M a t 1 ~(i, j) ~+~ M a t 2 ~(1, j) ~ ; ~$ |
| >rslt_ma | $(i, j) ~:=~ M a t 1 ~(i, j) ~+~ M a t 2 ~(j, ~ j) ~ ; ~$ |
| >rslt_ma | $(i, j) ~:=~ M a t 1 ~(i, j) ~+~ M a t 2 ~(i, M A T S I Z E) ~ ; ~$ |
| >rslt_m | $(i, j) ~:=~ M a t 1 ~(i, j) ~+~ M a t 2 ~(i, 1) ~ ; ~$ |
| >rslt_ma | (i,j) := Mat1 (i,j) + Mat2 (i,i) ; |
| >rslt_m |  |
| >rslt_m | (i,rslt_mat(i,j)) := Mat1 (i,j) + Mat2 (i, |
| >rslt_ma | (Mat1 (i, j) , j) := Mat1 (i, j) + Mat2 (i, j) ; |
| >rslt_m | (i,j) := Mat1 (i, Mat1 (i, j) ) + Mat2 (i,j) ; |
| >rslt_ma | (Mat2 (i, j) , j) := Mat1 (i, j) + Mat2 (i,j) ; |
| >rslt_ma | (i,Mat2 (i, j) ) := Mat1 (i,j) + Mat2 (i,j) ; |
| >rslt_ma | (i,j) := Mat1 (rslt_mat (i,j), j) + Mat2 (i,j) |
| -->rslt_ma | (i,j) := Mat1 (i,rslt_mat(i,j)) + Mat2 (i, |
| $\rightarrow$ rslt_ma | $(\mathrm{i}, \mathrm{j}) ~:=~ M a t 1 ~(M a t 1(i, j), j)+\operatorname{Mat2}(\mathrm{i}, \mathrm{j})$; |
| >rslt_ma |  |
| >rslt_ma | $(\mathrm{i}, \mathrm{j}) ~:=~ M a t 1 ~(M a t 2(i, j), j) ~+~ M a t 2 ~(i, j) ~ ; ~$ |
| ->rslt_ma | (i,j) := Mat1 (i, Mat2 (i, j) ) + Mat2 (i,j) ; |
| >rslt_m | (i,j) := Mat1 (i,j) + Mat2 (rslt_mat (i,j) , ${ }^{\text {( }}$ ) |
| -->rslt_ma | (i, j) := Mat1 (i, j) + Mat2 (i,rslt_mat (i,j) |
| >rslt_ma | (i,j) := Mat1 (i,j) + Mat2 (Mat1 (i, j) , j$)$; |
| -->rslt_m | (i,j) := Mat1 (i,j) + Mat2 (i, Mat1 (i,j)) ; |
| >rslt_ma | $(\mathrm{i}, \mathrm{j}):=\mathrm{Mat1}(\mathrm{i}, \mathrm{j})+\mathrm{Mat2}(\operatorname{Mat2}(\mathrm{i}, \mathrm{j}), \mathrm{j})$; |

OCA 179
OAC 204
OAC 205
OAC 206
OAC 207
OAC 208
OAC 209
OAC 210
OAC 211
OAA 222
OAA 223
OAA 224
OAA 225
OAN 238
OAN 239
OAN 240
OAN 241
OAN 242
OAN 243
OAN 244
OAN 245
OAN 246
OAN 247
OAN 248
OAN 249
OAN 250
OAN 251
OAN 252
OAN 253
SEE 321
SRN 336
SRR 369
SRR 370
SRE 392
EAI 459
EAI 460
EAI 461
ENI 469
ENI 470
ENI 471
EEZ 479
EEZ 480
EEZ 481
EOR 488
EOR 489
EOR 490
EOR 491
EOR 492
EOR 493
EUI 507
EUI 508
EUI 509
EDT 540
EDT 541
EEO 556
EEU 558


```
    ==>szi_lci8 := 0; -- Initialize loop counter.
