
CS483 Analysis of Algorithms

Lecture 04 – Graph *

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*this lecture note is based on *Algorithms* by S. Dasgupta, C.H. Papadimitriou, and U.V. Vazirani and *Introduction to the Design and Analysis of Algorithms* by Anany Levitin.

▷ Introduction

What can be represented as graph?

What are the problems that can be solved using graphs?

Graph Representation

Explore graphs

Topological sort

Strong connected components

Conclusion

Introduction

What can be represented as graph?

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What are the problems that can be solved using graphs?

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Network routing

Graph Representation

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What can be represented as graph?

What are the problems that can be solved using graphs?

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Adjacency matrix

– Space

Adjacency list

– Space

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▷ Explore graphs

Graph Search

Graph Search

Graph Search

Depth-first search

Depth-first search

DFS Application

DFS Application

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DFS Application

Topological sort

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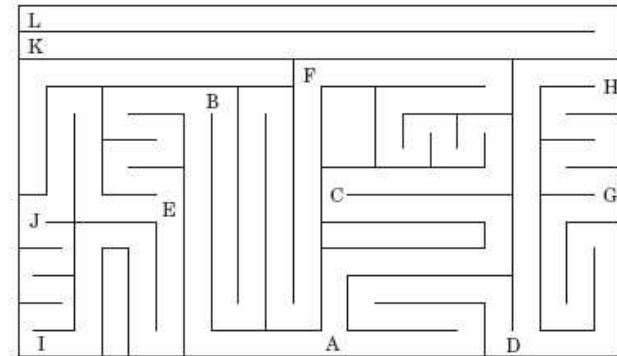
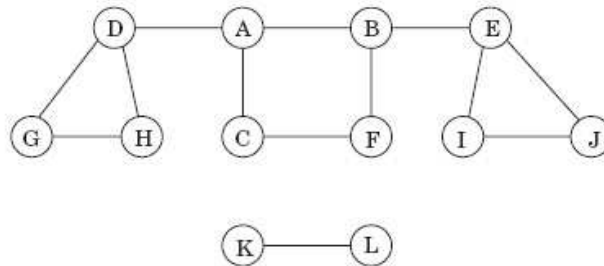
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Explore graphs

Graph Search

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 - Graph Search
 - Depth-first search
 - Depth-first search
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 - DFS Application
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- What parts of the graph are reachable from a given vertex? (i.e., connected components)
- Many problems require processing all graph vertices (and edges) in systematic fashion
- Basic tools to safely explore an unknown environment
-
-



Graph Search

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- Basic exploration algorithm

Algorithm 0.1: EXPLORE($G = \{V, E\}, v \in V$)

- Can the algorithm always work?
 - *proof*

Graph Search

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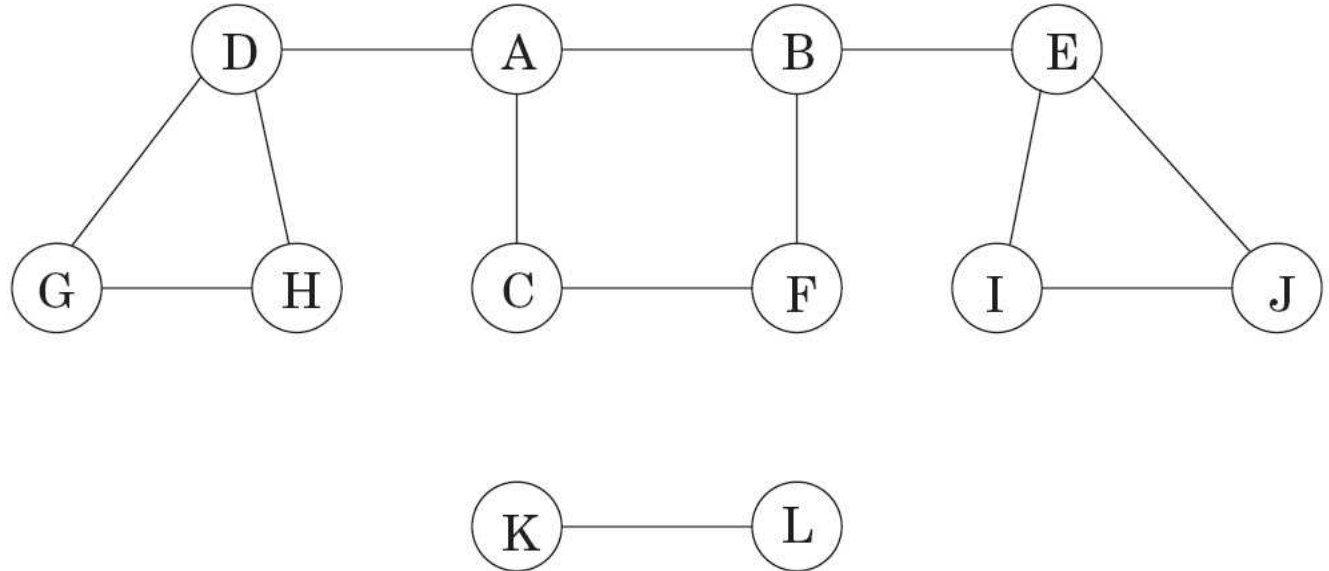
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□ Example: EXPLORE(B)



