CS 471 Operating Systems

Yue Cheng

George Mason University
Spring 2019

Introduction

- Instructor of Section 001
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 - Office: 5324 Engineering
 - Office hours: M 10:30am-12:30pm
 - Research interests: Distributed and storage systems, serverless and cloud computing, operating systems

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- Teaching assistant
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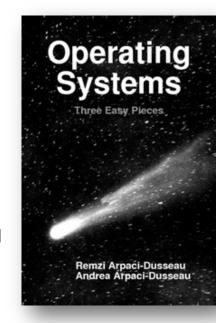


Administrivia

- Required textbook
 - Operating Systems: Three Easy Pieces,
 By Remzi H. Arpaci-Dusseau, Andrea C. Arpaci-Dusseau
- Recommended textbook
 - Operating Systems Principles & Practices
 By T. Anderson and M. Dahlin



- CS 310 Data Structures
- CS 367 Computer Systems & Programming -OR-
- ECE 445 Computer Organization
- Be comfortable with C programming language
- Class web page
 - https://cs.gmu.edu/~yuecheng/teaching/cs471_spring19/index.html
 - Class materials will all be available on the class web page



Administrivia (cont.)

Syllabus

https://cs.gmu.edu/media/syllabi/Spring2019/CS_471ChengY001.html

Grading

- 30% projects
- 10% homeworks
- 15%+15% two midterm exams
- 30% final exam

Reminders

- Honor code
- Late policy: 15% deducted each day. No credit after 3 days

Course format

- (Review) + lecture + (worksheets)
 - A short overview of the previous lecture to make sure the old content is not completely forgotten
 - Worksheet practices to make sure the lecture is well understood

OS/161 Projects

- Three coding projects
 - Proj 0: Intro to OS/161 (due Feb 15)
 - Proj 1: Synchronization (due Mar 22)
 - Proj 2: Syscalls and processes (due Apr 26)
 - Proj 3 (optional): Virtual memory (due May 10)

What is an OS?

What is an OS?

- OS manages resources
 - Memory, CPU, storage, network
 - Data (file systems, I/O)
- Provides low-level abstractions to applications
 - Files
 - Processes, threads
 - Virtual machines (VMs), containers

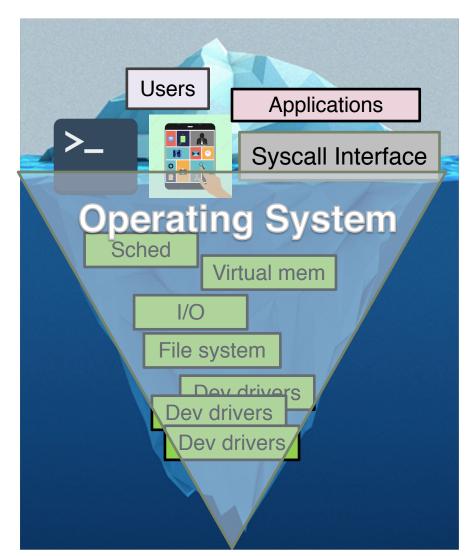
— ...

OS abstracts away low-level details



OS abstracts away low-level details

- User's perspective
 - User interface:
 - Terminal, GUI
 - Application interface:
 - System calls



The goals of an OS

- OS manages resources
 - Memory, CPU, storage, network
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 - **–** ...

