# Escaping Saddle Points in Constant Dimensional Spaces: an Agent-based Modeling Perspective 

Grant Schoenebeck, University of Michigan
Fang-Yi Yu, Harvard University

## Results

- Analyze the convergence rate of a family of stochastic processes
- Three related applications
- Evolutionary game theory
- Dynamics on social networks
- Stochastic Gradient Descent



## Target Audience



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## Target Audience (still not-to-scale)



## Outline

- Escaping saddle point

Evolutionary game theory

Dynamics
on social
networks

Stochastic Gradient Descent

## Outline

- Escaping saddle point
- Case study: dynamics on social networks


Upper bounds and lower bounds

## ESCAPING SADDLE POINTS

## Reinforced random walk with $F$

A discrete time stochastic process $\left\{X_{k}: k=0,1, \ldots\right\}$ in $\mathbb{R}^{d}$ that admits the following representation,

$$
X_{k+1}-X_{k}=\frac{1}{n}\left(F\left(X_{k}\right)+U_{k}\right)
$$

$X_{k}$


## Reinforced random walk with $F$

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- Unbiased noise (noise), $U_{k}$



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- Unbiased noise (noise), $U_{k}$
- Step size, $1 / n$



## Examples

A discrete time Markov process $\left\{X_{k}: k=0,1, \ldots\right\}$ in $\mathbb{R}^{d}$ that admits the following representation,

$$
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$$

- Agent based models with $n$ agents
- Evolutionary games
- Dynamics on social networks
- Heuristic local search algorithms with uniform step size $1 / n$


## Node Dynamic on complete graphs [SY18]

- Let $f_{N D}:[0,1] \rightarrow[0,1]$. $n$ agents interact on a complete graph
- Each agent $v$ has an initial binary state $C_{0}(v) \in\{0,1\}$
- At round $k$,
- Pick a node $v$ uniformly at random
- Compute the fraction of opinion $1, X_{k}=\frac{\left|C_{k}^{-1}(1)\right|}{n}<$ Complete graph
- Update $C_{k+1}(v)$ to 1 w.p. $f_{N D}\left(X_{k}\right) ; 0$ o.w.


## Node Dynamic

Includes several existing dynamics

- Voter model
- Iterative majority [Mossel et al 14]
- Iterative 3-majority [Doerr et al 11]


## Node Dynamic

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## Reinforced random walk on $\mathbb{R}$

- $X_{k}$ be the fraction of nodes in state 1 at $k$.


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## Reinforced random walk on $\mathbb{R}$

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- Given $X_{k}$, the expected number of nodes in state 1 after round $k$, is $\mathrm{E}\left[n X_{k+1} \mid X_{k}\right]=n X_{k}+\left(f_{N D}\left(X_{k}\right)-X_{k}\right)$.


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Updated to 1 from 1

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Drift $F\left(X_{k}\right)$


## Node Dynamic

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## Reinforced random walk on $\mathbb{R}$

- $X_{k}$ be the fraction of nodes in state 1 at $k$.
- $X_{k+1}-X_{k}=\frac{1}{n}\left(\left(f_{N D}\left(X_{k}\right)-X_{k}\right)+U_{k}\right)$.

Drift Noise


## Question

Given $F$ and $U$, what is the limit of $X_{k}$ for sufficiently large $n$ ?

$$
X_{k+1}-X_{k}=\frac{1}{n}\left(F\left(X_{k}\right)+U_{k}\right)
$$



## Mean field approximation

$$
X_{k+1}-X_{k}=\frac{1}{n}\left(F\left(X_{k}\right)+U\left(X_{k}\right)\right)
$$



$$
x^{\prime}=F(x)
$$

(

## Mean field approximation

If $n$ is large enough, for $k=O(n), X_{k} \approx x\left(\frac{k}{n}\right)$ by Wormald et al 95 .

## Regular point

If $n$ is large enough, for $k=O(n), X_{k} \approx x\left(\frac{k}{n}\right)$.