Depth-first search

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□ DFS

Algorithm 0.2: $\text{DFS}(G = \{V, E\}, v \in V)$

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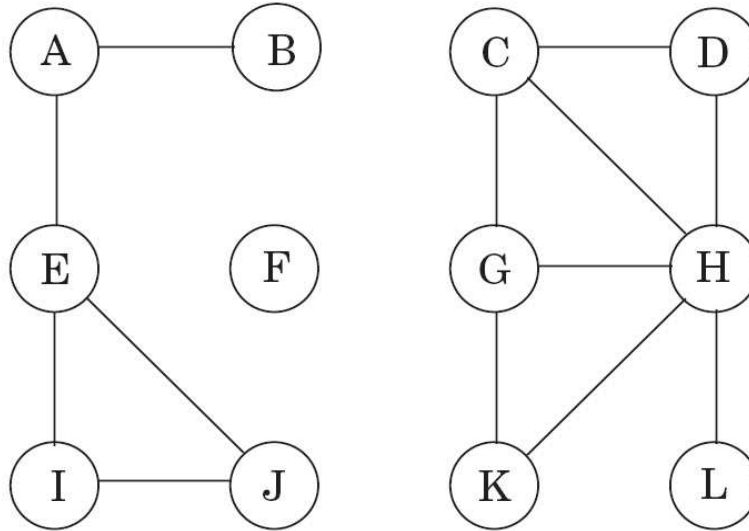
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□ Example:



□ Time complexity:

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□ Connect component

– Given a graph G , report the number of connect components in G .

– Given a graph G , can you preprocess G so that you can check if two nodes u and v from G are from the same connect component?

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- Ancestor/Descendant relationship of tree
 - Given a tree T , can you preprocess T so that you can answer if u is the ancestor of v in *constant time*, where u and v are two nodes from T .

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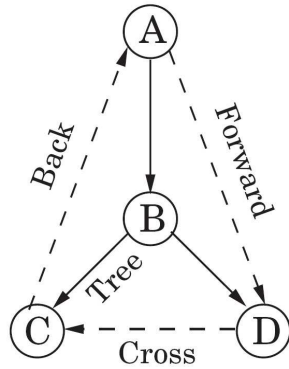
Conclusion

- Given a *directed* graph G , convert G to a tree whose nodes and edge are the vertices and edges of G

DFS Application

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- Types of edges
DFS tree



- To identify the type of an edge: pre/post ordering

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DAG and Topological Sort

Topological Sort: Using
DSF

Topological Sort: Using
Source Removal

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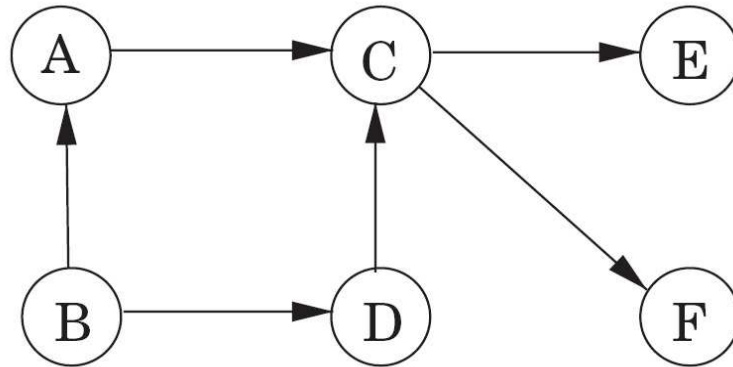
Topological Sort: Using
Source Removal

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Conclusion

- A graph G without (directed) cycle is a *directed acyclic graphs* (DAG)
- DAG can be found in modeling many problems that involve prerequisite constraints (construction projects, document version



control)