    FOR j IN 1..MATSIZE LOOP
    -->FOR j IN tmp_sum..MATSIZE LOOP
    -->FOR j IN 1..tmp_sum LOOP
    -->FOR j IN MATSIZE..MATSIZE LOOP
    -->FOR j IN O..MATSIZE LOOP
    -->FOR j IN i..MATSIZE LOOP
    -->FOR j IN 1..1 LOOP
    -->FOR j IN 1..0 LOOP
    -->FOR j IN 1..i LOOP
    -->Except_On_Exec (SEE+11);
    -->NULL; -- Replaces statements 98-104
    -->RETURN; -- Replaces statements 98-104
    -->RETURN (rslt_mat); -- Replaces statements 98-104
    -->EXIT; -- Replaces statements 98-104
    -->FOR j IN REVERSE 1..MATSIZE LOOP
        FOR j IN O..MATSIZE LOOP -- equiv to OCC 95
    -->FOR j IN 2..MATSIZE LOOP
    -->FOR j IN 1..MATSIZE-1 LOOP
    -->FOR j IN 1..MATSIZE+1 LOOP
99
    OVA 14
    OVA 15
    OVA 16
    OCV 33
    OCC 100
    OCC 101
    OCC 102
    OCC 103
    SRN }33
    SRR }37
    SRR }37
    SRE }39
    EDT }54
    SZI
100
    OCV 34
    OCV 35
    OCC 104
    OCC 105
    OCC 106
    OCC 107
    OCC 108
    OCC 109
    OCC 110
    OCC 111
    OCC 112
    SEE }32
    SRN 340
    SRR }37
    SRR }37
    SRE }39
    tmp_sum := 0;
    -->Mat1(i,k) := 0;
    -->Mat2(k,j) := 0;
    -->rslt_mat(i,j) := 0;
    -->tmp_sum := tmp_sum;
    -->tmp_sum := i;
    -->tmp_sum := j;
    -->tmp_sum := 1;
    -->tmp_sum := MATSIZE;
    -->NULL;
    -->RETURN;
    -->RETURN (rslt_mat);
    -->EXIT;
    -->tmp_sum := -1;
    ==>szi_lci9 := 0; -- Initialize loop counter.
    FOR k IN 1..MATSIZE LOOP
    -->FOR k IN tmp_sum..MATSIZE LOOP
    -->FOR k IN 1..tmp_sum LOOP
    -->FOR k IN MATSIZE..MATSIZE LOOP
    -->FOR k IN O..MATSIZE LOOP
    -->FOR k IN i..MATSIZE LOOP
    -->FOR k IN j..MATSIZE LOOP
    -->FOR k IN 1..1 LOOP
    -->FOR k IN 1..0 LOOP
    -->FOR k IN 1..i LOOP
    -->FOR k IN 1..j LOOP
    -->FOR k IN 1..MATSIZE LOOP
    -->Except_On_Exec (SEE+12);
    -->NULL; -- Replaces statements 100-102
    -->RETURN; -- Replaces statements 100-102
    -->RETURN (rslt_mat); -- Replaces statements 100-102
    -->EXIT; -- Replaces statements 100-102
```

SRI 449
SES 454
EDT 549
EDT 550
EDT 551
-->FOR k IN REVERSE 1..MATSIZE LOOP
-->END LOOP; -- Line 102 moved below 100.
-->FOR k IN 2..MATSIZE LOOP
-->FOR k IN 1..MATSIZE-1 LOOP
-->FOR k IN 1..MATSIZE+1 LOOP

OVA 17
OVA 18
OVA 19
OVA 20
OVA 21
OVA 22
OCV 36
OCV 37
OCV 38
OCV 39
OCC 113
OCC 114
OCC 115
OCC 116
OCC 117
OCC 118
OCC 119
OCC 120
OCC 121
OCC 122
OCC 123
OCC 124
OCC 125
OCC 126
OCC 127
OCC 128
OCC 129
OCC 130
OCC 131
OCC 132
OVC 153
OVC 154
OVC 155
OVC 156
OVC 157
OCA 180
OCA 181
OCA 182
OCA 183
OCA 184
OCA 185
OCA 186
OCA 187
OCA 188
OCA 189
OCA 190
OCA 191
OAV 198
OAV 199

```
tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (k,j);
-->Mat1(i,k) := tmp_sum + Mat1 (i,k) * Mat2 (k,j);
-->Mat2(k,j) := tmp_sum + Mat1 (i,k) * Mat2 (k,j);
-->rslt_mat(i,j) := tmp_sum + Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := Mat1(i,k) + Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := Mat2(k,j) + Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := rslt_mat(i,j) + Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (tmp_sum,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,tmp_sum) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (tmp_sum,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (k,tmp_sum);
-->tmp_sum := tmp_sum + Mat1 (j,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (1,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (MATSIZE,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (0,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (k,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,i) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,j) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,1) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,MATSIZE) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,0) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (i,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (j,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (1,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (MATSIZE,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (0,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (k,i);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (k,1);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (k,MATSIZE);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (k,0);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (k,k);
