

# CS 471 Operating Systems

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# Announcements

- HW1 posted yesterday on BB
  - Due end of Friday 09/20
- OS/161 PA1 (Synchronization) will be posted on BB this Thursday 09/19
  - Due end of day 10/18

# Review: Condition Variables

# Condition Variables

- CV: **Queue** of sleeping threads
- Threads add themselves to the **queue** with **wait**
- Threads wake up threads on the **queue** with **signal**



# Condition Variables

- `cond_wait(cond_t *cv, mutex_t *lock)`
  - assume the lock is held when `cond_wait()` is called
  - puts caller to sleep + **release** the lock (**atomically**)
  - when awoken, **reacquires** lock before returning
- `cond_signal(cond_t *cv)`
  - wake a **single** waiting thread (if  $\geq 1$  thread is waiting)
  - if there is no waiting thread, just return, doing nothing

# Review: Join

```
void thread_exit() {  
    Mutex_lock(&m);           // a  
    Cond_signal(&c);         // b  
    Mutex_unlock(&m);        // c  
}
```

```
void thread_join() {  
    Mutex_lock(&m);           // x  
    Cond_wait(&c, &m);       // y  
    Mutex_unlock(&m);        // z  
}
```

# Review: Join

Parent: x y z

Child: a b c

```
void thread_exit() {  
    Mutex_lock(&m);           // a  
    Cond_signal(&c);          // b  
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void thread_join() {  
    Mutex_lock(&m);           // x  
    Cond_wait(&c, &m);        // y  
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}
```

# Review: Join

Parent: x y z

Child: a b c

**GOOD!**

```
void thread_exit() {  
    Mutex_lock(&m);           // a  
    Cond_signal(&c);         // b  
    Mutex_unlock(&m);        // c  
}
```

```
void thread_join() {  
    Mutex_lock(&m);           // x  
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# Review: Join

Parent:                   x   y

Child:    a   b   c

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void thread_exit() {  
    Mutex_lock(&m);           // a  
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    Mutex_lock(&m);           // x  
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    Mutex_unlock(&m);        // z  
}
```

# Review: Join

Parent:                    x   y   ... *sleeeeeeeeeep forever* ...

Child:     a   b   c

```
void thread_exit() {  
    Mutex_lock(&m);           // a  
    Cond_signal(&c);         // b  
    Mutex_unlock(&m);        // c  
}
```

```
void thread_join() {  
    Mutex_lock(&m);           // x  
    Cond_wait(&c, &m);       // y  
    Mutex_unlock(&m);        // z  
}
```

# Review: Join

```
void thread_exit() {  
    done = 1;           // a  
    Cond_signal(&c);    // b  
}
```

```
void thread_join() {  
    Mutex_lock(&m);     // w  
    if (done == 0)     // x  
        Cond_wait(&c, &m); // y  
    Mutex_unlock(&m);  // z  
}
```



# Review: Join

Parent: w x y

Child: a b

```
void thread_exit() {  
    done = 1;           // a  
    Cond_signal(&c);    // b  
}
```

```
void thread_join() {  
    Mutex_lock(&m);     // w  
    if (done == 0)     // x  
        Cond_wait(&c, &m); // y  
    Mutex_unlock(&m);  // z  
}
```

# Review: Join

Parent: w x y ... *sleeeeeeeeeeep forever* ...

Child: a b

```
void thread_exit() {  
    done = 1;           // a  
    Cond_signal(&c);    // b  
}
```

```
void thread_join() {  
    Mutex_lock(&m);     // w  
    if (done == 0)     // x  
        Cond_wait(&c, &m); // y  
    Mutex_unlock(&m);  // z  
}
```

# Review: Join

Parent: w x y ... *sleeeeeeeeeeep forever* ...

Child: a b

```
void thread_exit() {  
    done = 1;           // a  
    Cond_signal(&c);    // b  
}  
  
void thread_join() {  
    Mutex_lock(&m);     // w  
    if (done == 0)     // x  
        Cond_wait(&c, &m); // y  
    Mutex_unlock(&m);  // z  
}
```

How to fix?

# Review: Join

Parent: w x y ... *sleeeeeeeeeeep forever* ...