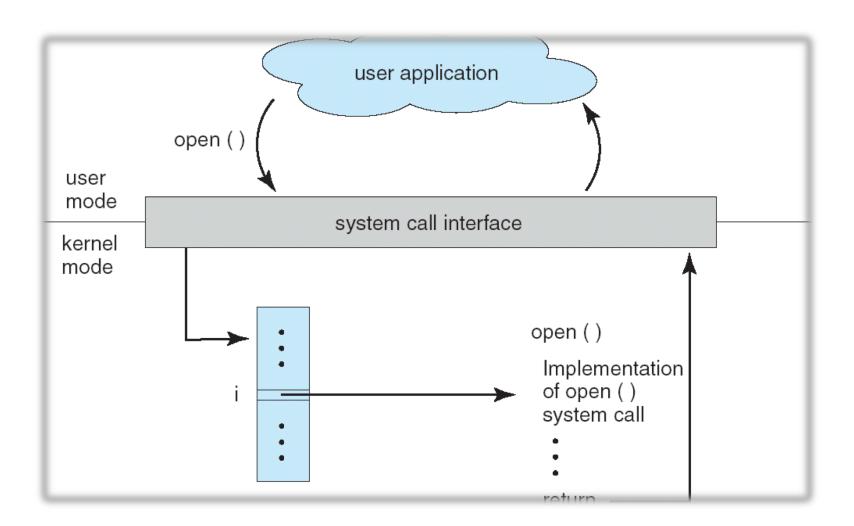
Goals

- Resource efficiency (resource virtualization)
- Easy-of-use (interfaces)
- Reliability (user-kernel space separation)

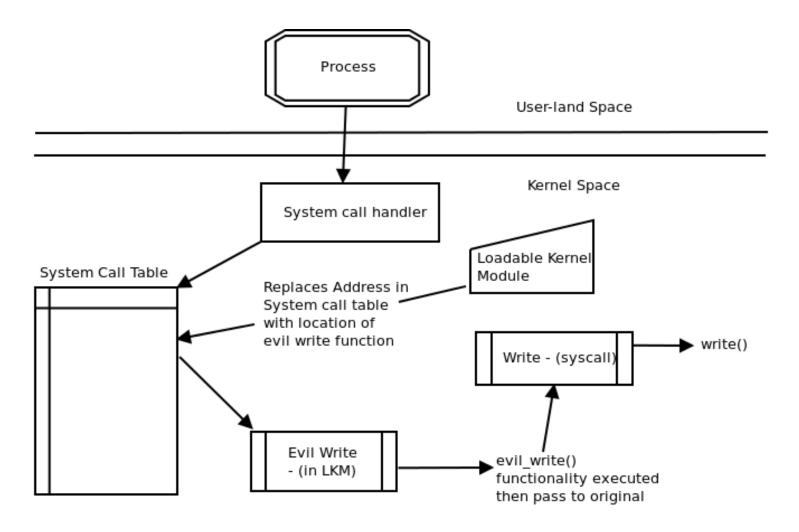
System Calls

- System calls provide the interface between a running program and the operating system
 - Generally available in routines written in C and C++
 - Certain low-level tasks may have to be written using assembly language
- Typically, application programmers design programs using an application programming interface (API)
- The run-time support system (run-time libraries) provides a system-call interface, that intercepts function calls in the API and invokes the necessary system call within the operating system
- Major differences in how they are implemented (e.g., Windows vs. Unix)

Example System Call Processing



Syscall: write(fd, buf, nbytes)