-->tmp_sum := i + Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := j + Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := 1 + Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := MATSIZE + Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := 0 + Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (Mat1(i,k),k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,Mat1(i,k)) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (Mat2(k,j),k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,Mat2(k,j)) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (rslt_mat(i,j),k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,rslt_mat(i,j)) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (Mat1(i,k),j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (k,Mat1(i,k));
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (Mat2(k,j),j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (k,Mat2(k,j));
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (rslt_mat(i,j),j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat2 (k,rslt_mat(i,j));
-->tmp_sum := tmp_sum + tmp_sum * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * tmp_sum;
```

OAC 212
OAC 213
OAC 214
OAC 215
OAC 216
OAC 217
OAC 218
OAC 219
OAC 220
OAC 221
OAA 226
OAA 227
OAA 228
OAA 229
OAN 260
OAN 261
OAN 262
OAN 263
OAN 264
OAN 265
OAN 266
OAN 267
OAN 268
OAN 269
OAN 270
OAN 271
SRN 341
SRR 380
SRR 381
SRE 397
EAI 462
EAI 463
EAI 464
EAI 465
EAI 466
ENI 472
ENI 473
ENI 474
ENI 475
ENI 476
EEZ 482
EEZ 483
EEZ 484
EEZ 485
EEZ 486
EOR 494
EOR 495
EOR 496
EOR 497
EOR 498
EOR 499
EOR 500
EOR 501
EOR 502
EOR 503
EOR 504

```
-->tmp_sum := tmp_sum + i * Mat2 (k,j);
-->tmp_sum := tmp_sum + j * Mat2 (k,j);
-->tmp_sum := tmp_sum + 1 * Mat2 (k,j);
-->tmp_sum := tmp_sum + 0 * Mat2 (k,j);
-->tmp_sum := tmp_sum + MATSIZE * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * i;
-->tmp_sum := tmp_sum + Mat1 (i,k) * j;
-->tmp_sum := tmp_sum + Mat1 (i,k) * 1;
-->tmp_sum := tmp_sum + Mat1 (i,k) * 0;
-->tmp_sum := tmp_sum + Mat1 (i,k) * MATSIZE;
-->tmp_sum := tmp_sum + rslt_mat (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat2 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * rslt_mat (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat1 (k,j);
-->tmp_sum := tmp_sum + Matrix_In1 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Matrix_In2 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat_Sum (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat_Prod (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat2 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + rslt_mat (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Matrix_In1 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Matrix_In1 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat_Sum (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat_Prod (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * Mat1 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * rslt_mat (k,j);
-->NULL;
-->RETURN;
-->RETURN (rslt_mat);
-->EXIT;
-->tmp_sum := ABS(tmp_sum + Mat1 (i,k) * Mat2 (k,j));
-->tmp_sum := ABS(tmp_sum) + Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + ABS(Mat1 (i,k) * Mat2 (k,j));
-->tmp_sum := tmp_sum + ABS(Mat1 (i,k)) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * ABS(Mat2 (k,j));
-->tmp_sum := -ABS(tmp_sum + Mat1 (i,k) * Mat2 (k,j));
-->tmp_sum := -ABS(tmp_sum) + Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + -ABS(Mat1 (i,k) * Mat2 (k,j));
-->tmp_sum := tmp_sum + -ABS(Mat1 (i,k)) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * -ABS(Mat2 (k,j));
-->tmp_sum := EEZ(tmp_sum + Mat1 (i,k) * Mat2 (k,j));
-->tmp_sum := EEZ(tmp_sum) + Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + EEZ(Mat1 (i,k) * Mat2 (k,j));
-->tmp_sum := tmp_sum + EEZ(Mat1 (i,k)) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) * EEZ(Mat2 (k,j));