Child: a b

```
void thread_exit() {  
    done = 1;           // a  
    Cond_signal(&c);    // b  
}  
    Mutex_unlock(&m);  
    }  
void thread_join() {  
    Mutex_lock(&m);     // w  
    while if (done == 0) // x  
        Cond_wait(&c, &m); // y  
    Mutex_unlock(&m);   // z  
}
```

# Good Rule of Thumb When Using CV

Always do **wait** and **signal while holding the lock**

# Good Rule of Thumb When Using CV

Always do **wait** and **signal while holding the lock**

Why? To prevent lost signals.

# Classical Problems of Synchronization

- Producer-Consumer Problem
  - CV-based version
- Readers-Writers Problem
- Dining-Philosophers Problem

# CV-based Producer-Consumer Implementation 1

## Single CV and if statement

```
cond_t  cond;
mutex_t mutex;

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        if (count == 1)                       // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                               // p4
        Pthread_cond_signal(&cond);          // p5
        Pthread_mutex_unlock(&mutex);        // p6
    }
}

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        if (count == 0)                       // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}
```

```
1  int buffer;
2  int count = 0; // initially, empty
3
4  void put(int value) {
5      assert(count == 0);
6      count = 1;
7      buffer = value;
8  }
9
10 int get() {
11     assert(count == 1);
12     count = 0;
13     return buffer;
14 }
```

Put and Get routines  
**Single buffer**



# CV-based Producer-Consumer Implementation 1

## Single CV and if statement

```
cond_t cond;
mutex_t mutex;

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        if (count == 1)                       // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                                // p4
        Pthread_cond_signal(&cond);           // p5
        Pthread_mutex_unlock(&mutex);         // p6
    }
}

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        if (count == 0)                       // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                       // c4
        Pthread_cond_signal(&cond);           // c5
        Pthread_mutex_unlock(&mutex);         // c6
        printf("%d\n", tmp);
    }
}
```


```
1  int buffer;
2  int count = 0; // initially, empty
3
4  void put(int value) {
5      assert(count == 0);
6      count = 1;
7      buffer = value;
8  }
9
10 int get() {
11     assert(count == 1);
12     count = 0;
13     return buffer;
14 }
```

## Put and Get routines Single buffer

**What's the problem of this approach?**

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) {  C1 running
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get

# CV-based Producer-Consumer Implementation 1

← P running

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}
```

```
void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}
```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	Nothing to get
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
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        printf("%d\n", tmp);
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```

```

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        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

← P running

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

```

void *producer(void *arg) {
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    }
}

```

← P running

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	$T_{c1}$ awoken



# CV-based Producer-Consumer Implementation 1

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
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        printf("%d\n", tmp);
    }
}
```

```
void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}
```

← P running

T <sub>c1</sub>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	T <sub>c1</sub> awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	

# CV-based Producer-Consumer Implementation 1

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}
```

```
void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}
```

← P running

T <sub>c1</sub>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	
	Ready		Ready	p5	Running	1	Buffer now full T <sub>c1</sub> awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) { ← C1 runnable
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	$T_{c1}$ awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep



# CV-based Producer-Consumer Implementation 1

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}
```

← C2 running

```
void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}
```

T <sub>c1</sub>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	T <sub>c1</sub> awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	T <sub>c2</sub> sneaks in ...

# CV-based Producer-Consumer Implementation 1

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}
```

← C2 running

```
void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}
```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	$T_{c1}$ awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	$T_{c2}$ sneaks in ...
	Ready	c2	Running		Sleep	1	
	Ready	c4	Running		Sleep	0	... and grabs data

# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) { ← C2 running
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	$T_{c1}$ awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	$T_{c2}$ sneaks in ...
	Ready	c2	Running		Sleep	1	
	Ready	c4	Running		Sleep	0	... and grabs data
	Ready	c5	Running		Ready	0	$T_p$ awoken



# CV-based Producer-Consumer Implementation 1

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        if (count == 0)                       // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        if (count == 1)                      // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);          // p5
        Pthread_mutex_unlock(&mutex);        // p6
    }
}

```

← C2 running

T <sub>c1</sub>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	T <sub>c1</sub> awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	T <sub>c2</sub> sneaks in ...
	Ready	c2	Running		Sleep	1	
	Ready	c4	Running		Sleep	0	... and grabs data
	Ready	c5	Running		Ready	0	T <sub>p</sub> awoken
	Ready	c6	Running		Ready	0	

# CV-based Producer-Consumer Implementation 1

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        if (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}
```

