# CS483 Analysis of Algorithms Lecture 08 – Dynamic Programming 02 \*

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<sup>\*</sup>this lecture note is based on *Algorithms* by S. Dasgupta, C.H. Papadimitriou, and U.V. Vazirani and *Introduction to the Design and Analysis of Algorithms* by Anany Levitin.

### **Edit Distance**

<ul> <li>Edit Distance</li> <li>Edit Distance</li> <li>Edit Distance</li> <li>Edit Distance and DAG</li> <li>Chain Matrix</li> <li>Multiplication</li> <li>Chain Matrix</li> <li>Multiplication</li> <li>Chain Matrix</li> <li>Multiplication</li> </ul>	Given two similar to Edit dist substitut S
Transitive closure Warshall's Algorithm Warshall's Algorithm Warshall's Algorithm All-pairs Shortest path problem Floyd's Algorithm Floyd's Algorithm Floyd's Algorithm	S How do – Brute
Travelling Salesman Problem (TSP) Travelling Salesman Problem (TSP) Conclusion Conclusion	– Gree

- Given two strings "lqorihten" and "algorithm", can you tell how similar these strings are?
- □ Edit distance is the number of operations (deletions, insertions, substitutions) that you can convert from one string to the other.

S		Ν	0	W	Y	—	S	Ν	0	W	_	Y
S	U	Ν	Ν	_	Y	S	U	Ν	_	_	Ν	Y
		Cos	st: 3					(	Cost:	5		

How do you compute the smallest edit distance between two strings?

– Brute force method? What's the time complexity?

– Greedy algorithm?

– Dynamic programming

### **Edit Distance**

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**Edit Distance** Edit Distance Edit Distance Edit Distance and DAG Chain Matrix Multiplication Chain Matrix Multiplication Chain Matrix Multiplication Transitive closure Warshall's Algorithm Warshall's Algorithm Warshall's Algorithm All-pairs Shortest path problem Floyd's Algorithm Floyd's Algorithm Floyd's Algorithm Travelling Salesman Problem (TSP) **Travelling Salesman** Problem (TSP) Conclusion Conclusion

**Observation**: Given two strings  $x[1 \cdots n]$  and  $y[1 \cdots m]$ . No matter how we match x to y, at the end of the match, we can only have:

- Question: Is it possible x[n] matches to y[i < m]? or y[m] matches to x[j < n]?
- □ **Example**: EXPONENTIAL vs. POLYNOMIAL
  - What are possible endings?
  - What are the subproblems we should consider?
  - How do we get an optimal answer?

### **Edit Distance**

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 $\Box$  Let E(i, j) be the edit distance for the subproblem of strings with lengths *i* and *j* 

 $\Box \quad E(i,j) = \min\{E(i-1,j-1) + \operatorname{diff}(x[i],y[j]), E(i-1,j) + 1, E(i,j-1) + 1\}$ 



## **Edit Distance and DAG**



Edit Distance Edit Distance Edit Distance Edit Distance and DAG Chain Matrix ▷ Multiplication Chain Matrix Multiplication Chain Matrix Multiplication Transitive closure Warshall's Algorithm Warshall's Algorithm Warshall's Algorithm All-pairs Shortest path problem Floyd's Algorithm Floyd's Algorithm Floyd's Algorithm Travelling Salesman Problem (TSP) Travelling Salesman Problem (TSP) Conclusion Conclusion

Given four matrices,  $A[50 \times 20]$ ,  $B[20 \times 1]$ ,  $C[1 \times 10]$ ,  $D[10 \times 100]$ , we wish to compute  $A \times B \times C \times D$ .

 $\Box \quad \text{If we compute } (((A \times B) \times C) \times D), \text{ we will perform } x \\ \text{multiplications?}$ 

 $\square \quad \text{What about } ((A \times B) \times (C \times D))?$ 

□ How do we find the best way to group matrices so that the number of multiplications is minimized?