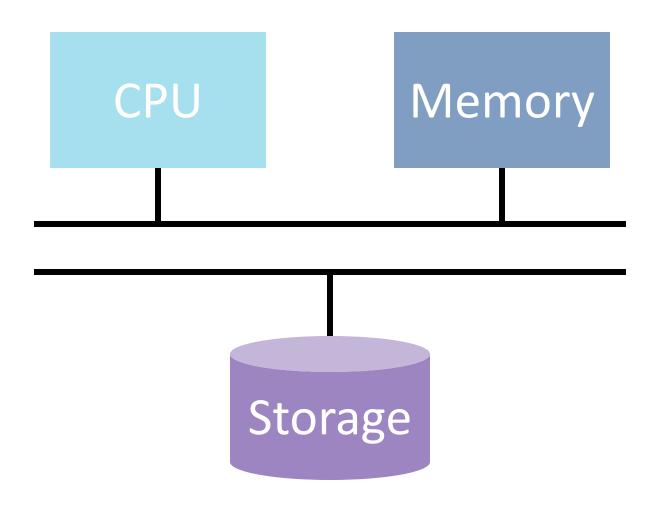


^{*}source: http://r00tkit.me/?p=46

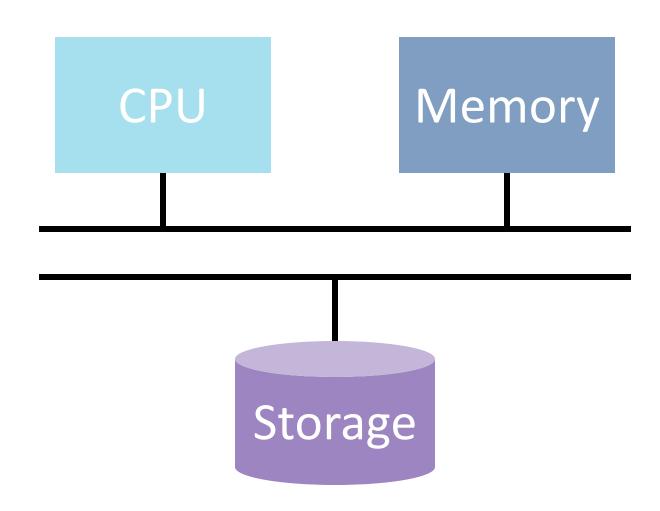
Major System Calls in Linux: File Management

- \circ fd = open(file, how, ...)
 - Open a file for reading, writing, or both
- \circ s = close(file)
 - Close an open file
- n = read(fd, buf, nbytes)
 - Read data from a file into a buffer
- n = write(fd, buf, nbytes)
 - Write data from a buffer into a file
- position = Iseek(fd, offset, whence)
 - Move the file pointer
- \circ s = stat(name, &buf)
 - Get a file's status info

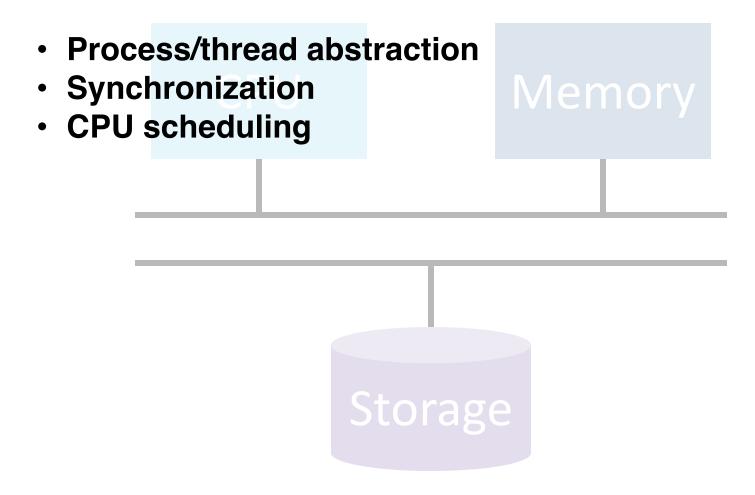
3 Major Topics



OS Provides Virtualization on Hardware



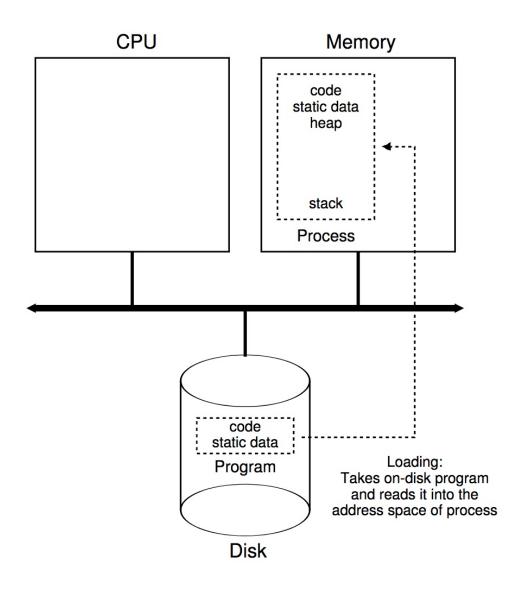
Topic 1: Concurrency, Synchronization, and CPU Scheduling