← C1 running

```
void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        if (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}
```

<u>T<sub>c1</sub></u>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep		Ready	p1	Running	0	
	Sleep		Ready	p2	Running	0	
	Sleep		Ready	p4	Running	1	Buffer now full
	Ready		Ready	p5	Running	1	T <sub>c1</sub> awoken
	Ready		Ready	p6	Running	1	
	Ready		Ready	p1	Running	1	
	Ready		Ready	p2	Running	1	
	Ready		Ready	p3	Sleep	1	Buffer full; sleep
	Ready	c1	Running		Sleep	1	T <sub>c2</sub> sneaks in ...
	Ready	c2	Running		Sleep	1	
	Ready	c4	Running		Sleep	0	... and grabs data
	Ready	c5	Running		Ready	0	T <sub>p</sub> awoken
	Ready	c6	Running		Ready	0	
c4	Running		Ready		Ready	0	Oh oh! No data

# CV-based Producer-Consumer Implementation 2

## Single CV and while

```
1  cond_t  cond;
2  mutex_t mutex;
3
4  void *producer(void *arg) {
5      int i;
6      for (i = 0; i < loops; i++) {
7          Pthread_mutex_lock(&mutex);           // p1
8          while (count == 1)                   // p2
9              Pthread_cond_wait(&cond, &mutex); // p3
10         put(i);                               // p4
11         Pthread_cond_signal(&cond);          // p5
12         Pthread_mutex_unlock(&mutex);        // p6
13     }
14 }
15
16 void *consumer(void *arg) {
17     int i;
18     for (i = 0; i < loops; i++) {
19         Pthread_mutex_lock(&mutex);           // c1
20         while (count == 0)                   // c2
21             Pthread_cond_wait(&cond, &mutex); // c3
22         int tmp = get();                      // c4
23         Pthread_cond_signal(&cond);          // c5
24         Pthread_mutex_unlock(&mutex);        // c6
25         printf("%d\n", tmp);
26     }
27 }
```

# CV-based Producer-Consumer Implementation 2

## Single CV and while

**What's the problem of this approach?**

```
1  cond_t  cond;
2  mutex_t mutex;
3
4  void *producer(void *arg) {
5      int i;
6      for (i = 0; i < loops; i++) {
7          Pthread_mutex_lock(&mutex);           // p1
8          while (count == 1)                   // p2
9              Pthread_cond_wait(&cond, &mutex); // p3
10         put(i);                               // p4
11         Pthread_cond_signal(&cond);          // p5
12         Pthread_mutex_unlock(&mutex);       // p6
13     }
14 }
15
16 void *consumer(void *arg) {
17     int i;
18     for (i = 0; i < loops; i++) {
19         Pthread_mutex_lock(&mutex);           // c1
20         while (count == 0)                   // c2
21             Pthread_cond_wait(&cond, &mutex); // c3
22         int tmp = get();                      // c4
23         Pthread_cond_signal(&cond);          // c5
24         Pthread_mutex_unlock(&mutex);       // c6
25         printf("%d\n", tmp);
26     }
27 }
```



```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

← **C1 running**

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	Nothing to get
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	



## C2 running



```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get

P running



```

void *consumer(void *arg) {
  int i;
  for (i = 0; i < loops; i++) {
    Pthread_mutex_lock(&mutex);           // c1
    while (count == 0)                    // c2
      Pthread_cond_wait(&cond, &mutex); // c3
    int tmp = get();                       // c4
    Pthread_cond_signal(&cond);           // c5
    Pthread_mutex_unlock(&mutex);         // c6
    printf("%d\n", tmp);
  }
}

```

```

void *producer(void *arg) {
  int i;
  for (i = 0; i < loops; i++) {
    Pthread_mutex_lock(&mutex);           // p1
    while (count == 1)                    // p2
      Pthread_cond_wait(&cond, &mutex); // p3
    put(i);                               // p4
    Pthread_cond_signal(&cond);           // p5
    Pthread_mutex_unlock(&mutex);         // p6
  }
}