- Given a *directed* graph G , identify *cycles* in G

– *proof*

DAG and Topological Sort

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Sort

Topological Sort: Using
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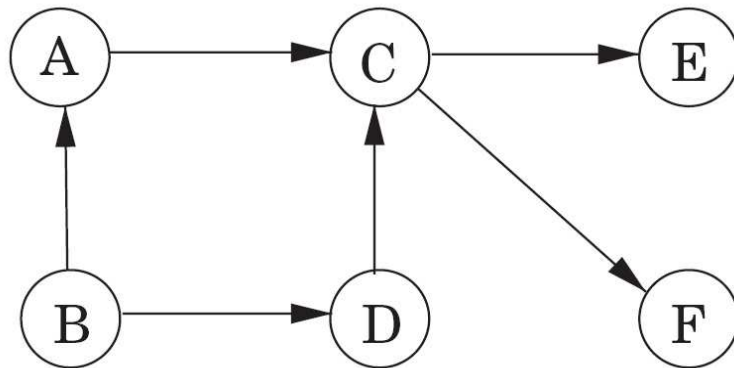
Topological Sort: Using
Source Removal

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Conclusion

- **Topological sorting or Linearization:** Vertices of a DAG can be linearly ordered so that:
 - Every edge its starting vertex is listed before its ending vertex
 - Being a DAG is also a necessary condition for topological sorting be possible
- **Example:**



Topological Sort: Using DSF

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- Compute DSF and reverse the visit order

Algorithm 0.3: $TS(G = \{V, E\})$

- Why does it work?
- Time complexity?

Topological Sort: Using Source Removal

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- Identify and remove sources iteratively.
 - A source is a vertex without incoming edges.

Algorithm 0.4: $TS(G = \{V, E\})$

- Why does it work?

- Time complexity?

Example

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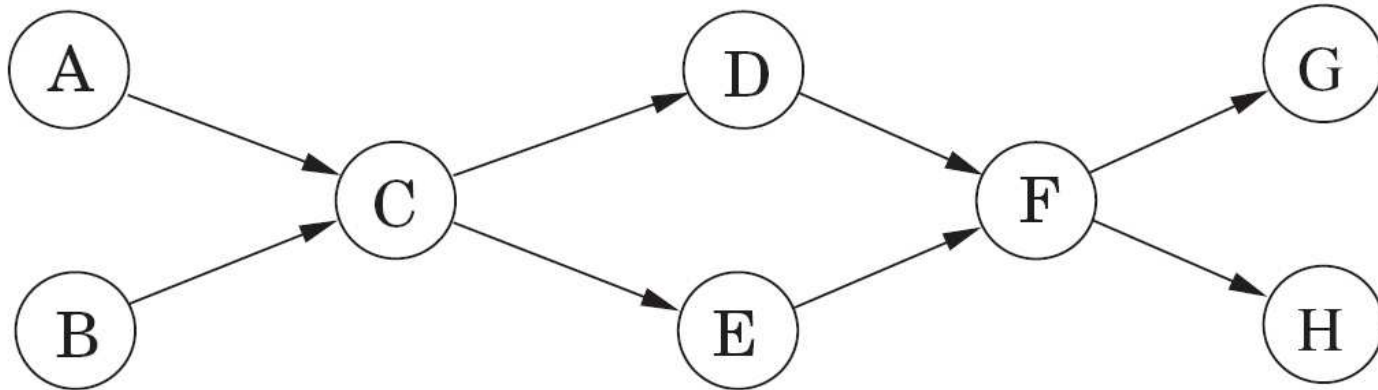
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□ Example:



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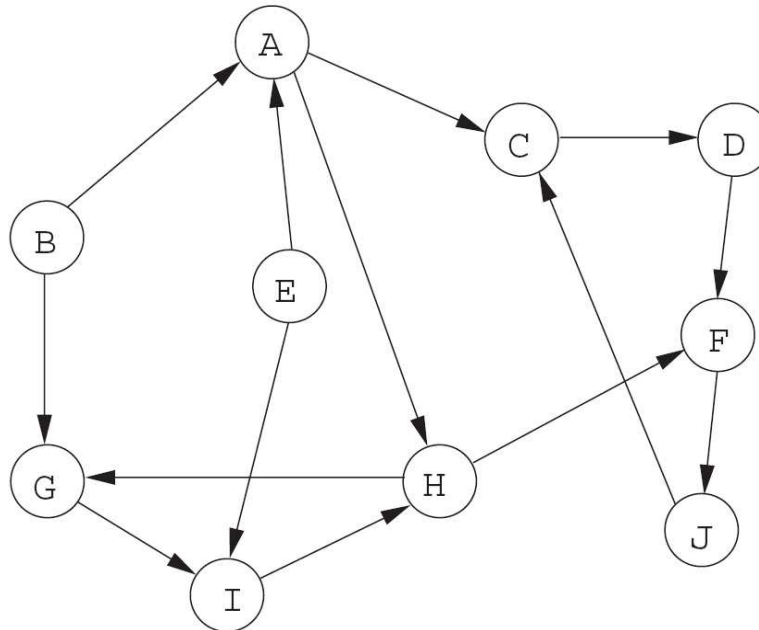
Strongly connected components and DAG

Strongly connected components and DAG

Strongly connected components and DAG

Conclusion

- **Definition:** Two nodes u and v are from the connected if and only if there is a path from u to v and a path from v to u .
- **Definition:** A set of vertices form a strongly connected component (SCC) iff any pairs of vertices are connected.