Process Abstraction

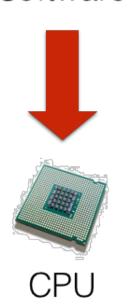
- A process is a program in execution
 - It is a unit of work within the system. A program is a passive entity, a process is an active entity.
- Process needs resources to accomplish its task
 - CPU, memory, I/O, files
 - Initialization data
- Process termination requires reclaim of any reusable resources
- Single-threaded process has one program counter specifying location of next instruction to execute
 - Process executes instructions sequentially, one at a time, until completion
- Multi-threaded process has one program counter per thread
- A software system may have many processes, some user, some operating system running concurrently on one or more CPUs
 - Concurrency by multiplexing the CPUs among the processes / threads

Loading from Program to Process

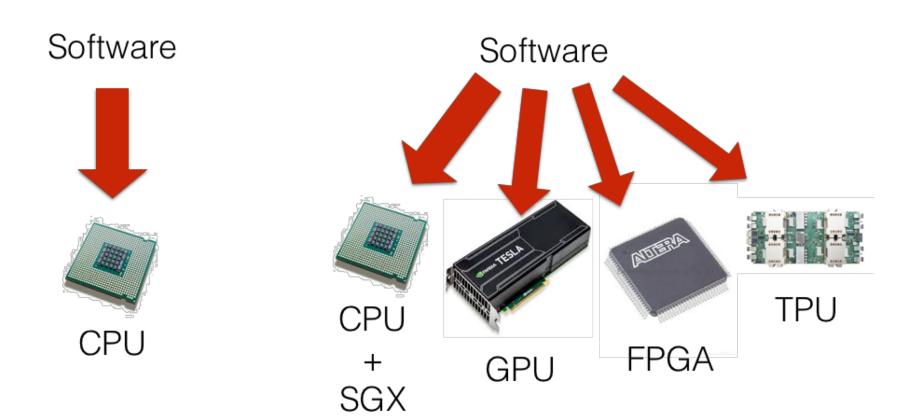


Computation – Increased Complexity

Software



Computation – Increased Complexity



Topic 2: Memory Management and Virtual Memory

Process/thread abstraction • Memory management Synchronization Virtual memory CPU scheduling

Memory Management

- All data in memory before and after processing
- All instructions in memory in order to execute
- Memory management determines what is in memory when
 - Optimizing CPU utilization and computer response to users
- Memory management activities
 - Keeping track of which parts of memory are currently being used and by whom
 - Deciding which processes (or parts thereof) and data to move into and out of memory
 - Allocating and deallocating memory space as needed
- Virtual memory management is an essential part of most operating systems

Topic 3: Storage, I/O, and Filesystems

- Process/thread abstraction Memory management
- Synchronization
- CPU scheduling

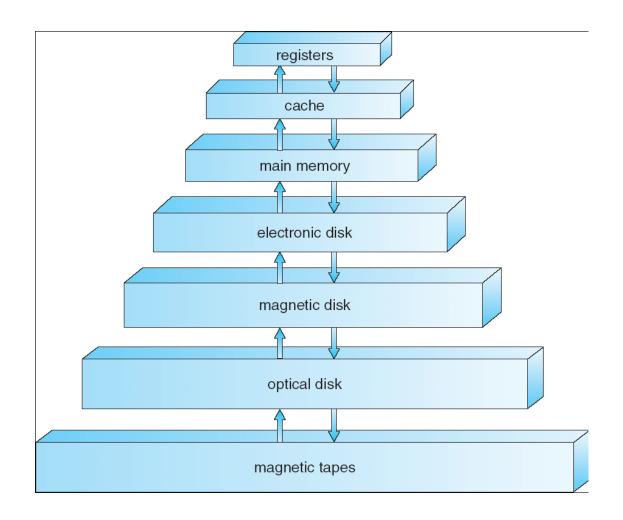
- Virtual memory

- Hard disk drives
- RAID
- Flash SSDs
- File and I/O systems

Storage Management

- OS provides a uniform, logical view of information storage
 - Abstracts physical properties to logical storage unit file
 - Each medium is controlled by device type (i.e., disk drive, tape drive)
 - Varying properties include access speed, capacity, datatransfer rate, access method (sequential or random)
- Filesystem management
 - Files usually organized into directories
 - Access control on most systems to determine who can access what
 - OS activities include
 - Creating and deleting files and directories
 - Primitives to manipulate files and dirs
 - Mapping files onto secondary storage
 - Backup files onto stable (non-volatile) storage media