```

T <sub>c1</sub>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full

P running



```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                   // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
        Pthread_cond_signal(&cond);         // c5
        Pthread_mutex_unlock(&mutex);       // c6
        printf("%d\n", tmp);
    }
}

```

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        while (count == 1)                   // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);         // p5
        Pthread_mutex_unlock(&mutex);       // p6
    }
}

```

T <sub>c1</sub>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	T <sub>c1</sub> awoken

P sleeping



```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                   // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}
```

```
void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        while (count == 1)                   // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);          // p5
        Pthread_mutex_unlock(&mutex);        // p6
    }
}
```

T <sub>c1</sub>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	T <sub>c1</sub> awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)



```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // c1
        while (count == 0) // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get(); // c4
        Pthread_cond_signal(&cond); // c5
        Pthread_mutex_unlock(&mutex); // c6
        printf("%d\n", tmp);
    }
}

```

← C1 running

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex); // p1
        while (count == 1) // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i); // p4
        Pthread_cond_signal(&cond); // p5
        Pthread_mutex_unlock(&mutex); // p6
    }
}

```

T <sub>c1</sub>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	T <sub>c1</sub> awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                    // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}

```

 **C1 running**

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        while (count == 1)                   // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);          // p5
        Pthread_mutex_unlock(&mutex);        // p6
    }
}

```

T <sub>c1</sub>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	T <sub>c1</sub> awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	T <sub>c1</sub> grabs data

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                    // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}

```

 **C1 running**

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        while (count == 1)                   // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);          // p5
        Pthread_mutex_unlock(&mutex);        // p6
    }
}

```

T <sub>c1</sub>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	T <sub>c1</sub> awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	T <sub>c1</sub> grabs data
c5	Running		Ready		Sleep	0	<b>Oops! Woke T<sub>c2</sub></b>





# C1 sleeping

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                   // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
        Pthread_cond_signal(&cond);         // c5
        Pthread_mutex_unlock(&mutex);       // c6
        printf("%d\n", tmp);
    }
}

```

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        while (count == 1)                   // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                             // p4
        Pthread_cond_signal(&cond);         // p5
        Pthread_mutex_unlock(&mutex);       // p6
    }
}

```

T <sub>c1</sub>	State	T <sub>c2</sub>	State	T <sub>p</sub>	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	T <sub>c1</sub> awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	T <sub>c1</sub> grabs data
c5	Running		Ready		Sleep	0	Oops! Woke T <sub>c2</sub>
c6	Running		Ready		Sleep	0	
c1	Running		Ready		Sleep	0	
c2	Running		Ready		Sleep	0	
c3	Sleep		Ready		Sleep	0	Nothing to get



 **C2 running**

```
void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                   // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                      // c4
        Pthread_cond_signal(&cond);          // c5
        Pthread_mutex_unlock(&mutex);        // c6
        printf("%d\n", tmp);
    }
}
```

```
void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        while (count == 1)                   // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);          // p5
        Pthread_mutex_unlock(&mutex);        // p6
    }
}
```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	$T_{c1}$ awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	$T_{c1}$ grabs data
c5	Running		Ready		Sleep	0	Oops! Woke $T_{c2}$
c6	Running		Ready		Sleep	0	
c1	Running		Ready		Sleep	0	
c2	Running		Ready		Sleep	0	
c3	Sleep		Ready		Sleep	0	Nothing to get
	Sleep	c2	Running		Sleep	0	
	Sleep	c3	Sleep		Sleep	0	Everyone asleep...

## C2 sleeping

```

void *consumer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // c1
        while (count == 0)                   // c2
            Pthread_cond_wait(&cond, &mutex); // c3
        int tmp = get();                     // c4
        Pthread_cond_signal(&cond);         // c5
        Pthread_mutex_unlock(&mutex);       // c6
        printf("%d\n", tmp);
    }
}

```

```

void *producer(void *arg) {
    int i;
    for (i = 0; i < loops; i++) {
        Pthread_mutex_lock(&mutex);           // p1
        while (count == 1)                   // p2
            Pthread_cond_wait(&cond, &mutex); // p3
        put(i);                              // p4
        Pthread_cond_signal(&cond);         // p5
        Pthread_mutex_unlock(&mutex);       // p6
    }
}