Strongly connected components

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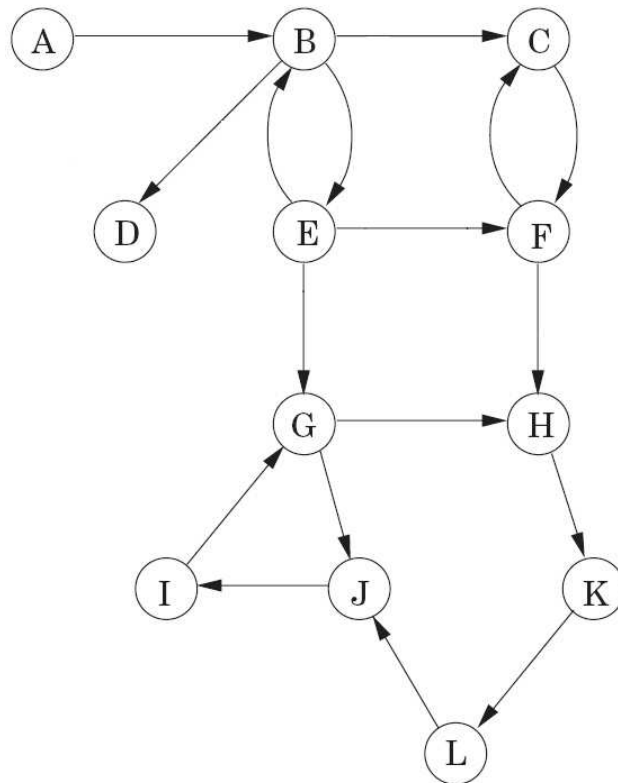
Strongly connected components and DAG

Strongly connected components and DAG

Strongly connected components and DAG

Conclusion

- Connected components in directed graph is less intuitive than that of undirected graph.
 - How many connected components are there in the graph below?



- How to compute SCCs from a directed graph?

Strongly connected components and DAG

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Conclusion

- Observation 1:**
- Observation 2:**

- Our strategy to find all SCC:**

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Strongly connected components and DAG

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Conclusion

- Task:** Find a vertex u in DAG sink node
- Fact:** It is easier to find a source node.

—

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▷ *proof:*

- How to find a vertex u in DAG sink node?

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Strongly connected components and DAG

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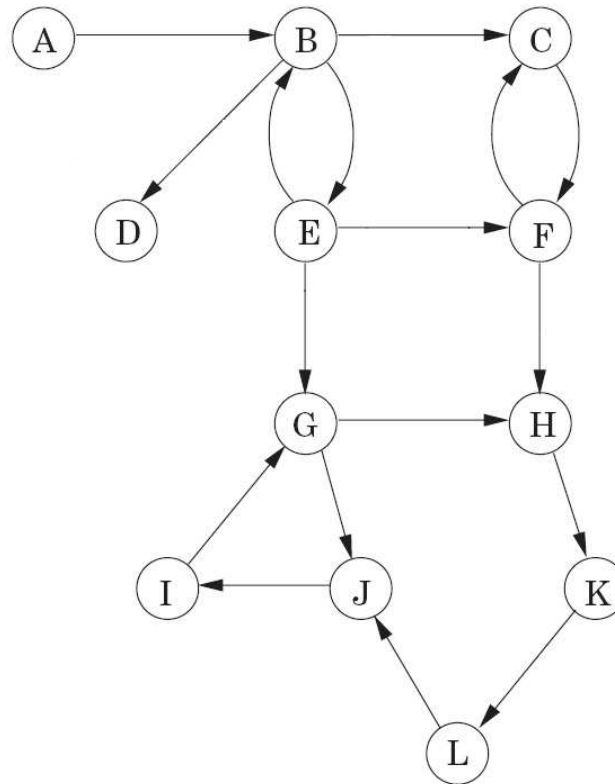
Strongly connected components and DAG

Conclusion

- **Task:** Remove all nodes from the previous SCC and identify a new sink node

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- **Example:**



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Summary

More graph algorithms

Conclusion

Summary

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▷ Summary

More graph algorithms

- Graphs can be very useful for many problems.
- DFS can be used for
 - Explore the graph
 - Reveal relationship between the graph nodes and types of edges
 - Linearization for DAG
 - Identify cycles, connected components, strongly connected components
- Assignment
 -

More graph algorithms

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Summary

▷ More graph algorithms

- Next week we will discuss problems related to paths in graph
 - shortest path in undirected and directed graphs
 - shortest path in weighted graphs
 - shortest path in graphs with negative edges
 - shortest path in DAG