Storage Hierarchy



Storage Structure

- Main memory relatively large storage media that the CPU can access directly
 - Small CPU cache memories are used to speed up average access time to the main memory at run-time
 - Volatile (data loss at power-off)
 - Byte-addressable
- Secondary storage extension of main memory that provides large nonvolatile storage capacity.
 - Magnetic disks
 - Electronic disks -- Solid state disks (SSDs)
 - Non-volatile (i.e., persistent)
 - Non byte-addressable

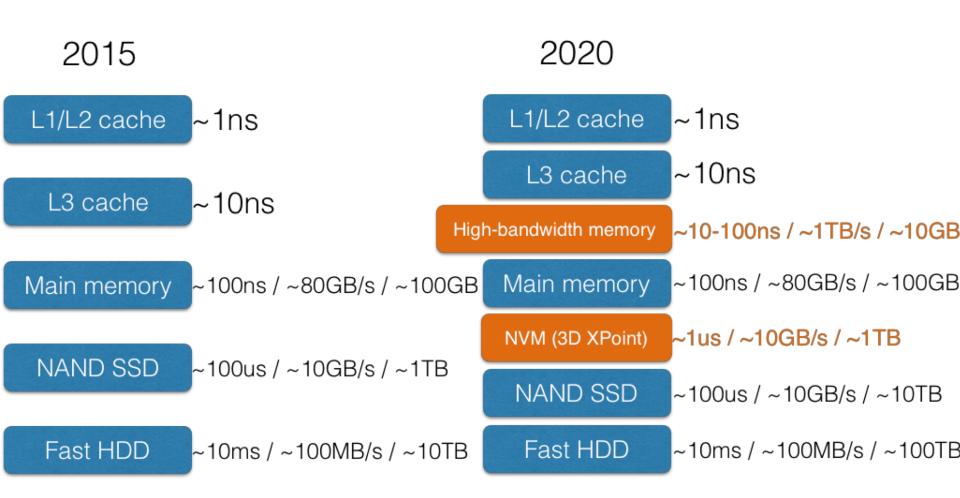
Storage Systems Tradeoffs

- Storage systems organized in hierarchy
 - Speed
 - Cost
 - Volatility
 - Density
- Faster access time, greater cost per bit
- Greater capacity (density), lower cost per bit
- Greater capacity (density), slower access speed

Storage hierarchy – Increased Complexity

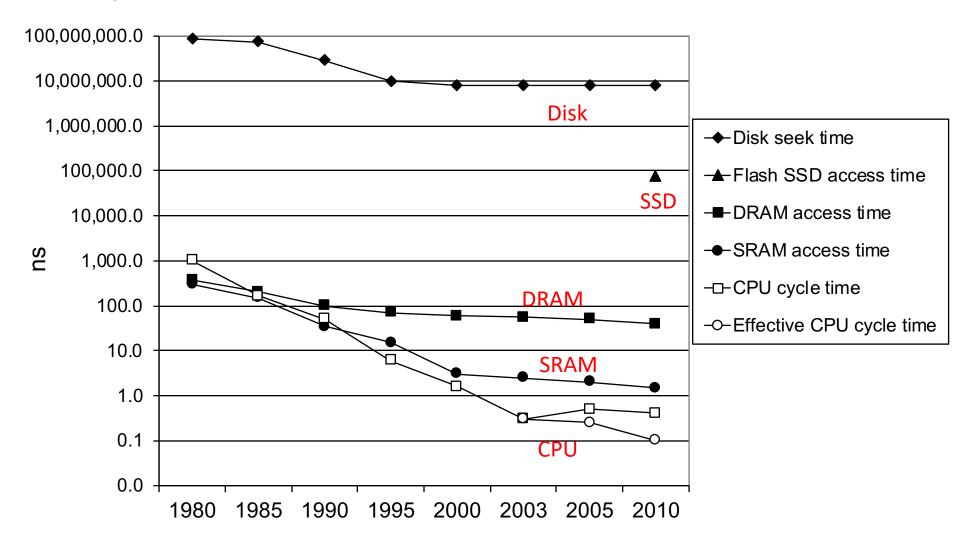
2015 L1/L2 cache ~1ns L3 cache ~10ns Main memory ~100ns / ~80GB/s / ~100GB NAND SSD ~100us / ~10GB/s / ~1TB Fast HDD ~10ms / ~100MB/s / ~10TB