```

$T_{c1}$	State	$T_{c2}$	State	$T_p$	State	Count	Comment
c1	Running		Ready		Ready	0	
c2	Running		Ready		Ready	0	
c3	Sleep		Ready		Ready	0	Nothing to get
	Sleep	c1	Running		Ready	0	
	Sleep	c2	Running		Ready	0	
	Sleep	c3	Sleep		Ready	0	Nothing to get
	Sleep		Sleep	p1	Running	0	
	Sleep		Sleep	p2	Running	0	
	Sleep		Sleep	p4	Running	1	Buffer now full
	Ready		Sleep	p5	Running	1	$T_{c1}$ awoken
	Ready		Sleep	p6	Running	1	
	Ready		Sleep	p1	Running	1	
	Ready		Sleep	p2	Running	1	
	Ready		Sleep	p3	Sleep	1	Must sleep (full)
c2	Running		Sleep		Sleep	1	Recheck condition
c4	Running		Sleep		Sleep	0	$T_{c1}$ grabs data
c5	Running		Ready		Sleep	0	Oops! Woke $T_{c2}$
c6	Running		Ready		Sleep	0	
c1	Running		Ready		Sleep	0	
c2	Running		Ready		Sleep	0	
c3	Sleep		Ready		Sleep	0	Nothing to get
	Sleep	c2	Running		Sleep	0	
	Sleep	c3	Sleep		Sleep	0	Everyone asleep...

# CV-based Producer-Consumer Implementation 3

```
1  cond_t empty, fill;
2  mutex_t mutex;
3
4  void *producer(void *arg) {
5      int i;
6      for (i = 0; i < loops; i++) {
7          Pthread_mutex_lock(&mutex);
8          while (count == 1)
9              Pthread_cond_wait(&empty, &mutex);
10         put(i);
11         Pthread_cond_signal(&fill);
12         Pthread_mutex_unlock(&mutex);
13     }
14 }
15
16 void *consumer(void *arg) {
17     int i;
18     for (i = 0; i < loops; i++) {
19         Pthread_mutex_lock(&mutex);
20         while (count == 0)
21             Pthread_cond_wait(&fill, &mutex);
22         int tmp = get();
23         Pthread_cond_signal(&empty);
24         Pthread_mutex_unlock(&mutex);
25         printf("%d\n", tmp);
26     }
27 }
```

**Two CVs and while**



# CV-based Producer-Consumer Implementation 3

```
1  cond_t empty, fill;
2  mutex_t mutex;
3
4  void *producer(void *arg) {
5      int i;
6      for (i = 0; i < loops; i++) {
7          Pthread_mutex_lock(&mutex);
8          while (count == 1)
9              Pthread_cond_wait(&empty, &mutex);
10         put(i);
11         Pthread_cond_signal(&fill);
12         Pthread_mutex_unlock(&mutex);
13     }
14 }
15
16 void *consumer(void *arg) {
17     int i;
18     for (i = 0; i < loops; i++) {
19         Pthread_mutex_lock(&mutex);
20         while (count == 0)
21             Pthread_cond_wait(&fill, &mutex);
22         int tmp = get();
23         Pthread_cond_signal(&empty);
24         Pthread_mutex_unlock(&mutex);
25         printf("%d\n", tmp);
26     }
27 }
```

## Two CVs and while

Using **two CVs** to distinguish two types of threads; in order to properly signal which thread should wake up

- Producer waits on **empty**
- Consumer waits on **full**

# Readers-Writers Problem

# Readers-Writers Problem

- A data object (e.g. a file) is to be shared among several concurrent processes/threads
- A **writer** process/thread must have exclusive access to the data object
- **Multiple reader** processes/threads may access the shared data simultaneously without a problem