– Brute force

- Greedy algorithm

– Dynamic programming

Edit Distance
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Edit Distance and DAG
Chain Matrix
Multiplication
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➢ Multiplication
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Multiplication
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Warshall's Algorithm
Warshall's Algorithm
All-pairs Shortest path
problem
Floyd's Algorithm
Floyd's Algorithm
Floyd's Algorithm
Travelling Salesman
Problem (TSP)
Travelling Salesman
Problem (TSP)
Conclusion
Conclusion

#### □ Dynamic programming and DAG

- 1. A pair of parentheses groups two matrices
- 2. The final matrix represents the root
- 3. Example:  $(((A \times B) \times C) \times D)$  and  $((A \times B) \times (C \times D))$

- $\Box$  So, our goal is to build an optimal binary tree
- Given four matrices,  $A[50 \times 20]$ ,  $B[20 \times 1]$ ,  $C[1 \times 10]$ ,  $D[10 \times 100]$ , we wish to compute  $A \times B \times C \times D$ .
  - 1. Subproblems with two matrices:
  - 2. Subproblems with three matrices:
  - 3. Subproblems with four matrices:

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General cases: Give a list of matrices  $\{A_i\}$ , C(i, j) be the minimum cost of  $A_i \times \cdots A_j$ , then

$$C(i,j) = \min_{i \le k < j} \{ C(i,k) + C(k+1,j) + m_{i-1} \cdot m_k \cdot m_j \}$$

 $\Box$  Example:  $A[50 \times 20], B[20 \times 1], C[1 \times 10], D[10 \times 100]$ 

0	j = 1	j = 2	j = 3	j = 4	j = 5
i = 1					
i=2					
i=3					
i=4					
i = 5					

## **Transitive closure**

**Edit Distance** Edit Distance **Edit Distance** Edit Distance and DAG Chain Matrix Multiplication Chain Matrix Multiplication Chain Matrix Multiplication  $\triangleright$  Transitive closure Warshall's Algorithm Warshall's Algorithm Warshall's Algorithm All-pairs Shortest path problem Floyd's Algorithm Floyd's Algorithm Floyd's Algorithm Travelling Salesman Problem (TSP) Travelling Salesman Problem (TSP)

Conclusion

Conclusion

- □ Transitive closure of a graph is a set of vertex pairs of a graph, which can be connected by one or multiple paths
- $\square$  We can represent the transitive closure using a matrix A. The element  $A_{i,j}$  is "1" if there are one or multiple paths between vertices i and j.

 $\Box$  Example:



□ Can you design a brute force algorithm? What is its time complexity?

## Warshall's Algorithm

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 $\Box$  Important observation: If there is a path from *a* to *z* via *s* then there must be a path from *a* to *s* and from *s* to *z* 

 $\Box$  Let  $A^k$  be the optimal answer when we only allow the first k nodes to be intermediate nodes in paths. We can compute the optimal solution for k + 1 nodes  $A^{k+1}$  efficiently

 $\Box$  What is  $A^0$ ?

 $\Box$  For k > 0,

$$A^{k+1}[i,j] = \begin{cases} 1 & (A^k[i,j]=1) \\ A^k[i,k] \text{ and } A^k[k,j] & (\text{otherwise}) \end{cases}$$

## Warshall's Algorithm

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 $\Box$  Example:



via Ø	A	B	C	D	E	via A	A	B	C	D	E
A						A					
B						В					
C						C					
D						D					
E						E					
via B	A	B	C	D	E	via C	$A$	B		D	
						A					
B											
C											
D											
E											
via D	$A$	B		D		via $E$	A	B	C	D	
A											
B						B					
C											
D						D					
E											

#### Warshall's Algorithm

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#### Algorithm

Algorithm 0.1: WARSHALL( $A[1 \cdots n]$ )

$$\begin{aligned} & \text{for } i \leftarrow \{1 \cdots n\} \\ & \text{do } \begin{cases} & \text{for } j \leftarrow \{1 \cdots n\} \\ & \text{do } \begin{cases} & \text{for } k \leftarrow \{1 \cdots n\} \\ & \text{do } A^k[i,j] \leftarrow (A^{k-1}[i,k] \text{ and } A^{k-1}[k,j]) \text{ or } A^{k-1}[i,j] \end{cases} \end{aligned}$$

 $\Box$  Time complexity?

