Network Programming using sockets

Distributed Software Systems
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Outline

- 1. Overview
 - (a) Motivation
 - (b) Terminology and Concepts
- 2. The socket interface
- 3. Client software design
- 4. Server software design
- 5. Examples (TFTP,HTTP)

Overview

- Common communication patterns in a distributed applications
 - Client-server
 - Group (multicast) communication
- Client: process that requests service
- Server: process that provides service
- client usually blocks until server responds
- client usually invoked by end users when they require service
- server usually waits for incoming requests
- server can have many clients making concurrent requests
- server often a program with special system privileges

Issues in Client design

- Must know or must find out the location of the server
- Which protocol to use: reliable or unreliable?
- blocking (synchronous) request or non-blocking (asynchronous)

Issues in Server Design

- Concurrent vs iterative servers: handle multiple requests concurrently or one after the other?
- Connection-oriented vs connection-less servers:
 TCP or UDP?
- Stateful vs stateless servers
- Multi-protocol servers
- Multi-service servers

Connection-less vs connection-oriented servers

- protocol used determines level of reliability
- TCP is reliable protocol:
 - verifies that data arrives at the other end, retransmits segments that don't
 - checks that data is not corrupted on the way
 - makes sure data arrives in order
 - eliminates duplicate packets
 - provides flow control to make sure sender does not send data faster than the receiver can consume it
 - informs both client and server if underlying network becomes inoperable

Connection-less vs connection-oriented servers

- UDP unreliable does not take any guarantees about reliable delivery
- UDP relies on application to take whatever actions are necessary for reliability
- UDP used if
 - application protocol designed to handle reliability and delivery errors
 - communication occurring over a LAN with hardware broadcast or multicast
 - overhead of TCP connections too much for application

Stateful vs stateless servers

- Information that server maintains about the status of ongoing interactions with clients \equiv *state* information
- stateful servers
 - can reduce size of messages from clients
 - state information can help server in performing request faster
- stateless servers
 - quicker and more reliable recovery after crashes
 - no problems because of lost, delayed, duplicate messages etc.
 - smaller memory requirements
- Example: stateful vs stateless file service
- stateless servers: application protocol should have idempotent operations

Concurrency in servers

- Concurrency needed if several clients and service is expensive
- Operating system support
 - Multiple processes
 - Threads
 - Asynchronous I/O: select system call

APIs for TCP/IP

- TCP/IP is a protocol designed to operate in multivendor environment
- interface between TCP/IP and applications loosely specified
- application interfaces
 - BSD UNIX: socket interface
 - AT&T: TLI interface
- TCP/IP software inside kernel invoked by system calls
- UNIX I/O facilities extended with TCP/IP specific calls

The Socket Interface

- provides functions that support network communication using many possible protocols
 - PF_INET is one protocol family supported by sockets
 - TCP and UDP are protocols in PF_INET family
- socket is the abstraction for network communication
- a socket is identified by socket descriptor
- system data structure for socket
 - family (e.g., PF_INET)
 - service (e.g., SOCK_STREAM)
 - Local IP address, Local Port
 - Remote IP address, Remote Port
- passive socket: socket used by a server to wait for incoming connections; active socket: socket used by client to initiate a connection

Endpoint Addresses

- TCP/IP protocols define a communication endpoint to consist of an IP address and a protocol port number
- other protocol families have other definitions
- socket abstractions supports the concept of address family which allows different protocols to have their own address representations
- TCP/IP protocols use a single address representation with address family denoted by AF_INET

Endpoint Addresses cont'd

structure for AF_INET addresses

```
*
                                                                                    *
/* struct to hold an address
                                                     /* protocol port number
/* IP address
                                                                                                         /* unused (set to zero)
                                        /* type of address
                    /* total length
                                                                                     struct in_addr sin_addr;
                                         u_short sin_family;
struct sockaddr_in {
                                                              u_short sin_port;
                                                                                                          char sin_zero[8];
                     u_char sin_len;
```

 if program using mixture of protocols, programmer must be careful since not all addresses have the same size

System Calls

socket

- used to create new socket
- arguments: protocol family (e.g. PF_INET),
 protocol or service (i.e., stream or datagram)
- returns socket descriptor

• connect:

- client calls connect to establish an active connection to the server
- argument to connect specifies remote endpoint

write

- servers and clients use write to send data across a TCP connection
- arguments: socket descriptor, address of data, length of data