-->tmp_sum := tmp_sum - Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum * Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum / Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum MOD Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum REM Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum ** Mat1 (i,k) * Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) + Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) - Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) / Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) MOD Mat2 (k,j);
-->tmp_sum := tmp_sum + Mat1 (i,k) REM Mat2 (k,j);
```



| SZI 422 | -->IF (szi_lci9 = 0) RAISE Mut_Trap; |
| :---: | :---: |
| SOI 431 | -->IF (szi_lci9 = 1) RAISE Mut_Trap; |
| SNI 440 | -->IF (szi_lci9 > 1) RAISE Mut_Trap; |
| 106 | RETURN (rslt_mat); |
| OAN 276 | -->RETURN (Matrix_In1); |
| OAN 277 | -->RETURN (Matrix_In2); |
| OAN 278 | -->RETURN (Mat_Sum) ; |
| OAN 279 | -->RETURN (Mat_Prod); |
| OAN 280 | -->RETURN (Mat1); |
| OAN 281 | -->RETURN (Mat2); |
| --SEE -- | no SEE mutant, function must have a RETURN. |
| --SRN -- | no SRN mutant, function must have a RETURN. |
| SRR 382 | -->RETURN; |
| --SRE - | no SRE mutant, function must have a RETURN. |
| SES 457 | -->END LOOP; -- Line 105 moved below 106. |
| 107 | END Multiply_Mat; |
| 108 |  |
| 109 |  |
| 110 |  |
| 111 | Main body of matrix |
| 112 | Open files, read, and and multiply the two matrices. |
| 113 | Close the files. |
| 114 |  |
| 115 B | GIN -- Matrix |
| 116 | Open (Input_File, In_File, INFILE) ; |
| OVV | -->Open (Output_File, In_File, INFILE); |
| OCC 141 | -->Open (Input_File, In_File, OUTFILE); |
| OCC 142 | -->Open (Input_File, In_File, "Sum of Matrices"); |
| OCC 143 | -->Open (Input_File, In_File, "Product of Matrices"); |
| SEE 326 | -->Except_On_Exec (SEE+14); |
| SRN 342 | -->NULL; |
| SRE 398 | -->EXIT; |
| 117 | Create (Output_File, Out_File, OUTFILE); |
| OVV 2 | -->Create (Input_File, Out_File, OUTFILE); |
| OCC 144 | -->Create (Output_File, Out_File, INFILE); |
| OCC 145 | -->Create (Output_File, Out_File, "Sum of Matrices"); |
| OCC 146 | -->Create (Output_File, Out_File, "Product of Matrices"); |
| SRN 343 | -->NULL; |
| SRE 399 | -->EXIT; |
| 118 |  |
| 119 | Read_Mat (Input_File, Matrix_In1) |
| OVV 3 | -->Read_Mat (Output_File, Matrix_In1) ; |
| OAN 282 | -->Read_Mat (Input_File, Matrix_In2); |
| OAN 283 | -->Read_Mat (Input_File, Mat_Sum) ; |
| OAN 284 | -->Read_Mat (Input_File, Mat_Prod); |
| SRN 344 | -->NULL; |
| SRE 400 | -->EXIT; |
| ESR 516 | -->Write_Mat (Input_File, Matrix_In1) ; |
| 120 | Read_Mat (Input_File, Matrix_In2) |
| OVV 4 | -->Read_Mat (Output_File, Matrix_In2) ; |
| OAN 285 | -->Read_Mat (Input_File, Matrix_In1) ; |
| OAN 286 | -->Read_Mat (Input_File, Mat_Sum) ; |
| OAN 287 | -->Read_Mat (Input_File, Mat_Prod); |
| SRN 345 | -->NULL; |



```
    OAN 304 -->Mat_Prod := Multiply_Mat (Matrix_Sum, Matrix_In2);
    OAN 305 -->Mat_Prod := Multiply_Mat (Matrix_Prod, Matrix_In2);
    OAN 306 -->Mat_Prod := Multiply_Mat (Matrix_In1, Matrix_In1);
    OAN 307 -->Mat_Prod := Multiply_Mat (Matrix_In1, Matrix_Sum);
    OAN 308 -->Mat_Prod := Multiply_Mat (Matrix_In1, Matrix_Prod);
    SRN 352 -->NULL;
    SRE 408 -->EXIT;
    ESR 520 -->Mat_Prod := Add_Mat (Matrix_In1, Matrix_In2);
134 Put (Output_File, "Product of matrices:");
    OCC 150 -->Put (Output_File, OUTFILE);
    OCC 151 -->Put (Output_File, INFILE);
    OCC 152 -->Put (Output_File, "Sum of matrices:");
    SRN 353 -->NULL;
    SRE 409 -->EXIT;
1 3 5
    -->New line
    SRN 354 -->NULL;
    SRE 410 -->EXIT;
136 Write_Mat (Output_File, Mat_Prod);
    OVV 11 -->Write_Mat (Input_File, Mat_Prod);
    OAN 309 -->Write_Mat (Output_File, Matrix_In1);
    OAN 310 -->Write_Mat (Output_File, Matrix_In2);
    OAN 311 -->Write_Mat (Output_File, Mat_Sum);
    SRN 355 -->NULL;
    SRE 411 -->EXIT;
    ESR 521 -->Read_Mat (Output_File, Mat_Prod);
137 New_line (Output_File);
    OVV 12 -->New_line (Input_File);
    SRN 356 -->NULL;
    SRE 412 -->EXIT;
138
139 Close (Output_File);
    OVV 13 -->Close (Input_File);
    SRN 357 -->NULL;
    SRE 413 -->EXIT;
140
    END Matrix;
```