Edit Distance Edit Distance Edit Distance Edit Distance and DAG Chain Matrix Multiplication Chain Matrix Multiplication Chain Matrix Multiplication Transitive closure Warshall's Algorithm Warshall's Algorithm Warshall's Algorithm All-pairs Shortest path  $\triangleright$  problem Floyd's Algorithm Floyd's Algorithm Floyd's Algorithm Travelling Salesman Problem (TSP) Travelling Salesman Problem (TSP) Conclusion Conclusion

□ In this problem we want to find the shortest paths connecting all possible pairs of vertices of a graph

 $\Box$  Example:



 $\Box$  What is the brute force algorithm and its time complexity?

# **Floyd's Algorithm**

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A.k.a. Floyd-Warshall algorithm (or the Roy-Floyd algorithm)
 Robert Floyd (1936-2001)



(Robert Floyd, 1972, from http://sigact.acm.org/floyd/)
 The algorithm is very similar to Warshall's algorithm
 Basic idea: Let A<sup>k-1</sup> be the optimal answer when we only allow the first k - 1 nodes to be intermediate nodes in paths. We can compute the optimal solution for k nodes A<sup>k</sup> efficiently



# **Floyd's Algorithm**

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 $\Box$  Example:



via Ø	A	B	C	D	E	via A	A	B	C	D	E
A						A					
B						В					
C						C					
D						D					
E						E					
via B	$A$	B	C	D		via C	A	B	C	D	E
A						A					
B						B					
D						D					
via D	A	B	C	D	E	via E	A	B		D	
A											
D						D					

# **Floyd's Algorithm**

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**Edit Distance** Edit Distance **Edit Distance** Edit Distance and DAG Chain Matrix Multiplication Chain Matrix Multiplication Chain Matrix Multiplication Transitive closure Warshall's Algorithm Warshall's Algorithm Warshall's Algorithm All-pairs Shortest path problem Floyd's Algorithm Floyd's Algorithm ▷ Floyd's Algorithm Travelling Salesman Problem (TSP) Travelling Salesman Problem (TSP)

Conclusion

Conclusion

#### Algorithm

**Algorithm 0.2:**  $FLOYD(A[1 \cdots n])$ 

$$\begin{aligned} & \text{for } i \leftarrow \{1 \cdots n\} \\ & \text{do } \begin{cases} & \text{for } j \leftarrow \{1 \cdots n\} \\ & \text{do } \begin{cases} & \text{for } k \leftarrow \{1 \cdots n\} \\ & \text{do } A^k[i,j] \leftarrow \min\{(A^{k-1}[i,k] + A^{k-1}[k,j]), A^{k-1}[i,j]\} \end{cases} \end{aligned}$$

 $\Box \quad \text{Time complexity}?$ 

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 $\square$  **Problem**: Find the shortest path from A to A by visiting each vertex exactly once



 $\Box$  Brute force:

 $\Box$  Greedy:

□ Dynamic programming:

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Conclusion

Conclusion

□ Given a graph with n nodes and starting vertex is 1.
 □ Algorithm

Algorithm 0.3:  $FLOYD(A[1 \cdots n])$ 

 $\begin{aligned} & \text{for } s \leftarrow \{2 \cdots n\} \\ & \text{do } \begin{cases} \text{for all subsets } S \subset \{1, 2, \cdots, n\} \text{ of size } s \text{ and containing } 1 \\ & \text{do } \begin{cases} \text{for } j \in S \text{ and } j \neq 1 \\ & \text{do } C(S, j) = \min_{i \in S, i \neq j} \{C(S - \{j\}, i) + d_{ij}\} \end{cases} \end{aligned}$ 

☐ Time complexity?

#### Conclusion



**Travelling Salesman Problem** (from Randall Munroe, creator of xkcd)

... EXACTLY? UHH ...

LISTEN, I HAVE SIX OTHER

O(1)

#### Conclusion

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 $\Box$  We have solved the following problems using dynamic programming

- Longest increasing sequence
- Binomial coefficients of  $(a + b)^n$  (Pascal's triangle)
- Knapsack problem
- Edit distance
- Matrix multiplication chain (optimal binary tree)
- Transitive closure (Warshall's algorithm)
- All pairs shortest paths (Floyd's algorithm)
- TSP
- □ It is usually more difficult to represent a problem as a set of sub-problems
- □ Next couple of weeks: Linear programming.