System Calls cont'd

read

- used to receive data from a TCP connection
- arguments: socket, buffer, length of buffer
- read blocks if no data; if more data than fits in buffer, it only extracts enough to fill the buffer; if less than buffer length, it extracts all the data and returns number of bytes read
- read and write can also be used with UDP but different behavior
- close: used to deallocate socket; deleted when last process that is using socket does a close

bind

- used to specify a local endpoint address for a socket
- uses sockaddr_in structure

System Calls cont'd

listen

- used by connection-oriented servers to put socket in passive mode
- arguments: socket, size of queue for socket connection requests

accept

- creates a new socket for each connection request
- returns descriptor of new socket to its calller

• UDP calls:

- send, sendto, sendmsg
- recv, recvfrom, recvmsg

Integer Conversion

- standard representation for binary integers used in TCP/IP protocol headers: network byte order, MSB first
- e.g. the protocol port field in struct sockaddr_in uses network byte order
- host's integer representation maybe different
- conversion routines: htons, htonl, ntohl, ntohls should be used for portability

Client Software

conceptually simpler than servers because

- do not have to handle concurrent interactions with multiple servers
- usually not privileged software ⇒ don't have to be as careful
- no authentication, protection, etc.

Locating the server

server's IP address and port number needed

- can be specified as a constant in the program
- have the user specify it as an argument when invoking client
- read from a file on disk
- use a protocol to find the server (e.g. a broadcast message to which servers respond)

Parsing address argument

- address argument typically is a hostname like cs.gmu.edu or IP address in dotted decimal notation like 129.174.29.34
- need to specify address using structure sockaddr_in
- library routines inet_addr and gethostbyname used for conversions

```
struct hostent {
    char *hname;
    char **h_aliases;
    int h_addrtype;
    int h_length;
    char **h_addr_list;
};
#define h_addr h_addr_list[0];
```

EXAMPLE:

```
struct hostent *hptr;
char *name = ''cs.gmu.edu'';

if ( hptr = gethostbyname(name)) {
    /* IP address is in hptr->h_addr */
} else {
    /* handle error */
}
```

inet_addr converts dotted decimal IP address into binary

Client Software cont'd

- looking up a well known port by name
- struct servent defined in *netdb.h* in the same way as struct hostent

```
struct servent *sptr;

if (sptr = getservbyname(''smtp'',''tcp'')){
    /* port number is now in sptr->s_port */
} else {
    /* handle error */
}
```

 NOTE: getservbyname returns protocol port in network byte order

Client Software cont'd

- looking up a protocol by name
- struct protoent defined in *netdb.h*

```
struct protoent *pptr;

if (pptr = getprotobyname(''udp'')){
   /* official protocol number is in pptr->p_proto */
} else {
   /* handle error */
}
```

TCP client algorithm

- 1. Find IP address and protocol number of server
- 2. allocate a socket
- 3. specify that the connection needs an arbitrary, unused protocol port on local machine and allow TCP to select one
- 4. Connect the socket to the server
- 5. Communicate with the server using application-level protocol
- 6. Close the connection

TCP client cont'd

Allocating a socket

```
#include <sys/types.h>
#include <sys/socket.h>

int s; /* socket descripto */

s = socket(PF_INET,SOCK_STREAM, 0);
```

- Choosing a local port number
 - conflicts have to be avoided
 - happens as a side-effect to connect call
- choosing a local IP address
 - a problem for hosts connected to multiple networks
 - chosen automatically by TCP/IP at time of connection

Connecting a TCP socket to a server

retcode = connect(s,remaddr,remaddrlen)

- connect performs four tasks
 - 1. tests specified socket is valid and not already connected
 - 2. fills in remote address in socket from second argument
 - 3. chooses a local endpoint address for socket (if it does not have one)
 - 4. initiates a connection and returns value to the caller

Communicating with the server using TCP: Example

```
#define BLEN 120
char *req = ''request of some sort'';
char buf[BLEN];
char *bptr;
int n;
int buflen;
bptr = buf;
buflen = BLEN;
/* send request */
write(s,req,strlen(req);
/* read response (may come in several pieces) */
while ((n = read(s,bptr,buflen) > 0) {
   bptr += n;
  buflen -= n;
}
```