## Fixed point, $\boldsymbol{F}\left(\boldsymbol{x}^{*}\right)=\mathbf{0}$

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## Escaping non-attracting fixed point

When can the process escape a non-attracting fixed point?


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1. $\Theta(n)$
2. $\Theta(n \log n)$
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## Escaping non-attracting fixed point

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## Lower bound

Escaping saddle point region takes at least $\Omega(n \log n)$ steps.


## Upper bound

Escaping saddle point region takes at most $\mathrm{O}(n \log n)$ steps. If


## Upper bound

Escaping saddle point region takes at most $\mathrm{O}(n \log n)$ steps.
If

- Noise, $U_{k}$
- Martingale difference
- bounded
- Noisy (covariance matrix is large)
- Expected difference, $F \in \mathcal{C}^{2}$
$-x^{*}$ is hyperbolic



## Gradient-like dynamics

Converges to an attracting fixed-point region in $\mathrm{O}(n \log n)$ steps.
If

- Noise, $U_{k}$
- Martingale difference
- bounded
- Noisy
- Expected difference, $F \in \mathcal{C}^{2}$
- Fixed points are hyperbolic
- Potential function



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Dynamics on social networks

## (DIS)AGREEMENT IN PLANTED COMMUNITY NETWORKS

## Echo chamber

## Beliefs are amplified through interactions in segregated systems



## Echo chamber

## Beliefs are amplified through interactions in segregated systems



## Echo chamber

## Beliefs are amplified through interactions in segregated systems




## Question

What is the consensus time given a rich-get-richer opinion formation and the level of intercommunity connectivity?


## Node Dynamic [Schoenebeck, Yu 18]

- Fixed a graph $G=(V, E)$ opinion set \{0,1\}
- Given an initial configuration

$$
X_{0}: V \mapsto\{0,1\}
$$

- At round $t$,
- A node $v$ is picked uniformly at random

The update of opinion only depends on the fraction of opinions amongst its neighbors


## Node Dynamic ND $\left(G, f_{N D}, X_{0}\right)$

- Fixed a (weighted) graph $G=(V, E)$ opinion set $\{0,1\}$, an update function $\boldsymbol{f}_{\boldsymbol{N D}}$
- Given an initial configuration $X_{0}: V \mapsto\{0,1\}$
- At roundt,
- A node $v$ is picked uniformly at random
- $X_{t}(v)=1$ w.p. $\boldsymbol{f}_{N D}\left(r_{X_{t-1}(v)}\right)$;

$$
=0 \text { otherwise }
$$

## Planted Community

- A weighted complete graph with n nodes, $K(n, p)$
- Two communities with equal size
- An edge has weight $p$ if in the same community and $1-p$ o.w.



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Complete graph

## Question

- What is the interaction between rich-get-richer opinion formation and the level of intercommunity connectivity?



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## Strong Community Structure

- There exists an initial state such that the process cannot reach consensus fast.

-3-Majority



## Weak Community Structure

- For all initial states, the process reaches consensus fast.




## Our Dichotomy Theorem

- Given a smooth rich-get-richer function $f_{N D} \in \mathcal{C}^{2}$, and a planted community graph $G=K(n, p)$. The maximum expected consensus time of $\operatorname{ND}\left(G, f_{N D}, X_{\mathbf{0}}\right)$ has two cases:


Complete graph
Two isolated complete graphs

## Node dynamic

- A Markov chain on 2-d grid
- $(0,0)$ and ( 1,1 ) are consensus states



## Our Dichotomy Theorem





## Our Dichotomy Theorem



## Our Dichotomy Theorem



## Our Dichotomy Theorem



Attracting point


## Fast consensus

## $\vec{X}_{k+1}-\vec{X}_{k}=\frac{1}{n}\left(F_{N D}\left(\vec{X}_{k}\right)+U\left(\vec{X}_{k}\right)\right)$ reach an attracting fixed point in $O(n \log n)$



## Question?