Storage hierarchy – Increased Complexity



The CPU-Memory Gap

The gap widens between memory, disk, and CPU speeds.

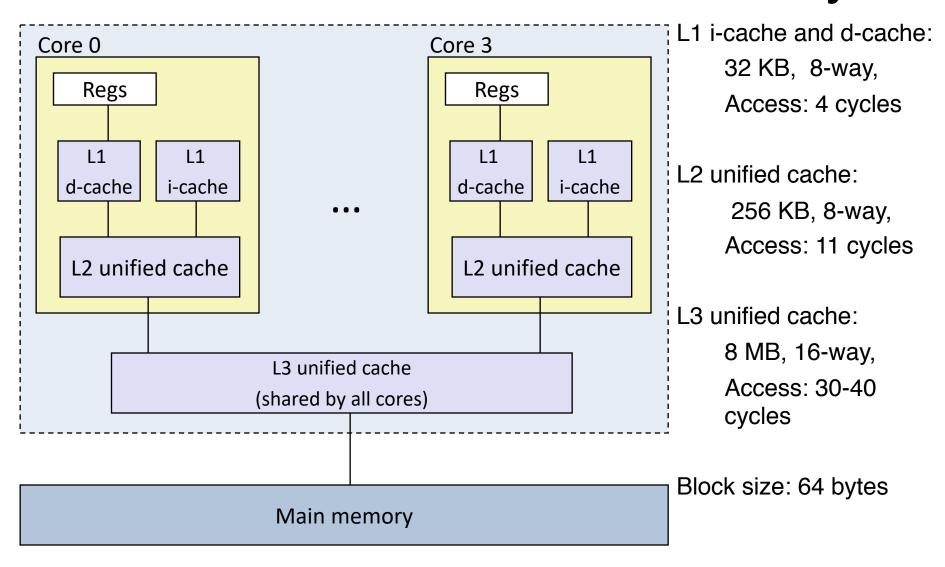


Data decades ago, but trends are the same

Caching

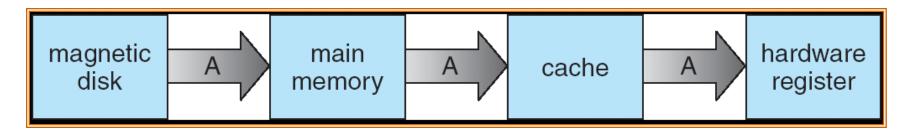
- Skew rule: 80% requests hit on 20% hottest data
- Important principle, performed at many levels in a computer (in hardware, operating system, software)
- Information in use copied from slower to faster storage temporarily
- Faster storage (cache) checked first to determine if information is there
 - If it is, information used directly from the cache (fast)
 - If not, data copied to cache and used there
- Cache smaller than storage being cached
 - Cache management important design problem
 - Cache size and replacement policy

Intel Core i7 Cache Hierarchy



Migration of Integer A from Disk to Register

 Multitasking environments must be careful to use most recent value, no matter where it is stored in the storage hierarchy



- Multiprocessor environment must provide cache coherency in hardware such that all CPUs have the most recent value in their cache
- Distributed environment situation even more complex
 - Several copies of a piece of data can exist

Advanced Topics (Miscellaneous)

Process/thread abstraction • Memory management
 Synchronization • Virtual memory
 CPU scheduling