Closing a TCP connection

- partial close needed because client may not know when all the data from the server has arrived and server may not know if client will send another request
- shutdown call

```
errcode = shutdown(s,direction);
```

• direction = 0: no further input, 1: no further output, 2: shutdown in both directions

Programming a UDP client

- 1. Find IP address and protocol number of server
- 2. Allocate a socket
- 3. Specify that the connection needs an arbitrary, unused protocol port on local machine and allow UDP to select one
- 4. Specify the server to which messages must be sent
- 5. Communicate with the server using application-level protocol
- 6. Close the socket

Connected and Unconnected UDP sockets

- with UDP, connected sockets do not mean a "connection" was established
- onnected sockets ⇒ server specified once
- unconnected sockets ⇒ server specified each time
- read and write: message transfer NOT streams
- close does not inform remote endpoint of any actions
- UDP is unreliable

Examples

- TCP and UDP clients for services
 - DAYTIME
 - TIME
 - ECHO
- connectTCP and connectUDP procedures invoke connectsock

Issues in Server Design

- Concurrent vs iterative servers: handle multiple requests concurrently or one after the other?
- Connection-oriented vs connection-less servers:
 TCP or UDP?
- Stateful vs stateless servers

Iterative, connection-oriented server

Algorithm

- 1. Create a socket and bind to the well-known address for the service being offered
- 2. Place the socket in passive mode
- 3. Accept the next connection request from the socket, and obtain a new socket for the connection
- 4. Repeatedly read a request from the client, formulate a response, and send a reply back to the client according to the application protocol
- 5. When finished with a particular client, close the connection and return to step 3 to accept a new connection
- servers should specify INADDR_ANY as internet address while binding
- needed for hosts with multiple IP addresses

Iterative, connection-less servers

Algorithm

- 1. Create a socket and bind to the well-known address for the service being offered
- 2. Repeatedly read the next request from a client, formulate a response, and send a reply back to the client according to the application protocol
- cannot use connect (unlike clients)
- use sendto and recyfrom

Concurrent, Connection-less servers

Algorithm

- **Master 1.** Create a socket and bind to the well-known address for the service being offered. Leave the socket unconnected.
- **Master 2.** Repeatedly call recvfrom to receive the next request from a client, and create a new slave thread/process to handle the response
- **Slave 1.** Receive a specific request upon creation as well as access to the socket
- **Slave 2.** Form a reply according to the application protocol and send it back to the client using sendto

Slave 3. Exit

- cost of process/thread creation for each client request
- while using threads, use thread-safe functions and be careful while passing arguments to threads

Concurrent, Connection-oriented servers

Algorithm

- **Master 1.** Create a socket and bind to the well-known address for the service being offered. Leave the socket unconnected.
- Master 2. Place the socket in passive mode.
- **Master 3.** Repeatedly call accept to receive the next request from a client, and create a new slave process/thread to handle the response
- **Slave 1.** Receive a connection request (i.e., socket for connection) upon creation
- **Slave 2.** Interact with the client using the connection: read request(s) and send back response(s)
- **Slave 3.** Close the connection and exit
- processes created using fork; can also use execve

Apparent concurrency using a single process

- multiple processes ⇒ need to use shared memory IPC facilities if data structures shared among processes
- creating processes can be expensive
- threads make this easier
- can also achieve the same goal using a single process and asynchronous I/O using select

Apparent concurrency using a single process

Algorithm

- 1. Create a socket and bind to the well-known port for the service. Add the socket to the list of those on which I/O is possible
- 2. Use select to wait for I/O on existing sockets
- 3. If original socket is ready, use accept to obtain the next connection, and add the new socket to the list of those on which I/O is possible
- 4. If some socket other than the original is ready, use read to obtain the next request, form a response, and use write to send the response back to the client
- 5. Continue processing with step 2.