# Reader-Writer Lock

```
1  typedef struct _rwlock_t {
2      sem_t lock;          // binary semaphore (basic lock)
3      sem_t writelock;    // used to allow ONE writer or MANY readers
4      int  readers;       // count of readers reading in critical section
5  } rwlock_t;
6
7  void rwlock_init(rwlock_t *rw) {
8      rw->readers = 0;
9      sem_init(&rw->lock, 0, 1);
10     sem_init(&rw->writelock, 0, 1);
11 }
12
13 void rwlock_acquire_readlock(rwlock_t *rw) {
14     sem_wait(&rw->lock);
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16     if (rw->readers == 1)
17         sem_wait(&rw->writelock); // first reader acquires writelock
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19 }
20
21 void rwlock_release_readlock(rwlock_t *rw) {
22     sem_wait(&rw->lock);
23     rw->readers--;
24     if (rw->readers == 0)
25         sem_post(&rw->writelock); // last reader releases writelock
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Initially, # readers is 0  
binary sem lock set to 1  
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35 }
```

Initially, # readers is 0  
binary sem lock set to 1  
writelock set to 1

**Writer cannot  
be in CS when  
readers are!**

# Readers-Writers Problem: Writer Thread

```
rwlock_acquire_writelock(rw);  
    ...  
    write is performed  
    ...  
rwlock_release_writelock(rw);
```

# Readers-Writers Problem: Reader Thread

```
rwlock_acquire_readlock(rw)
    ...
    read is performed
    ...
rwlock_release_readlock(rw)
```

**Well, is this solution Okay?**

# Readers-Writers Problem: Reader Thread

```
rwlock_acquire_readlock(rw)
    ...
    read is performed
    ...
rwlock_release_readlock(rw)
```

**Well, is this solution Okay?**

A: Technically it works. But **starvation** may happen

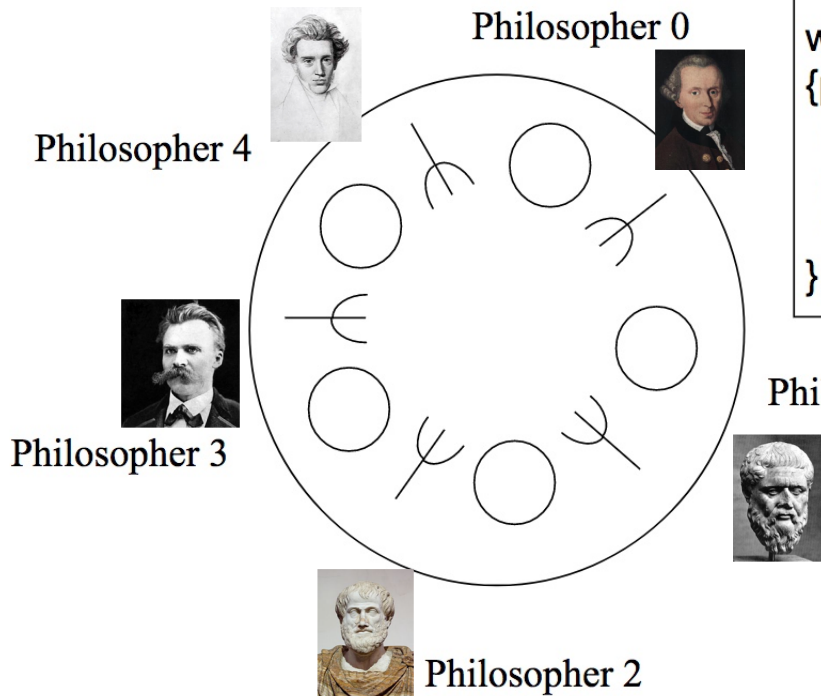


# Starvation

- A process/thread that is forced to wait **indefinitely** in a synchronization program is said to be subject to **starvation**
  - In some execution scenarios, that process does not make any progress
  - **Deadlocks imply starvation, but the reverse is not true**

# Dining-Philosophers Problem

# Dining-Philosophers Problem



```
while(food available)
{pick up 2 adj. forks;
 eat;
 put down forks;
 think awhile;
}
```

- 5 philosophers share a common circular table. There are 5 forks (or chopsticks) and food (in the middle). When a philosopher gets hungry, he tries to pick up the closest forks
- A philosopher may pick up only one fork at a time, and cannot pick up a fork already in use. When done, he puts down both of his forks, one after the other

Shared data

```
sem_t forks[5];
```

Initially all semaphore values are 1

# Dining-Philosophers Problem

- The basic loop of a philosopher

```
while (1) {  
    think();  
    getforks();  
    eat();  
    putforks();  
}
```

Annotations for the code:

- Red box around `getforks();` with a red arrow pointing to `??`
- Blue arrow pointing from `eat();` to **Critical section**
- Red box around `putforks();` with a red arrow pointing to `??`

# The Helper Functions

```
int left(int p) { return p; }
int right(int p) { return (p + 1) % 5; }
```

`sem_t forks[5]`

- Each fork initialized to **1**

```
1 void getforks() {
2     sem_wait(forks[left(p)]);
3     sem_wait(forks[right(p)]);
4 }
5
6 void putforks() {
7     sem_post(forks[left(p)]);
8     sem_post(forks[right(p)]);
9 }
```

**Is this solution correct?**

# Simplest Example of A Deadlock

Thread 0

```
sem_wait(fork[0])  
sem_wait(fork[1])  
sem_signal(fork[0])  
sem_signal(fork[1])
```

Interleaving

Thread 1

```
sem_wait(fork[1])  
sem_wait(fork[0])  
sem_signal(fork[1])  
sem_signal(fork[0])
```



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sem_wait(fork[0])  
sem_signal(fork[1])  
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```
sem_wait(fork[0])  
sem_wait(fork[1])  
sem_signal(fork[0])  
sem_signal(fork[1])
```

Interleaving

```
sem_wait(fork[0])  
  
sem_wait(fork[1])  
  
sem_wait(fork[0])  
wait...  
  
sem_wait(fork[1])
```

Thread 1

```
sem_wait(fork[1])  
sem_wait(fork[0])  
sem_signal(fork[1])  
sem_signal(fork[0])
```

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```

Thread 1

```
sem_wait(fork[1])  
sem_wait(fork[0])  
sem_signal(fork[1])  
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```

# Simplest Example of A Deadlock

**Q: Would the previous 5DP implementation cause exactly the same form of a deadlock as shown below?**

Thread 0

```
sem_wait(fork[0])  
sem_wait(fork[1])  
sem_signal(fork[0])  
sem_signal(fork[1])
```

Interleaving

```
sem_wait(fork[0])  
  
sem_wait(fork[1])  
  
sem_wait(fork[0])  
wait...  
  
sem_wait(fork[1])  
wait...
```

Thread 1

```
sem_wait(fork[1])  
sem_wait(fork[0])  
sem_signal(fork[1])  
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```



# Review: Conditions for Deadlocks

- Mutually exclusive access of shared resources
  - Binary semaphore `fork[0]` and `fork[1]`

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  - Thread 0 waits for Thread 1 to `signal(fork[1])` and
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  - Thread 1 waits for Thread 0 to `signal(fork[0])`
- Hold and wait
  - Holding either `fork[0]` or `fork[1]` while waiting on the other
- No preemption
  - Neither `fork[0]` and `fork[1]` can be removed from their respective holding threads

# Why 5DP is Interesting?

- How to eat with your fellows without causing deadlocks
  - Circular arguments (the **circular wait condition**)
  - Not giving up on firmly held things (**no preemption**)
  - Infinite patience with half-baked schemes (**hold some & wait for more**)
- Why starvation exists and what we can do about it?

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- ~~How to eat with your fellows without causing deadlocks~~ — **How to mess with your fellows!**
  - Circular arguments (the **circular wait condition**)
  - Not giving up on firmly held things (**no preemption**)
  - Infinite patience with half-baked schemes (**hold some & wait for more**)
- Why starvation exists and what we can do about it?



# Dijkstra's Solution:

## Break the Circular Wait Condition

- Change how forks are acquired by at least one of the philosophers
- Assume P0 – P4, 4 is the highest number

```
1 void getforks() {
2     if (p == 4) {
3         sem_wait(forks[right(p)]);
4         sem_wait(forks[left(p)]);
5     } else {
6         sem_wait(forks[left(p)]);
7         sem_wait(forks[right(p)]);
8     }
9 }
```

# Again, Starvation

- Subtle difference between deadlock and starvation
  - Once a set of processes are in a deadlock, there is **no future execution sequence** that can get them out of it!
  - In starvation, there does exist **hope** – some execution order may be favorable to the starving process although no guarantee it would ever occur
  - Rollback and retry are prone to starvation
  - Continuous arrival of higher priority process is another common starvation situation

# Building a Semaphore w/ CV Worksheet