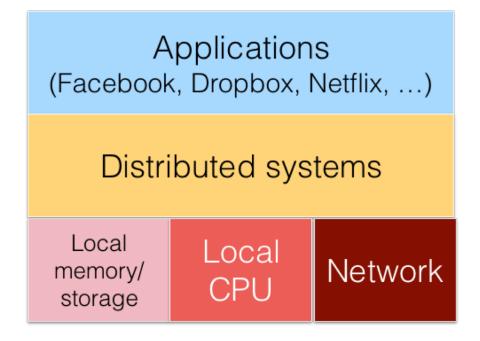
- Hard disk drives
- RAID
- Flash SSDs
- File and I/O systems

Distributed systems

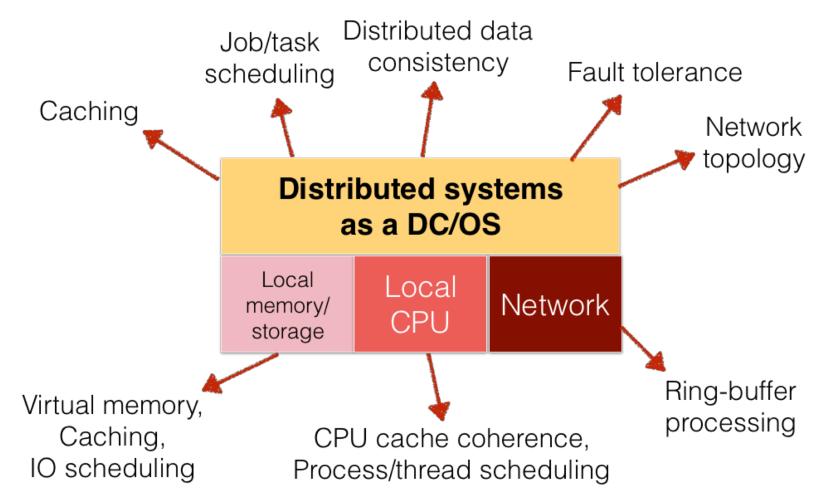
Distributed Systems as a DC/OS



Distributed Systems as a DC/OS



Distributed Systems as a DC/OS



Why do you take this course?

General Learning Goals

- Grasp basic knowledge about Operating Systems and Computer Systems software
- 2. Learn important systems concepts in general
 - Multi-processing/threading, synchronization
 - Scheduling
 - Caching, memory, storage
 - And more...
- Gain hands-on experience in writing/hacking/designing large systems software

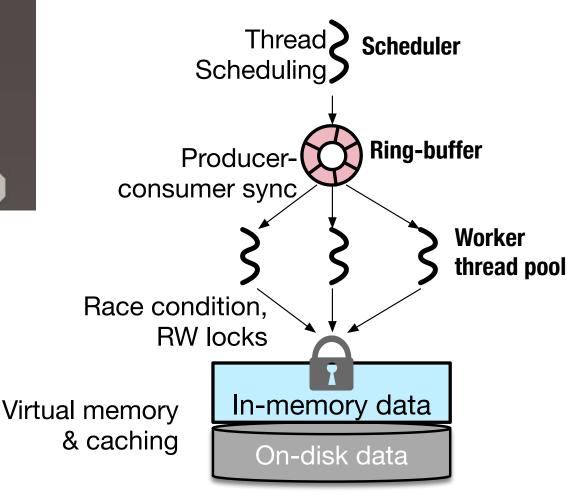
Why do you take this course?

- The OS concepts are everywhere
 - Fundamental OS techniques broadly generalize to widely-used systems technique
 - Scheduling
 - Concurrency
 - Memory management
 - Caching
 - ...

One example: Memcached



- Memcached is a distributed in-memory object cache system
 - Written in C
 - In-memory hash table
 - Multi-threading



Memcached can be treated as a user-space mini-OS

Homework #0

- Sign-up for Piazza
- Go over OS/161 Proj 0
- Fill out the Google Form for OS/161 project team composition

Next class...

- Intro of OS/161 (Proj 0)
- Process abstraction