The Problem of Server Deadlock

- iterative server: suppose client creates a connection but does not send any requests
- suppose client does not consume responses
- ullet connection-oriented servers will block on write if local buffer full \Rightarrow deadlock in single process servers

Multi-protocol Server Design

- multiprotocol server handles service requests over both UDP and TCP
- Motivation: allows the use of shared code for service
- asynchronous I/O needed (select system call)
- design can be iterative or concurrent (multi-process or single-process)

Multi-service Server Design

- single server for multiple services
- Motivation: conserve system resources and make maintenance easier
- Design: Iterative, concurrent, or single process concurrent
- Connection-less or Connection-oriented
- Multi-service, Multi-protocol "super servers", e.g. UNIX inetd
- Static or dynamic server configuration

UNIX inetd super server

- configuration file /etc/inetd.conf
- entries: service name (from /etc/services), socket type, protocol, wait status, userid, server program, arguments

Management of Server Concurrency

- Iterative vs Concurrent: tradeoff depends on time to service request, time to create processes, time between requests
- Process/thread Preallocation for improving performance
 - especially useful if service involves I/O
 - In UNIX, slave processes can take advantage of socket sharing
 - NFS uses preallocation
 - Preallocation on multiprocessors
- Delayed Process/thread Allocation for improving performance
 - wait until sure that servicing request will take a long time before creating a new process/thread for handling request
- can combine delayed allocation with preallocation!

Java sockets API

- TCP socket classes
 - Socket
 - ServerSocket
 - InetAddress
- UDP classes
 - DatagramPacket
 - DatagramSocket

Java Examples

A TCP Client for the Echo service

```
import java.io.*;
import java.net.*;
public class EchoClient {
  public static void main(String[] args) throws IOException {
     Socket echoSocket = null:
     PrintWriter out = null;
     BufferedReader in = null:
     try {
        echoSocket = new Socket("taranis", 7);
        out = new PrintWriter(echoSocket.getOutputStream(), true);
        in = new BufferedReader(new InputStreamReader(
                           echoSocket.getInputStream()));
     } catch (UnknownHostException e) {
        System.err.println("Don't know about host: taranis.");
        System.exit(1);
     } catch (IOException e) {
        System.err.println("Couldn't get I/O for "
                     + "the connection to: taranis.");
        System.exit(1);
     }
     BufferedReader stdln = new BufferedReader(
                       new InputStreamReader(System.in));
     String userInput;
     while ((userInput = stdIn.readLine()) != null) {
        out.println(userInput);
        System.out.println("echo: " + in.readLine());
     out.close();
     in.close();
     stdln.close();
     echoSocket.close();
  }
}
```

A TCP Client for the Daytime service

```
import java.net.*;
import java.io.*;
public class DayClient1 {
 public static final int DAYTIME_PORT = 13;
 String host;
 Socket s:
 public static void main(String args[]) throws
IOException {
  DayClient1 that = new DayClient1(args[0]);
  that.go();
 public DayClient1(String host) {
  this.host = host:
 public void go() throws IOException {
  s = new Socket(host, DAYTIME_PORT);
  BufferedReader i = new BufferedReader(
    new InputStreamReader(s.getInputStream()));
  System.out.println(i.readLine());
  i.close();
  s.close();
```

A TCP Server for the Daytime service

```
import java.io.*;
import java.net.*;
import java.util.*;
public class DayServer1 {
 private ServerSocket ss;
 public static final int DAYTIME_PORT = 13;
 public static void main(String args[]) throws
IOException {
  DayServer1 d = new DayServer1();
  d.go();
 public void go() throws IOException {
  Socket s = null;
  ss = new ServerSocket(DAYTIME_PORT, 5);
  for (;;) {
   s = ss.accept();
   BufferedWriter out = new BufferedWriter(
      new OutputStreamWriter(s.getOutputStream(),"8859_1"));
   out.write("Java Daytime server: " +
             (new Date()).toString() + "\n");
   out.close();
   s.close();
```

A Multithreaded TCP server

```
public class MultiServe implements Runnable {
 private ServerSocket ss;
 public static void main(String args[]) throws Exception {
  MultiServe m = new MultiServe();
  m.go();
 }
 public void go() throws Exception {
  ss = new ServerSocket(DayClient2.DAYTIME_PORT, 5);
  Thread t1 = new Thread(this, "1");
  Thread t2 = new Thread(this, "2");
  Thread t3 = new Thread(this, "3");
  t1.start(); t2.start(); t3.start();
 }
 public void run() {
  Socket s = null:
  BufferedWriter out = null;
  String myname = Thread.currentThread().getName();
  for (;;) {
   try {
     System.out.println("thread " + myname + " about to accept..");
     s = ss.accept();
     System.out.println("thread " + myname +
                         "accepted a connection");
     out = new BufferedWriter(
          new OutputStreamWriter(s.getOutputStream()));
     out.write(myname + " " + new Date());
     Thread.sleep(10000);
     out.write("\n");
     out.close();
    catch (Exception e) {
     e.printStackTrace();
  }
 }
```

