# CS311 Data Structures <br> Lecture 11 - Red-Black Trees (Deletion) 

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## Red-Black Tree Review

## Insertion and Invariants

1. Color swap between grandparent and (parent and uncle) does not change black height
2. Rotation at grandparent followed by recoloring does not change black height


## Red-Black Tree Deletion

## number of kids

Let $x$ be the node in RB tree to be deleted

1. $x$ has two children nodes

- swap value with the node $y$ with the largest value in the left subtree
- then delete node $y$
- node $y$ must have one or zero child node

2. $x$ has one child
3. $x$ has zero child

## color of node

Let $x$ be the node in RB tree to be deleted

1. $x$ is red, delete as usual

- black depth remains the same
- no two adjacent nodes are red

2. $x$ is black, many hairy cases

- black depth along the path containing $x$ is one less
- two adjacent nodes may be red


## Case Analysis

## $x$ is black

1. deleting $x$ is going to reducing the black height of the path containing $x$ by one
2. key idea: recoloring $x$ 's sibling $s$ to red, and reducing the black height of the path containing $s$ by one

- Question: is it possible that $s$ is null?
- Question: what to do if $s$ is already red?


## Some simple practice

Questions

1. How to get your sibling?
2. How to get your uncle?

## More Case Analysis

Now we reduced our analysis to: $x$ is black and $x$ 's kid $k$ is also black (or null). We further let $s$ be the sibling, $L$ and $R$ are the left and right children of $s$, respectively, and let $P$ be the parent.

## Five possible cases

1. $s$ is red
1.1 What are the colors of $P, L, R$ ?
2. $s$ is black
2.1 Both kids $L$, and $R$ are all black
2.2 Both $R$ and $L$ are red
$2.3 R$ is red, $L$ is black
2.4 $L$ is red, $R$ is black