Another Multi-threaded Server Example

```
import java.net.*;
import java.io.*;
public class KKMultiServer {
  public static void main(String[] args) throws IOException {
     ServerSocket serverSocket = null;
     boolean listening = true;
     try {
        serverSocket = new ServerSocket(4444);
     } catch (IOException e) {
        System.err.println("Could not listen on port: 4444.");
       System.exit(-1);
     while (listening)
        new KKMultiServerThread( serverSocket.accept()).start();
     serverSocket.close();
  }
}
```

```
import java.net.*;
import java.io.*;
public class KKMultiServerThread extends Thread {
  private Socket socket = null;
  public KKMultiServerThread(Socket socket) {
     super("KKMultiServerThread");
     this.socket = socket;
  }
  public void run() {
     try {
        PrintWriter out = new PrintWriter(socket.getOutputStream(), true);
        BufferedReader in = new BufferedReader(
                        new InputStreamReader(
                        socket.getInputStream()));
        String inputLine, outputLine;
        KnockKnockProtocol kkp = new KnockKnockProtocol();
        outputLine = kkp.processInput(null);
        out.println(outputLine);
        while ((inputLine = in.readLine()) != null) {
           outputLine = kkp.processInput(inputLine);
           out.println(outputLine);
           if (outputLine.equals("Bye"))
             break;
        out.close();
        in.close();
        socket.close();
     } catch (IOException e) {
        e.printStackTrace();
  }
}
```

A UDP Client

```
import java.io.*;
import java.net.*;
import java.util.*;
public class QuoteClient {
  public static void main(String[] args) throws IOException {
     if (args.length != 1) {
        System.out.println("Usage: java QuoteClient < hostname>");
        return:
     }
     // get a datagram socket
     DatagramSocket socket = new DatagramSocket();
     // send request
     byte[] buf = new byte[256];
     InetAddress address = InetAddress.getByName(args[0]);
     DatagramPacket packet =
                     new DatagramPacket(buf, buf.length, address, 4445
     socket.send(packet);
     // get response
     packet = new DatagramPacket(buf, buf.length);
     socket.receive(packet);
     // display response
     String received = new String(packet.getData(), 0);
     System.out.println("Quote of the Moment: " + received);
     socket.close();
  }
}
```

A UDP Quote Server

```
import java.io.*;
public class QuoteServer {
  public static void main(String[] args) throws IOException {
     new QuoteServerThread().start();
}
import java.io.*;
import java.net.*;
import java.util.*;
public class QuoteServerThread extends Thread {
  protected DatagramSocket socket = null;
  protected BufferedReader in = null;
  protected boolean moreQuotes = true;
  public QuoteServerThread() throws IOException {
     this("QuoteServerThread");
  }
  public QuoteServerThread(String name) throws IOException {
     super(name);
     socket = new DatagramSocket(4445);
     try {
       in = new BufferedReader(new FileReader("one-liners.txt"));
     } catch (FileNotFoundException e) {
       System.err.println("Could not open quote file. Serving time
instead.");
  }
```

```
public void run() {
  while (moreQuotes) {
     try {
        byte[] buf = new byte[256];
           // receive request
        DatagramPacket packet = new DatagramPacket(buf, buf.length);
        socket.receive(packet);
           // figure out response
        String dString = null;
        if (in == null)
          dString = new Date().toString();
        else
          dString = getNextQuote();
        buf = dString.getBytes();
            // send the response to the client at "address" and "port"
        InetAddress address = packet.getAddress();
        int port = packet.getPort();
        packet = new DatagramPacket(buf, buf.length, address, port);
        socket.send(packet);
     } catch (IOException e) {
        e.printStackTrace();
          moreQuotes = false;
     }
  socket.close();
}
protected String getNextQuote() {
  String returnValue = null;
  try {
     if ((returnValue = in.readLine()) == null) {
        in.close();
          moreQuotes = false;
        returnValue = "No more quotes. Goodbye.";
  } catch (IOException e) {
     returnValue = "IOException occurred in server.";
  return return Value:
}
```

}